

Vanilla Skype part 2

Fabrice DESCLAUX Kostya KORTCHINSKY

`serpilliere(at)droids-corp.org - fabrice.desclaux(at)eads.net`
`recca(at)rstack.org - kostya.kortchinsky(at)eads.net`
EADS Corporate Research Center — DCR/STI/C
SSI Lab
Suresnes, FRANCE

RECON2006, June 17th 2006



Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Introduction

Reverse engineering Skype

- Skype is a gold mine for reverse engineers
 - Binary protected against static and dynamic analysis
 - Almost everything is proprietary
 - Heavy use of cryptography
 - Binary loaded with hidden and undocumented features
- The work to carry out is far from easy

What to look for ?

- Find some ways to divert Skype from its original usage
 - Fun things to do with Skype
- Clarify some common beliefs
- Identify cryptographic flaws

Skype versions

A large variety of flavours...

- Skype v2.0.0.*
- PChome-Skype v2.0.1.*
- **TOM-Skype** v2.0.4.*
- livedoor-Skype v2.0.6.*
- Buffalo-Skype v2.0.7.*
- Daum-Skype v1.4.9.*
- HGC-Skype v2.0.10.*
- Onet-Skype v2.0.11.*
- Jubii-Skype v2.0.12.*
- eBay-Skype v2.0.13.*
- U3-Skype v1.4.14.*
- Maktoob-Skype v2.0.15.*
- Chinagate-Skype v2.0.16.*
- PacNet-Skype v2.0.17.*
- eBay.es-Skype v2.0.18.*
- eBay.it-Skype v2.0.19.*
- eBay.co.uk-Skype v2.0.20.*
- eBay.de-Skype v2.0.21.*
- eBay.fr-Skype v2.0.22.*
- Bebo-Skype v2.0.24.*
- eBay.nl-Skype v2.0.26.*
- **eBay.cn-Skype** v2.0.29.*

Downloading a particular version

<http://www.skype.com/go/getskype-⟨keyword⟩>

Example: <http://www.skype.com/go/getskype-ebay-fr>

Disclaimer

What Skype, Inc. does not tell you

A lot of "features" are silently fixed by Skype, Inc. with the numerous subversion updates that are published almost weekly. Since it is rather difficult to follow *everything*, some of the stuff described hereafter might not be totally accurate in the latest versions.

Outline

- 1 Introduction
- 2 Networking**
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

For P in packets: zip P

Packet compression

- Each packet can be compressed
- The algorithm used: arithmetic compression
- Zip would have been too easy 😊

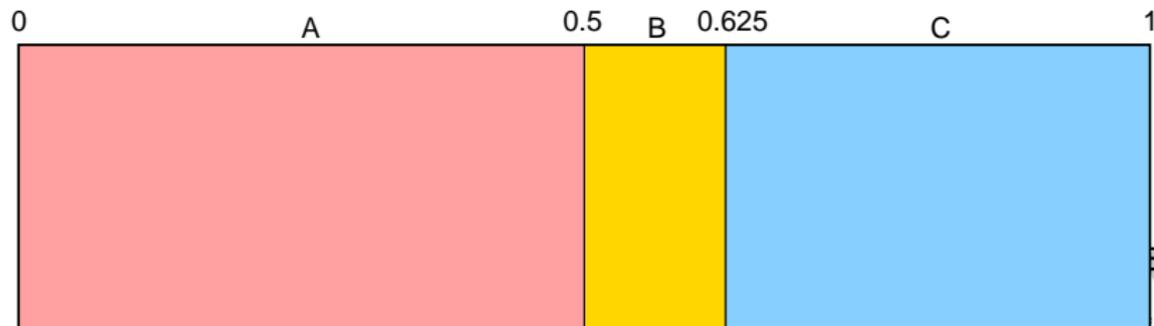
Principle

- Close to Huffman algorithm
- Reals are used instead of bits

Arithmetic compression

Example

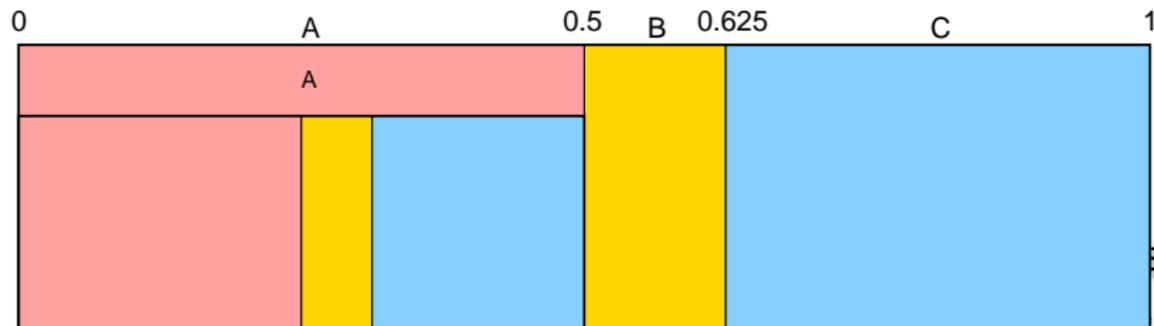
- $[0, 1]$ is split in subintervals for each symbol according to their frequency
- First symbol is A . We subdivide its interval
- Then comes C
- Then A again
- Then B
- Each real enclosed into this small interval can encode $ACAB$



Arithmetic compression

Example

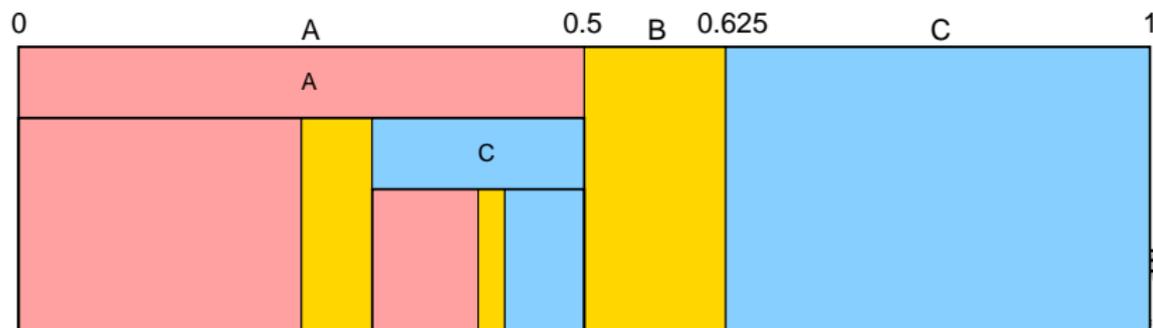
- $[0, 1]$ is split in subintervals for each symbol according to their frequency
- First symbol is A . We subdivide its interval
 - Then comes C
 - Then A again
 - Then B
 - Each real enclosed into this small interval can encode $ACAB$



