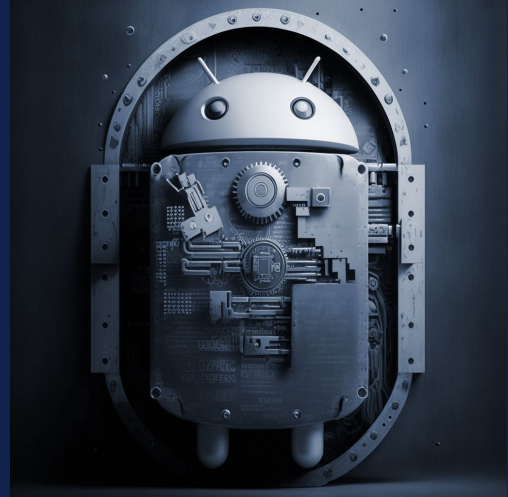


Dissecting the Modern Android Data Encryption Scheme

Maxime Rossi Bellom
Damiano Melotti



Who we are



- [@DamianoMelotti](#)
- Security researcher @ Quarkslab
- Interested in low-level mobile security and fuzzing

- [@max_r_b](#)
- Security researcher and R&D leader @ Quarkslab
- Working on mobile and embedded software security



The trigger

> Hey! My device fell into water and the main SoC is dead. However the Titan M^{1,2} chip seems to be alive and well, do you think you would be able to help me recover my data on the phone?

[1]: [2021: A Titan M Odyssey](#) (Maxime Rossi Bellom, Philippe Teuwen, Damiano Melotti)

[2]: [Attack on Titan M, Reloaded: Vulnerability Research on a Modern Security Chip](#) (Damiano Melotti, Maxime Rossi Bellom)

The trigger

- Our answer: no, the main SoC is still essential for disk encryption/decryption
 - ... but up to what extent?
- Objective of this research: find out exactly
- Offensive approach:
 - What would a forensic analyst do?
 - Assuming infinite vulnerabilities, what can you do to get the secrets out?
 - Do you still need to bruteforce credentials?

Data Encryption at Rest 101

- Idea: no sensitive plaintext files in storage
 - Attackers must not find files in clear on disk
- Threat model: full physical access to powered-off device
- Data is automatically encrypted when written and automatically decrypted when read
- How?
 - Android: Full-Disk Encryption and File-Based Encryption (required from Android 10)
 - Underneath: dm-crypt for FDE, fscrypt for FBE

File-Based Encryption at Rest 101

- Relies on fscrypt, implemented in the Linux kernel
 - It supports Ext4, F2FS, and UBIFS
- Operates at the filesystem level
 - Allows files encrypted with different keys or unencrypted in a file system
- A master key is provided for directory tree
- Then derived per file keys (for regular file, directory, and symbolic link)
- Metadata are not encrypted by fscrypt

Android File-Based Encryption

- Each file has its own key
- Direct Boot and multi-user support
- Two encryption levels:
 - Credential Encrypted (CE), available only after authentication
 - Device Encrypted (DE), available also during boot
- In short, 2 “main” keys
 - DE key, for data decrypted at boot
 - CE key, available after authentication, protecting user data
- DE key is automatically decrypted using HW-backed keys