Arithmetic compression

Example

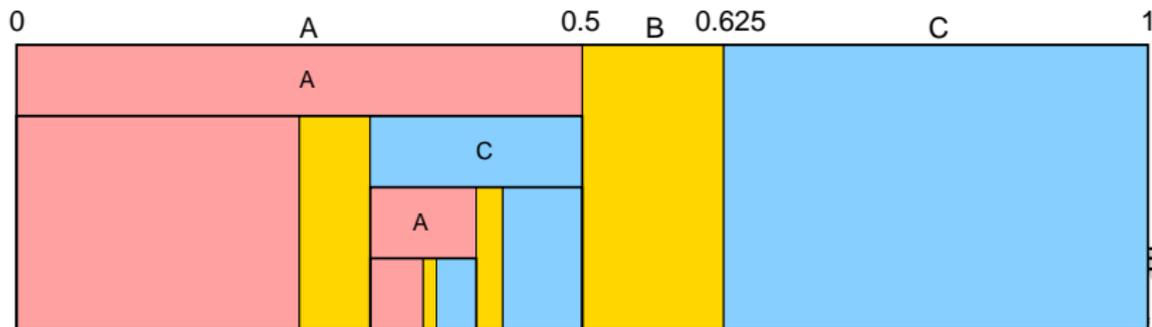
- $[0, 1]$ is split in subintervals for each symbol according to their frequency
- First symbol is *A*. We subdivide its interval
- Then comes *C*
- Then *A* again
- Then *B*
- Each real enclosed into this small interval can encode *ACAB*



Arithmetic compression

Example

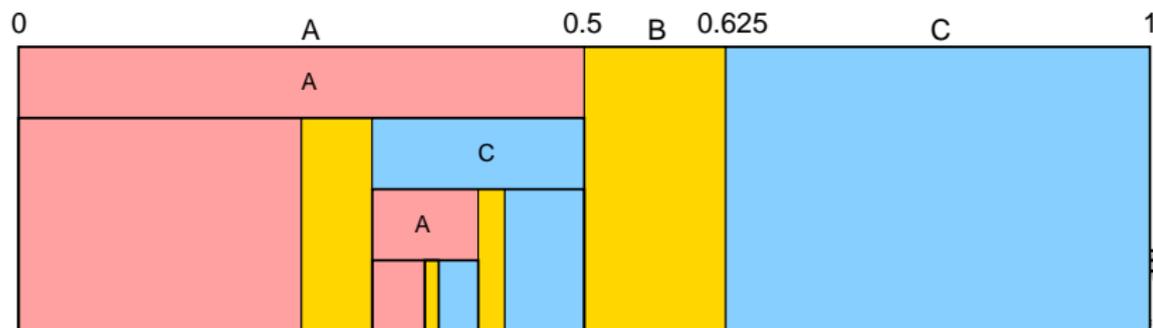
- $[0, 1]$ is split in subintervals for each symbol according to their frequency
- First symbol is *A*. We subdivide its interval
- Then comes *C*
- Then *A* again
- Then *B*
- Each real enclosed into this small interval can encode *ACAB*



Arithmetic compression

Example

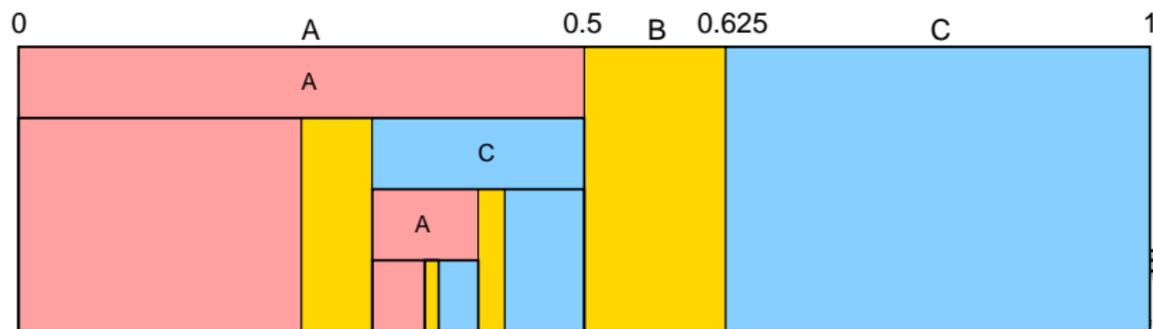
- $[0, 1]$ is split in subintervals for each symbol according to their frequency
- First symbol is A . We subdivide its interval
- Then comes C
- Then A again
- Then B
- Each real enclosed into this small interval can encode $ACAB$



Arithmetic compression

Example

- $[0, 1]$ is split in subintervals for each symbol according to their frequency
- First symbol is A . We subdivide its interval
- Then comes C
- Then A again
- Then B
- Each real enclosed into this small interval can encode $ACAB$



Outline

- 1 Introduction
- 2 Networking**
 - Compression
 - **Analysis of the login phase**
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Trusted data

Embedded trusted data

In order to recognize Skype authority, the binary has 14 moduli.

Moduli

- Two 4096 bits moduli
- Nine 2048 bits moduli
- Three 1536 bits moduli

RSA moduli example

- 0xba7463f3...c4aa7b63
- ...
- 0xc095de9e...73df2ea7

Finding friends

Embedded data

For the very first connection, IP/PORT are stored in the binary

Login servers

```
push    offset aLibConnectionL ; "*Lib/Connection/LoginServers"
push    45h
push    offset a195_215_8_1413 ; "195.215.8.141:33033 212.72.49.14"
mov     ecx, eax
call    sub_7B8440
```

Supernodes

- A list of 200 supernodes is hardcoded in the binary
- It changes in every version and subversion of Skype

Phase 0: Hypothesis

Trusted data

- Each message signed by one of the Skype modulus is trusted
- The client and the Login server have a shared secret
 - A MD5 hash of the user's information

Phase 1: Key generation

Session parameters

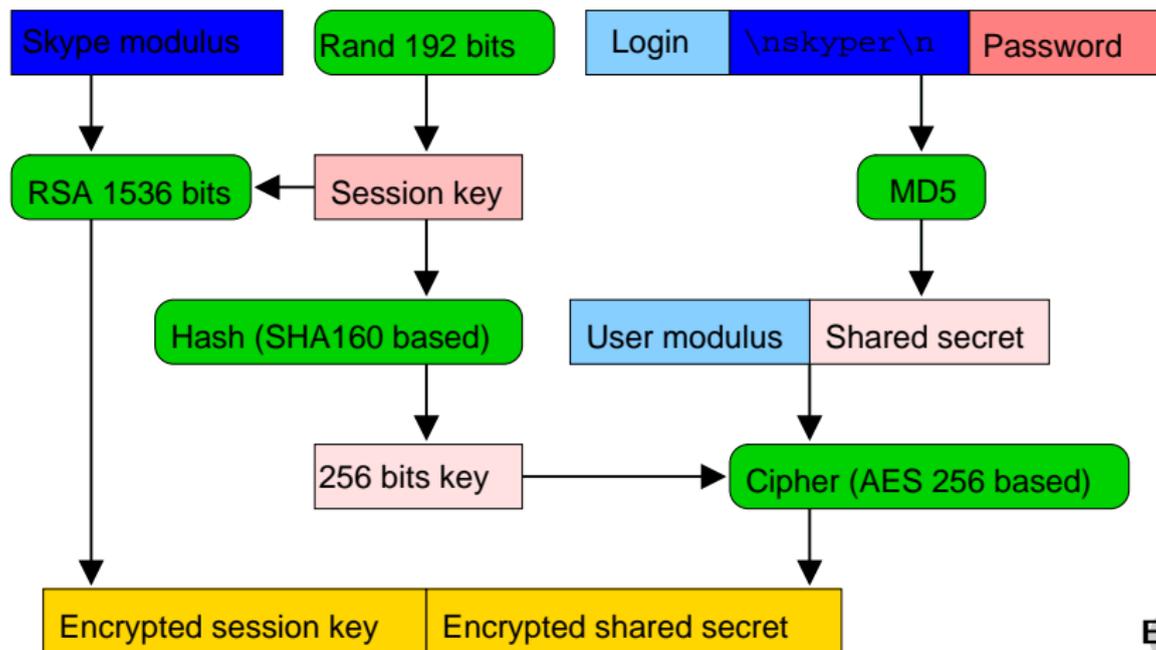
- When a client logs in, Skype will generate two 512 bits length primes
- This will give 1024 bits length RSA private/public keys
- Those keys represent the user for the time of his connection
 - Or longer if the user chooses to save them
- The client generates a symmetric session key K

Phase 2: Authentication

Key exchange

- The client hashes its `login||\nskyper\n||password` with MD5
- The client ciphers its public modulus and the resulting hash with K
- The client encrypts K using RSA with one of the trusted Skype modulus
- He sends the encrypted session key K and the ciphered data to the login server

Phase 2: Authentication



Phase 3: Running

Session behavior

- If the hash of the password matches, the login associated with the public key is dispatched to the supernodes
- This information is signed by the Skype server
- Note that private informations are signed by each user

Search for buddy

- If you search for a login name, a supernode will send back his couple
- You receive the public key of the desired buddy
- The whole packet is signed by a Skype modulus

Example of encrypted stuff

Public blob

```

0 |4bbbbbbb bbbbbbbb bbbbbbbb bbbbbbbb| K.....
10 |bbba4104 03007265 63636137 37000003| ..A...recca77...
20 |00040180 01d3e860 164f8a1b 0a771e5b| .....'.0...w.[
30 |d74e1548 b96fa8bb 712167c9 0273003b| .N.H.o..q!g..s.;
40 |e201d464 d92d2d13 073a6622 5aae2c28| ...d.--.:f"Z.,(
50 |f80640ff 40b9327e 98781fe5 9b6dadfa| ..@.@.2~.x...m..
60 |b7fbcbf7 84a4bf66 051682fc 4dadae53| .....f....M..S
70 |3317c5bf 5be61f2f 7458a133 faa61731| 3...[../tX.3...1
80 |ed910a83 abc70cd1 cf7c2876 e23f60bc| .....|(v.¿.
90 |667d0533 8ce755a8 c66e463b 6d60b13a| f}.3..U..nF;m'.:
a0 |2d0a107c 2900048c 84950926 5fb26626| -..|).....&_.f&
b0 |4ea8968c 6a7a6d2c 97c78ae4 ed967fbc| N...jzm,.....
  
```

Phase 4: Communicating

Inter client session

- Both clients' public keys are exchanged
- Those keys are signed by Skype authority
- Each client sends a 8 bytes challenge to sign
- Clients are then authenticated and can choose a session key

Some strings to guide you

```
db 'session_manager: [%04x] remote party sent wrong identity',0Ah,0
db 'session_manager: [%04x] remote party failed challenge',0Ah,0
db 'session_manager: [%04x] missing challenge response',0Ah,0
db 'session_manager: [%04x] remote UIC has expired',0Ah,0
db 'session_manager: [%04x] no encryption key in reply',0Ah,0
```



Outline

- 1 Introduction
- 2 Networking**
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic**
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Detecting Skype Traffic

Some ideas to detect Skype traffic without deobfuscation

- Most of the traffic is crypted . . . But not all.
- UDP communications imply clear traffic to learn the public IP
- TCP communications use the same RC4 stream twice !