Android File-Based Encryption

```
a22x:/ # ls /data/data | head -n 20
+BWqxAAAAAQGI050Z57bZxMl6oTCNKCs
+geIWCAAAAgIsCJB+mpPqIQYOH0FrnojC1KJ8e0lvGZWJYXWFTc0WD
+lE39AAAAAwnVineHXItKqcE1Glo9+EinbUaa3wvp,fXEjaJ3r4BiC
+ostbBAAAAwZ2g592ei8GG2DMfc4y6H94jxkEfoRzUDdlQZn0YZKqA5VCUgi89sf2JN8yCBFraC
,0wF+BAAAAgQtSCvHd5rRmVC2lvxEht1Eo50M9kma1oe+vkWT176K8dxofVp5RcmnLv0xC7WMLJ
,WGCADAAAAg66o3+mZ7f009fBNp8zQSdWBqRJiWPUZHafQiTu7khxB
,eGnGDAAAAw+ZAptd14KRyth5ncJmJkYAZBTBW7DoUNpMamRGj05MA
,vu29CAAAAgKApb3amvoChi0pYILwr5xxvUtf0TpZ2h6Cr0wG4CR5D
0MQ5GCAAAQAAtXNLN524L3GwuGAek+rVelNjgE5mwqljis0kydZa2FKJlVD6ezJJGkcjcRTTD7B
0QcmpBAAAAwMxohHqP+AT3ktRtIAJk9,bu2g04xwjLzf,vgN71QQ2B
0WZ9mBAAAAA5KuJUTY261eUrgnVJAloJnZDm3b7JybYYPw8xhLrf0D
11ppkCAAAAI4M8ENP6k0t3rA9kIN9bepZ0FdpiYDp6bQ5Cj2IU2NB
13nD7DAAAAw,NoobtA0XesI0kFqFC4MIwoKYAfdBQrGTm8rfbbdsNC
1K0haDAAAAAIbCCdZY7LE6W4+DAXnAiv,PEhu4CiUL6ofn9ZuL6buC
1V6luAAAAAwmrNINeM84y83,LYzeNEkdCcACNxLAhJ+sI7VzbV51yuebJw9M0PTrjRi9cv391kD
1fMeDAAAAAgPnPYDGu0WWC8uWiojo+nhhnqWk+x0ZerDfrWY3ZVIbB
1hGHtCAAAAwfxP3KkXNUYQ0SwoBzqZcsbu2g04xwjLzf,vgN71QQ2B
2d3fjDAAAAGf7ilkMftxkJiD0JbRx9,dtg04Rk9L1P5CCl7JuqqewNpAl07TQq0KJh2yXzcj9,8JfbjIHwnpEmj7FQ0ixXBP
2v,UXBAAAAQloHyVK08uify8onfmrQXJawy1Dqwwq6kc2g2BKr0H48D
2wjJhBAAAAGQ0MvGF3NqTpssmh6XgAWw1jHW986vSwQMXIPQQz6ghD
```