Detecting Skype Traffic

TCP traffic

- TCP stream begin with a 14 byte long payload
- From which we can recover 10 bytes of RC4 stream
- RC4 stream is used twice and we know 10 of the 14 first bytes

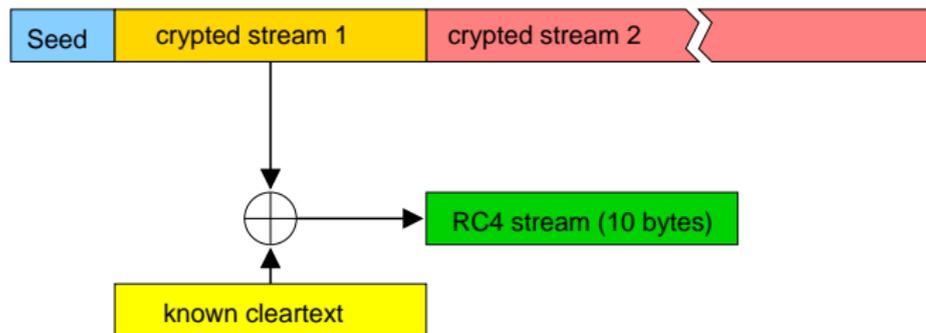


known cleartext

Detecting Skype Traffic

TCP traffic

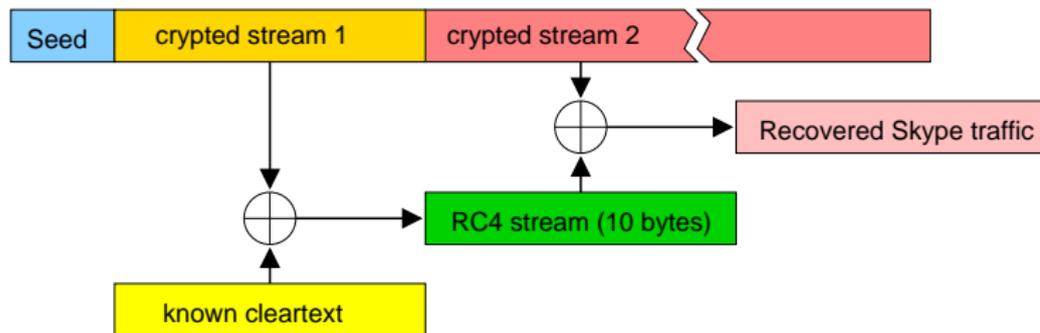
- TCP stream begin with a 14 byte long payload
- From which we can recover 10 bytes of RC4 stream
- RC4 stream is used twice and we know 10 of the 14 first bytes



Detecting Skype Traffic

TCP traffic

- TCP stream begin with a 14 byte long payload
- From which we can recover 10 bytes of RC4 stream
- RC4 stream is used twice and we know 10 of the 14 first bytes

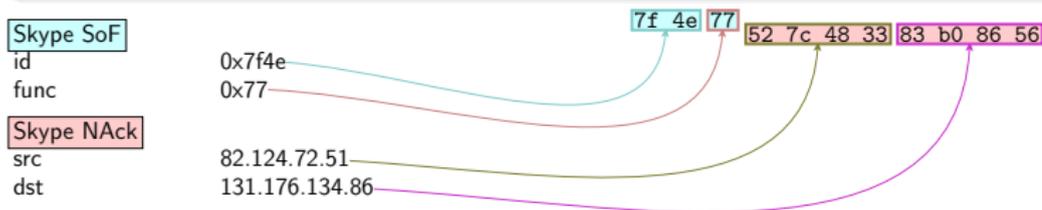


Detecting Skype Traffic

UDP traffic

Skype NACK packet characteristics

- $28+11=39$ byte long packet
- Function & $0x8f = 7$
- Bytes 31-34 are (one of) the public IP of the network



Detecting Skype Traffic

Blocking UDP traffic

On the use of NACK packets...

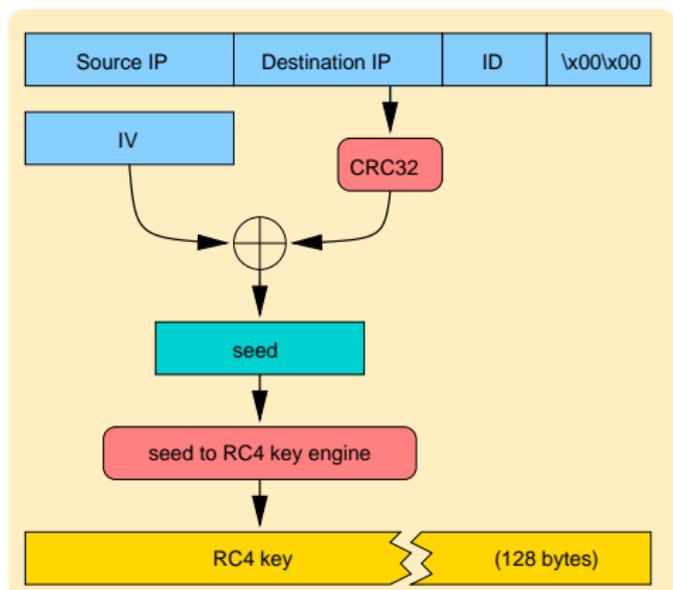
- The very first UDP packet received by a Skype client will be a NACK
- This packet is not crypted
- This packet is used to set up the obfuscation layer
- Skype can't communicate on UDP without receiving this one

How to block Skype UDP traffic with one rule

```
iptables -I FORWARD -p udp -m length --length 39 -m u32 \
  --u32 '27&0x8f=7' --u32 '31=0x527c4833' -j DROP
```

How to generate traffic without the *seed* to RC4 key engine

- Get the RC4 key for a given seed for once
- Always use this key to encrypt
- Calculate the CRC stuff
- Use $IV = seed \oplus crc$



Outline

- 1 Introduction
- 2 Networking**
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - **Nice commands**
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Firewall testing (a.k.a remote scan)

Let's TCP ping Slashdot

```
>>> send(IP(src="1.2.3.4",dst="172.16.72.19")/UDP(sport=1234,dport=1146)
/Skype_SoF(id=RandShort())/Skype_Enc()/Skype_Cmd(cmd=41, is_req=0,
is_b0=1, val=Skype_Encod(encod=0x41)/Skype_Objects_Set(objnb=1)
/Skype_Obj_INET(id=0x11, ip="slashdot.org", port=80)))
```

A TCP connect scan from the inside

```
>>> send(IP(src="1.2.3.4",dst="172.16.72.19")/UDP(sport=1234,dport=1146)
/Skype_SoF(id=RandShort())/Skype_Enc()/Skype_Cmd(cmd=41, is_req=0,
is_b0=1, val=Skype_Encod(encod=0x41)/Skype_Objects_Set(objnb=1)
/Skype_Obj_INET(id=0x11, ip="172.16.72.1", port=(0,1024)))
```

A look for MS SQL from the inside

```
>>> send(IP(src="1.2.3.4",dst="172.16.72.19")/UDP(sport=1234,dport=1146)
/Skype_SoF(id=RandShort())/Skype_Enc()/Skype_Cmd(cmd=41, is_req=0,
is_b0=1, val=Skype_Encod(encod=0x41)/Skype_Objects_Set(objnb=1)
/Skype_Obj_INET(id=0x11, ip="172.16.72.*", port=1433)))
```



Firewall testing (a.k.a remote scan)

Me: *Say hello to slashdot.org:80*

IP 1.2.3.4.1234 > 172.16.72.19.1146: UDP, length: 24

Skype: *Yes, master*

IP 172.16.72.19.1146 > 1.2.3.4.1234: UDP, length: 11

Skype: *Hello! (in UDP)*

IP 172.16.72.19.1146 > 66.35.250.151.80: UDP, length: 20

Skype: *connecting to slashdot in TCP*

IP 172.16.72.19.3776 > 66.35.250.151.80: S 0:0(0)

IP 66.35.250.151.80 > 172.16.72.19.3776: S 0:1(0) ack 0

IP 172.16.72.19.3776 > 66.35.250.151.80: . ack 1

Skype: *Hello! (in TCP). Do you speak Skype ?*

IP 172.16.72.19.3776 > 66.35.250.151.80: P 1:15(14) ack 1

IP 66.35.250.151.80 > 172.16.72.19.3776: . ack 15

Skype: *Mmmh, no. Goodbye.*

IP 172.16.72.19.3776 > 66.35.250.151.80: F 15:15(0) ack 1

IP 66.35.250.151.80 > 172.16.72.19.3776: F 1:1(0) ack 16

Skype Network

Supernodes

- Each skype client can relay communications to help unfortunates behind a firewall
- When a skype client has a good score (bandwidth+no firewall+good cpu) he can be promoted to supernode

Slots and blocks

- Supernodes are grouped by slots
- You usually find 9 or 10 supernodes by slot
- You have 8 slots per block

Who are the supernodes ?

Just ask

- Each supernode knows almost all other supernodes
- This command actually ask for at most 100 supernodes from slot 201

```
>>> sr1(IP(dst="67.172.146.158")/UDP(sport=31337,dport=4344)/Skype_SoF(
    id=RandShort())/Skype_Enc()/Skype_Cmd(cmd=6, reqid=RandShort(),
    val=Skype_Encod(encod=0x41)/Skype_Objects_Set(objnb=2)
    /Skype_Obj_Num(id=0,val=201)/Skype_Obj_Num(id=5,val=100))
```

- Nowadays there are ~ 2050 slots
- That means $\sim 20k$ supernodes in the world

More commands

Related to supernodes

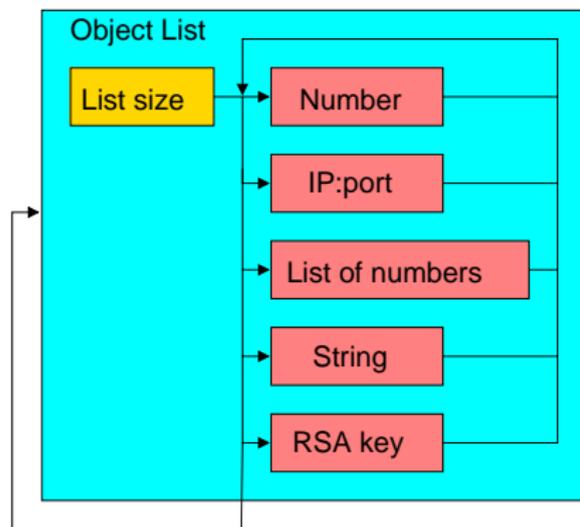
- Promote any client to a supernode
- Ask for supernode clients information
 - Bandwidth
 - Memory
 - OS version
 - Skype version
- Ban any supernode for one hour

Outline

- 1 Introduction
- 2 Networking**
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - **Remote exploit**
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Object lists

- An object can be a number, a string, an IP:port, or even another object list
- Each object has an ID
- Skype knows which object corresponds to which command's parameter from its ID



Space allocation

Algorithm

```
lea    ecx, [esp+arg_4]
push   ecx
call   get_uint
add    esp, 0Ch
test   al, al
jz     parse_end
mov    edx, [esp+arg_4]
lea    eax, ds:0[edx*4]
push   eax
mov    [esi+10h], eax
call   LocalAlloc
mov    ecx, [esp+arg_4]
mov    [esi+0Ch], eax
```

- 1 Read an unsigned int *NUM* from the packet
- 2 This integer is the number of unsigned int to read next
- 3 *malloc* $4*NUM$ for storing those data

Data reading

Algorithm

```
read_int_loop :  
push    ebx  
push    edi  
push    ebp  
call    get_uint  
add     esp, 0Ch  
test    al, al  
jz      parse_end  
mov     eax, [esp+arg-4]  
inc     esi  
add     ebp, 4  
cmp     esi, eax  
jb      read_int_loop
```

- 1 For each *NUM* we read an unsigned int
- 2 And we store it in the array freshly allocated

Heap overflow

How to exploit that?