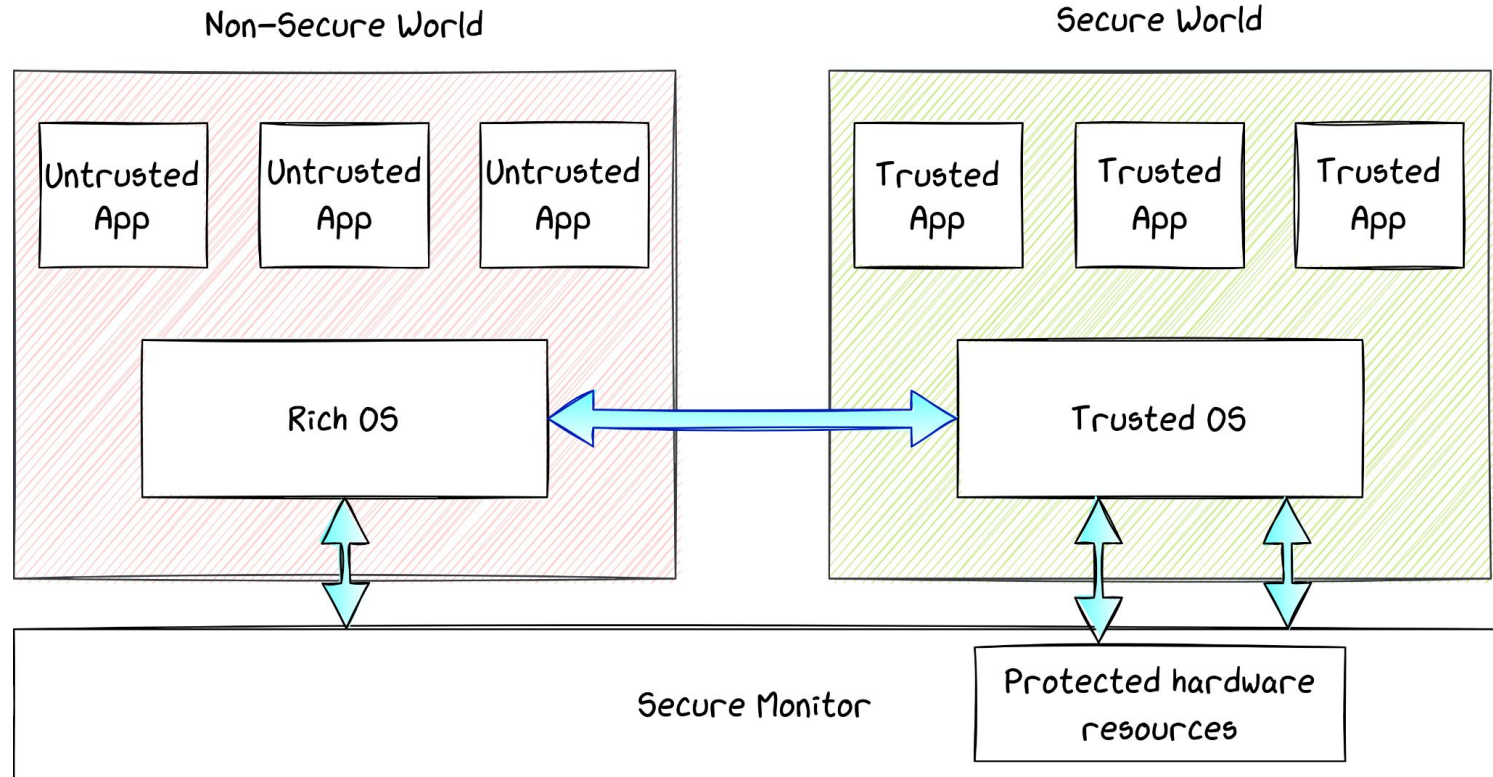

FBE key derivation

- We focus only on the CE key
- Complex derivation steps
 - Start from DE files owned by privileged users

```
a22x:/ # ls -l /data/system_de/0/spblob/
total 32
-rw----- 1 system system 58 2022-06-29 21:59 000000000000000000.handle
-rw----- 1 system system 72 2022-06-29 21:59 921e9ab09afd8d9d.metrics
-rw----- 1 system system 93 2022-06-29 21:59 921e9ab09afd8d9d.pwd
-rw----- 1 system system 16384 2022-06-29 21:59 921e9ab09afd8d9d.secdisk
-rw----- 1 system system 186 2022-06-29 21:59 921e9ab09afd8d9d.spblob
```

- User credentials are used in the process
 - No matter how many bugs an attacker has, bruteforcing remains necessary!