- If $NUM = 0x80000001$
 - The multiplication by 4 will overflow :
 - $0x80000001 * 4 = 0x00000004$
 - So Skype will allocate $0x00000004$ bytes
 - But it will read NUM integers
- ⇒ Skype will overflow the heap

Exploiting

Reliability

- In theory, exploiting a heap on Windows XP SP2 is not very stable
- But Skype has some Oriented Object parts
- It has some structures with functions pointers in the heap
- If the allocation of the heap is close from this structure, the overflow can smash function pointers
- And those functions are often called

⇒ Even on XP SP2, the exploit is possible 😊

Remote code execution

Loving OOP

Here is the code responsible for the function pointer call

```
push    esi
push    edi
lea     ecx , [ebx+eax]
call   ebp
mov     eax , [ecx]
jmp     dword ptr [eax+8]
```

Skype patch

Code

```
cmp     edi , 3FFFFFFh
jbe     short loc_72F52B
push   offset aAlistSetsizeA1 ; "alist::SetSize(): alloc size ove
```

About the patch

- The same piece of code is present about 60 times
- Each time a comparison with $0x3FFFFFFF$ is done
- Sometimes, the register is not multiplied by 4, but by 5 or more

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 **Skype API**
 - **Filtering**
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Chat filtering

Chinese censorship

- TOM-Skype and eBay.cn-Skype censor **incoming** text messages on behalf of the Chinese government
- Both versions are shipped with a `ContentFilter.exe` binary
- It is a plugin that is verified and loaded automatically by Skype
- Words are matched against an encrypted list of simplified chinese expressions

Undocumented API

- A filtering API is activated in those Skype versions
 - `FILTERING ON` will start a message redirection mechanism
 - `FILTER n OK` or `FILTER n BLOCK` will allow or block a message submitted to the filtering plugin

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 **Skype API**
 - Filtering
 - **AP2AP**
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Hiding behind Skype

AP2AP

An interesting feature of the API is the **Application to Application** protocol, which allows two applications to communicate through Skype

- They benefit from Skype NAT and Proxy bypassing abilities
- The data is encrypted by Skype itself
- The remote endpoint is only identified by a login and not an IP address

Uses

- Exfiltration
- Discrete remote control of the machine
- File transfers
- Network connections tunneling

Encrypted tunnels

Sample applications

- AP2AP remote cmd.exe
- AP2AP socks v4, v4a and v5 proxy
- AP2AP key logging

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun**
 - Randomness**
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Random number generator (1/2)

Code

```
unsigned __int64 Skype_8ByteRandom(void)
{
    BYTE pbBuffer[1124];
    SHA1_CTX SHA1Context;

    memcpy(&pbBuffer[16], Skype_RandomSeed, SHA1_DIGLEN);
    GlobalMemoryStatus((LPMEMORYSTATUS)&pbBuffer[36]);
    UuidCreate((UUID *)&pbBuffer[64]);
    GetCursorPos((LPPPOINT)&pbBuffer[76]);
    *(DWORD *)&pbBuffer[80] = GetTickCount();
    *(DWORD *)&pbBuffer[84] = GetMessageTime();
    *(DWORD *)&pbBuffer[88] = GetCurrentThreadId();
    *(DWORD *)&pbBuffer[92] = GetCurrentProcessId();
    QueryPerformanceCounter((LARGE_INTEGER *)&pbBuffer[96]);
    SHA1_Init(&SHA1Context);
    SHA1_Update(&SHA1Context, &pbBuffer[0], 1124);
    SHA1_Update(&SHA1Context, "additional salt...", 19);
    SHA1_Final(Skype_RandomSeed, &SHA1Context);

    return Skype_8ByteSHA1(&pbBuffer[0], 1124);
}
```

Random number generator (2/2)

Code

```
static BYTE Skype_RandomSeed[SHA1_DIGLEN];

unsigned __int64 Skype_8ByteSHA1(BYTE *pbData, DWORD dwLength)
{
    SHA1_CTX SHA1Context;
    BYTE pbHash[SHA1_DIGLEN];

    SHA1_Init(&SHA1Context);
    SHA1_Update(&SHA1Context, &pbData[0], dwLength);
    SHA1_Final(pbHash, &SHA1Context);

    return *(unsigned __int64 *)&pbHash[0];
}
```

My 2 cents

- The random number generator implementation is quite strong, thus giving a good base to all the overlying cryptography
- Surprisingly, some parts of the structures used are overwritten

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun**
 - Randomness
 - Easter eggs**
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Eggy

Easter egg in the chat module

- Removed in version 2.0.0.103 and later
 - **Skype people do read our slides !**
- Triggered by a command in a chat window
 - /eggy <secret>
- Decrypts and displays one of two texts given <secret>
 - 1st if (length == 6 && crc32 == 0xb836ac79)
 - 2nd if (length == 14 && crc32 == 0x0407aac1)

Decryption algorithm

```
for (i = 0, x = 0; i < (strlen(szInput) >> 1); i++) {  
    szOutput[i] = ((szInput[(i << 1) + 1] << 4) |  
        (szInput[i << 1] & 0xbf)) ^ x ^ szKey[i % strlen(szKey)];  
    x ^= szOutput[i];  
}
```

Breaking the egg

Dictionary and bruteforce attack

- Based on length and crc32 values
- Decrypted text will allow to settle in the event of collisions
- 1st secret found : **prayer**

Cryptanalysis

- Model the cipher like a usual one time pad with a known key length
 - $c'_i = c_i \oplus c_{i-1}$ with $c'_1 = c_1$
 - $k'_i = k_i \oplus k_{i-1}$ with $k'_1 = k_1$
- Carry on with a usual statistical cryptanalysis attack
- 2nd secret found : **indrek@mare.ee**

Example

Crypted text

```
MCBEMCK@LF@ADENA@FBAHFND@FBANCKEDCJDDCDEKAFANFEAGFL  
@NB@DHCJEBBJELBNDEDOALGMAAFCDFFA@NGIELCLDKGFBFFBCND  
HCO@GBD@EFMAFCLAIFFFAMGCCLFCAABLCNCKAOGA@CFB@DCNFA@D  
DM@CGE@BCAEKBBAIBGAMCF@ACLD CAGEGCHDOGEEBGKAAFC@FCI@
```

Key

```
indrek@mare.ee
```

Decrypted text

*The programmer behind the internal workings of Skype chat,
cheers! Indrek Mandre (1979 - still alive?)*

Example

Crypted text

```
MCBEMCK@LF@ADENA@FBAHFND@FBANCKEDCJDDCDEKAFANFEAGFL  
@NB@DHCJEJBBJELBNDEDOALGMAAFCDFFA@NGIELCLDKGFBFFBCND  
HCO@GBD@EFMAFCLAIFFFAMGCCLFCAABLNCNKAOGA@CFB@DCNFA@D  
DM@CGE@BCAEKBBAIBGAMCF@ACLD CAGEGCHDOGEEBGKAAFC@FCI@
```

Key

```
indrek@mare.ee
```

Decrypted text

*The programmer behind the internal workings of Skype chat,
cheers! Indrek Mandre (1979 - still alive?)*

Example

Crypted text

```
MCBEMCK@LF@ADENA@FBAHFND@FBANCKEDCJDDCDEKAFANFEAGFL  
@NB@DHCJEJBBJELBNDEDOALGMAAFCDFFA@NGIELCLDKGFBFFBCND  
HCO@GBD@EFMAFCLAIFFFAMGCCLFCAABLCNCKAOGA@CFB@DCNFA@D  
DM@CGE@BCAEKBBAIBGAMCF@ACLD CAGEGCHDOGEEBGKAAFC@FCI@
```

Key

```
indrek@mare.ee
```

Decrypted text

*The programmer behind the internal workings of Skype chat,
cheers! Indrek Mandre (1979 - still alive?)*

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun**
 - Randomness
 - Easter eggs
 - Debug logs**
 - Plugins
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Logs

Debug logs

- Skype can generate debug logs if some registry keys are set to the correct values in `HKCU\Software\Skype\Phone\UI\General`
 - Logging for encrypted log files
 - Logging2 for clear text log files
- Only the MD5 hashes of the correct values appear in the Windows binary

Enabling logs

- Patch the binary
 - One needs to get rid of all the integrity checks first
- Recover the correct values, which are out of bruteforcing range

Log encryption

Cipher

- Skype generates a 128 bit RC4 key to encrypt logs on the fly
- It is formatted, then encrypted using a 1024 bit RSA public key ($e = 3$), and stored at the beginning of the log file

Encrypted data format



RC4 key

Key format



Recovering the key

- The clear text log file format is known
- The log file name already contains the year, month and day
- The only things remaining are
 - The seconds (0 to 59)
 - The value of `GetTickCount()` (usually $< 2^{24}$)
- If Skype is automatically launched at Windows startup, recovery is instantaneous

RC4 key

Key format



Recovering the key

- The clear text log file format is known
- The log file name already contains the year, month and day
- The only things remaining are
 - The seconds (0 to 59)
 - The value of `GetTickCount()` (usually $< 2^{24}$)
- If Skype is automatically launched at Windows startup, recovery is instantaneous

"Logging"

Hint

```
http://download.skype.com/logging-on-off.zip
```

Traces

Trace file

- Skype voice engine can generate encrypted trace files if Logging and Logging2 are set
- Encryption is much simpler, a basic XOR with a 31 byte key

Decryption algorithm

```
for (i = 0, j = 0; i < strlen(pBuffer); i++, j++, k = (k + 1) % 31) {  
    if (pBuffer[i] == 1) {  
        pBuffer[j] = (127 - pBuffer[i + 1]) ^ pXORTable[k];  
        i++;  
    } else if (pBuffer[i] == 2) {  
        pBuffer[j] = pBuffer[i + 1];  
        i++;  
    } else  
        pBuffer[j] = pBuffer[i] ^ pXORTable[k];  
}  
pBuffer[j] = '\0';
```

Some things you can find in logs

Stack dumps

```
11:35:40 Mutex::Acquire: possible deadlock. Stack dump:  
11:35:40 0012fb8c: 00aebaa4 02057030 001e1d63 0012fc54 0074861b 0012fb8c 0012fd  
11:35:40 0012fbac: 02057030 001e1d63 0012fc00 009eb048 ffffffff 0072fa98 000000  
11:35:40 0012fbcc: 0012fbe4 03ce2540 007274ad 000000f5 0012fbe4 0043d7e8 000000  
...
```

Assert failures

```
10:21:38 Call #2: StartPlayout (1 1)  
10:21:38 Call #2: setting audio bandwidth to 2625 pkt 60ms  
10:21:38 ASSERTFAILURE(Channel && VE->EngineInited && Recording) in D:\Src\GI\S
```

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun**
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins**
 - Chinese Blacklist
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Plugin "signing"

Skype plugins ACL

- Skype implements an ACL-like system to allow or disallow plugins to attach themselves to a running instance
- A plugin "signature" is added to the configuration file based on the user reply to a warning dialog

Example entry

```
<AccessControlList>
  <Client1>
    <Key1>623df12b13d8dea5e32ea1f8467f3d2f040f662d0e604032a08cca9cd243
    <Key2>31823a73a63c38a2e7ead0a2408a7f2a</Key2>
    <Key3>263594</Key3>
    <Path>D:\Skype\Plugins\plugin_master.exe</Path>
  </Client1>
</AccessControlList>
```

Warning dialog

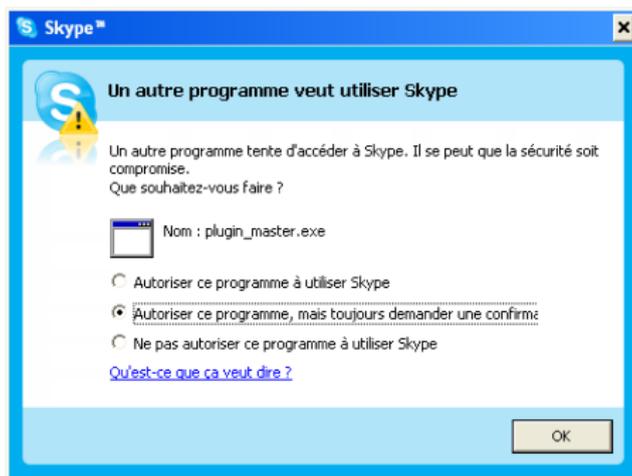


Figure: "Permit", "Ask" or "Ban" a plugin

"Signing" plugins

Hashes to hashes

- The "signature" mechanism is just about MD5 hashes of the full path, the binary, and the ACL specified by the user
- Nothing much can stop us from writing our own and add it to the configuration file !

Pseudo-code (*'.'* is concatenation)

```
szSalt = "Element'ry!penguInS;-)SingingHareKrishna_"
szKey1 = Str(Md5(Str(Md5(Upr(szPath) . szSalt))))
        . Str(Md5(Str(Md5(pBinary)) . szSalt))
szKey2 = Str(Md5("Per" . Upr(szPath) . "mit"))
szKey3 = "0" // Last Hwnd of the plugin
```



Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun**
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - **Chinese Blacklist**
- 5 Credentials
 - More networking
 - Credentials
- 6 Conclusion

Encrypted blacklist

Keyfile

- On startup, TOM-Skype ContentFilter.exe fetches an encrypted keywords list file at `http://skypetools.tom.com/agent/keyfile`
- Each line is an AES encrypted regular expression
- A 32 character key is hardcoded in unicode in the binary
 - Only the 1st 32 bytes are used

Extract

```
[\.*\ \,*\;]*t[\.*\ \,*\;]*e[\.*\ \,*\;]*s[\.*\ \,*\;]*t[\.*\ \,*\;]*i[\.*\ \,*\;]*
[\.*\ \,*\;]*f[\.*\ \,*\;]*u[\.*\ \,*\;]*c[\.*\ \,*\;]*k[\.*\ \,*\;]*
[\.*\ \,*\;]*6[\.*\ \,*\;]*2[\.*\ \,*\;]*7[\.*\ \,*\;]*9[\.*\ \,*\;]*7[\.*\ \,*\;]*
```

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 **Credentials**
 - **More networking**
 - Credentials
- 6 Conclusion

Session half key exchange

How does it work ?

- Each peer generates a 128 bit random nonce, extends it to 1024 bits by repeating it
- The extend nonce is encrypted using the RSA public key of the other peer
- Each peer decrypts the received data and computes 128 bits of the 256 bit AES session key

Some maths

- $C = 1 + 2^{128} + 2^{256} + 2^{384} + 2^{512} + 2^{640} + 2^{768} + 2^{896}$
- $m = x * C$ and $m' = m^e \bmod n$, so $m' = x^e * C^e \bmod n$
- $m'' = x^e \bmod n$ with $m'' = m' * C^{-e} \bmod n$

The "weakness"

- Best known attack is in 2^{64}
 - <http://citeseer.ist.psu.edu/boneh00why.html>
- NSA can probably do better ☺

Outline

- 1 Introduction
- 2 Networking
 - Compression
 - Analysis of the login phase
 - Playing with Skype Traffic
 - Nice commands
 - Remote exploit
- 3 Skype API
 - Filtering
 - AP2AP
- 4 Skype cryptography fun
 - Randomness
 - Easter eggs
 - Debug logs
 - Plugins
 - Chinese Blacklist
- 5 **Credentials**
 - More networking
 - **Credentials**
- 6 Conclusion

Saved credentials

What does Skype save ?

- If told to, Skype will save in the `config.xml` file
 - The login MD5 hash (`username\nskype\r\npassword`)
 - The generated RSA **private** key
 - The Skype encrypted corresponding RSA public key
- Everything is heavily encrypted, but in a symmetric way :)
- The following algorithms are used
 - `CryptProtectData()`, `CryptUnprotectData()`
 - SHA-1
 - AES-256
 - "FastTrack cipher"
 - 1024+ bit RSA

Credentials structure

Version 1

- 16 bytes for login MD5 hash
- 128 bytes for user RSA private key (D) (1024 bits)
- 4 bytes for Skype RSA key ID
- 192+ bytes for RSA block encrypted with Skype RSA key
 - Padding
 - Skype encoded data
 - User name
 - 1 dword
 - User RSA public key (N) (1024 bits)
 - 1 dword
 - SHA-1 hash of Skype encoded data
 - 1 byte = 0xbc
- 2 bytes for CRC32 (reduced to 16 bits)

Decrypting the credentials 1/2

Recovering the AES 256 bit key

- Unprotect the token from
HKCU\Software\Skype\ProtectedStorage
- Use incremental counter mode SHA-1 to create a 32 byte key from the token

Decrypting the 1st layer

- Use incremental counter mode AES to decrypt the credentials
- **Login MD5 hash** is now decrypted

Decrypting the 2nd layer

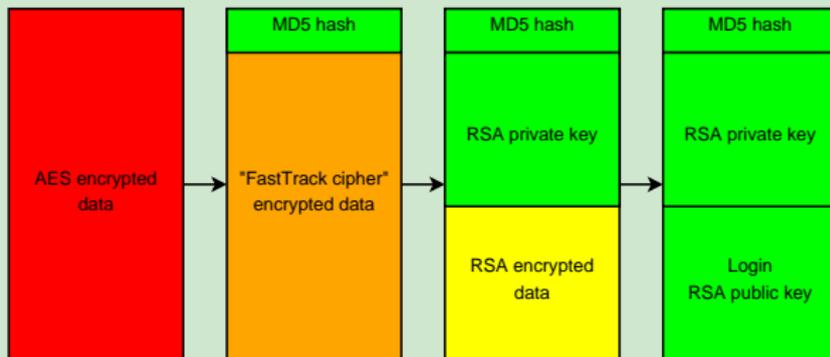
- Use the login MD5 hash as key for the "FastTrack cipher"
- Decrypt the rest of credentials data
- **RSA private key** is now decrypted

Decrypting the credentials 2/2

Decrypting the 3rd layer

- Use the correct Skype public key to decrypt the remaining RSA block
- **RSA public key** is now decrypted

Graphical summary



Saved credentials usage

Login MD5 hash

- Skype password recovery
 - Dictionary attack
 - Bruteforce attack

RSA private key

- Sniffed session half key recovery
 - Decrypt the 128 bit random nonce exchanged
 - Compute half of the AES-256 session key
- Complete sniffed session key recovery
 - If both RSA private keys are recovered
- \implies Sniffed conversation decryption

Conclusion

Auditing a software

- Auditing a binary in its complete form is much more accurate than auditing a portion of the sources
- Skype, Inc. clearly doesn't tell you everything

Skype v2.5

- The developers have silently modified the behaviour of Skype carefully following the BlackHat talk points
 - Most of the sensitive commands are now TCP only
 - Some *very* sensitive commands are only accepted when coming from the currently-connected-to supernode only
 - Some features have simply been trashed

Acknowledgements

Shouts to

Phil, Imad, Dave, Halvar, Gera, Team Rstack, Microsoft

MD5ed props to (from a former life)

17f063b9c9f793dc841c7fee0f76eede

Questions ?