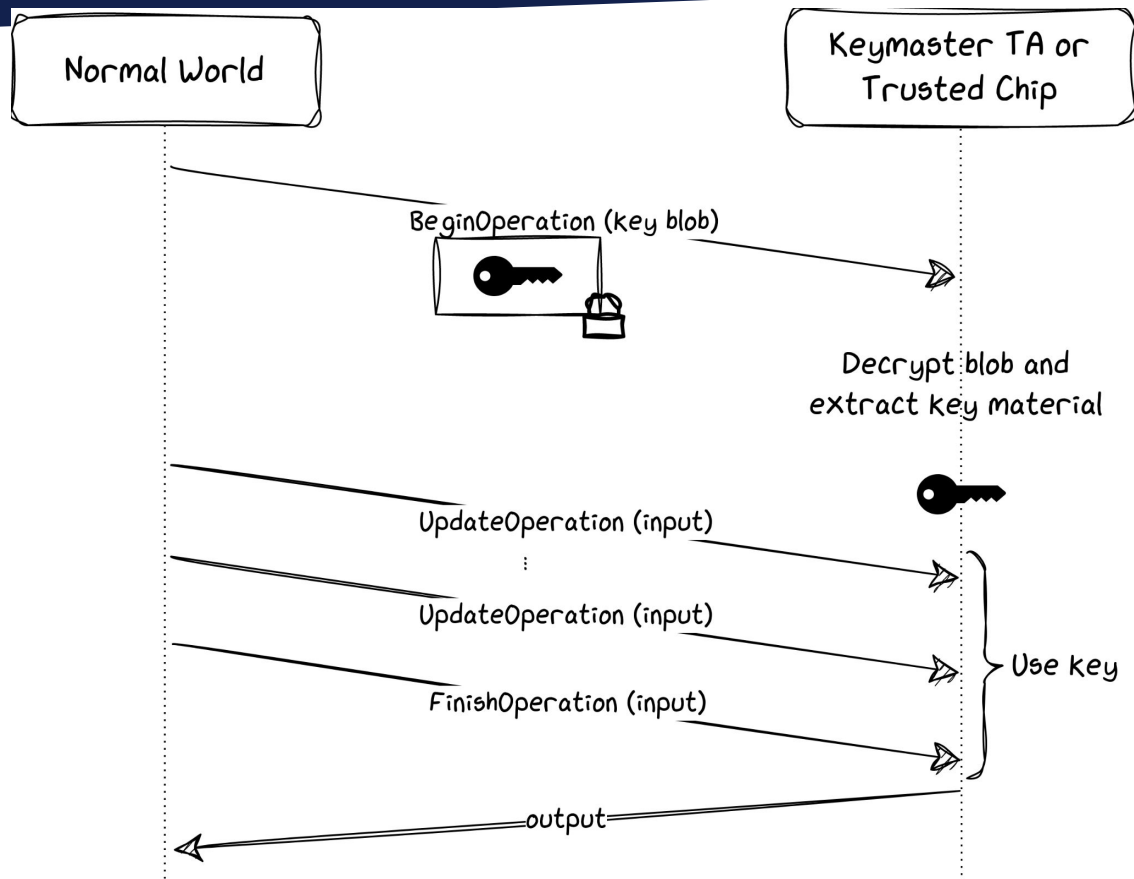
ARM TrustZone

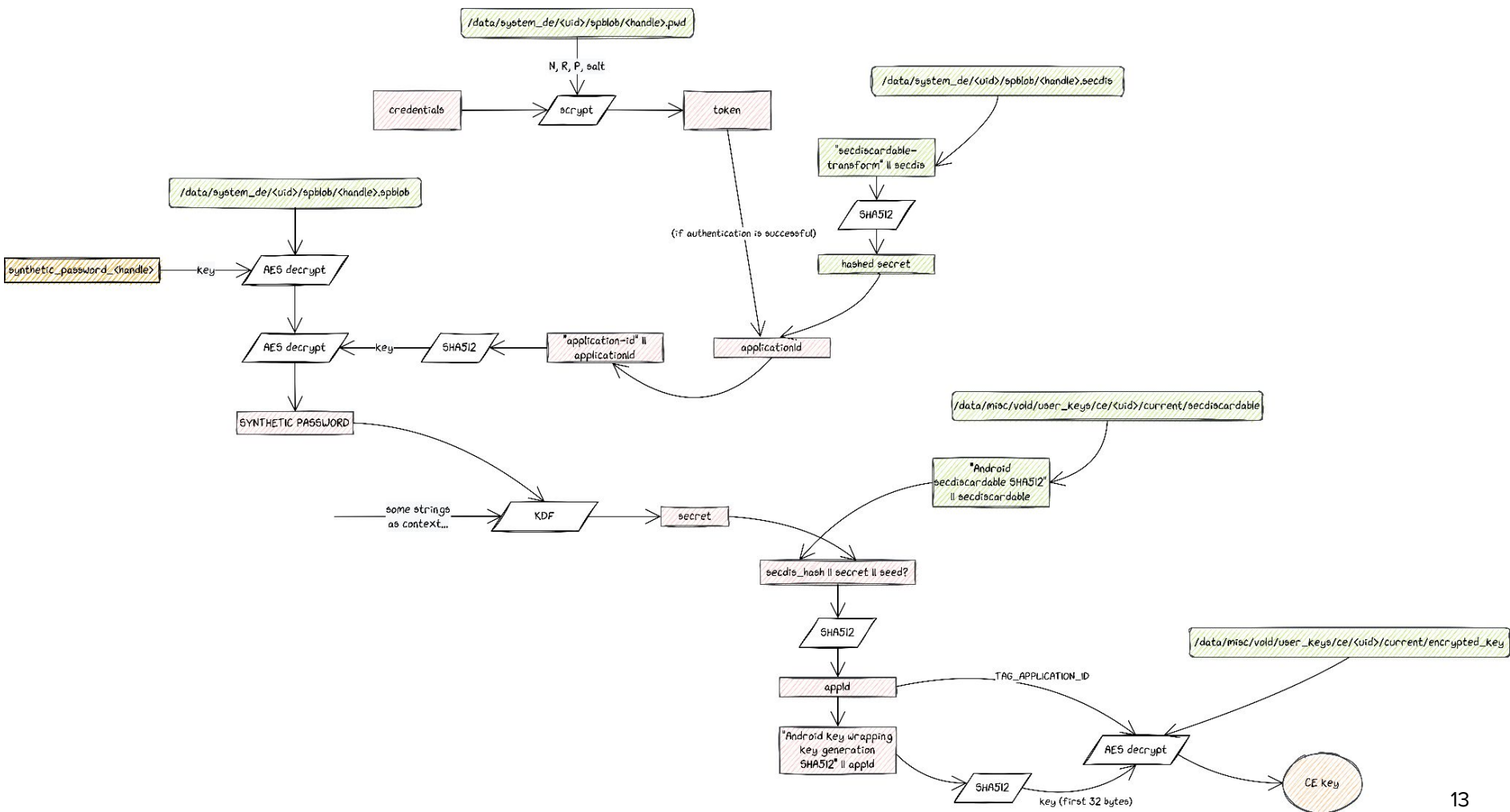


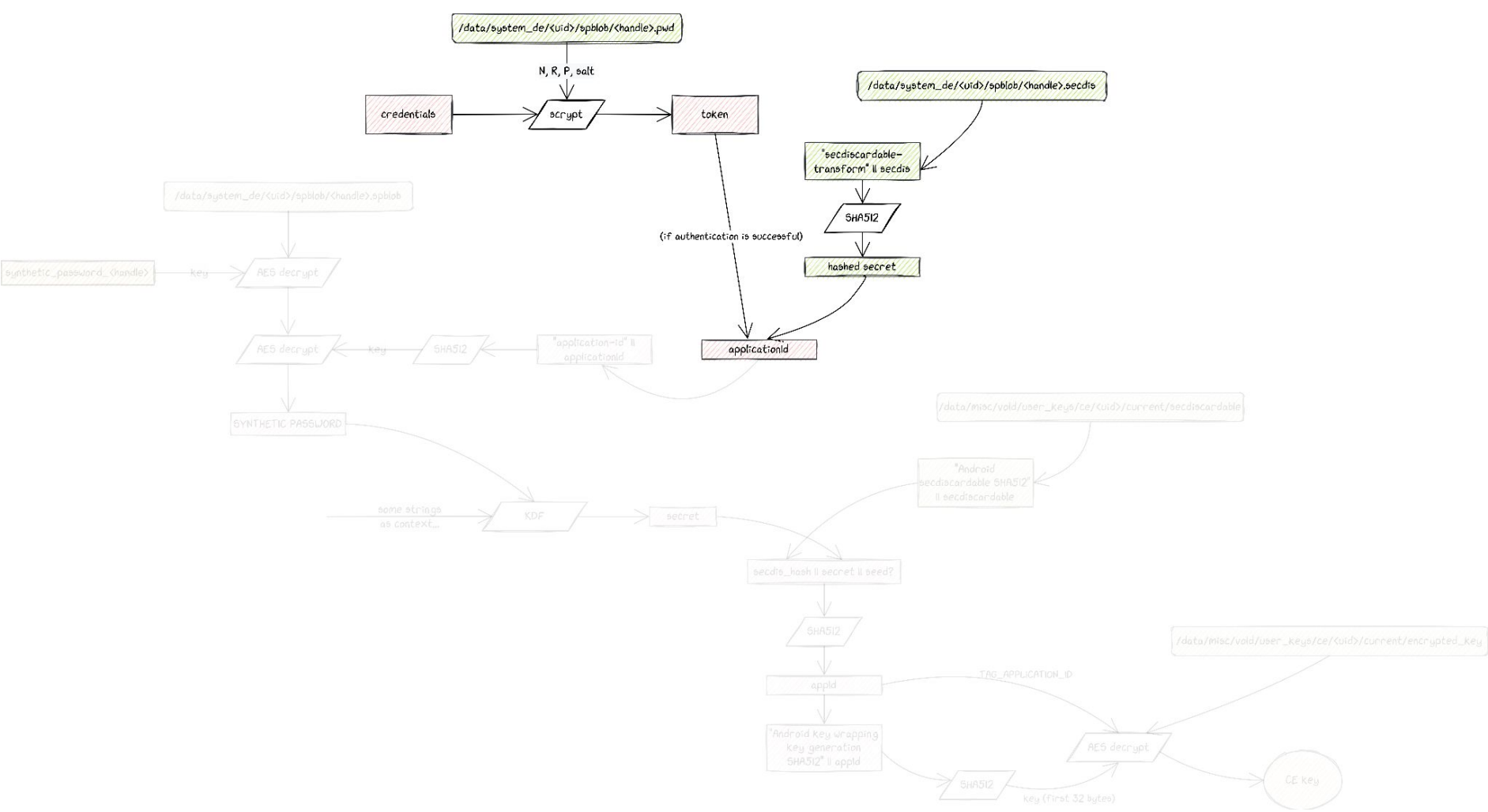
Android Keystore system

- Key storage and crypto services
- Keys are stored as *key blobs*
- Three protection levels:
 - Software only
 - TEE (default)
 - Hardware-backed (StrongBox)
- Raw key should never leave protected environment

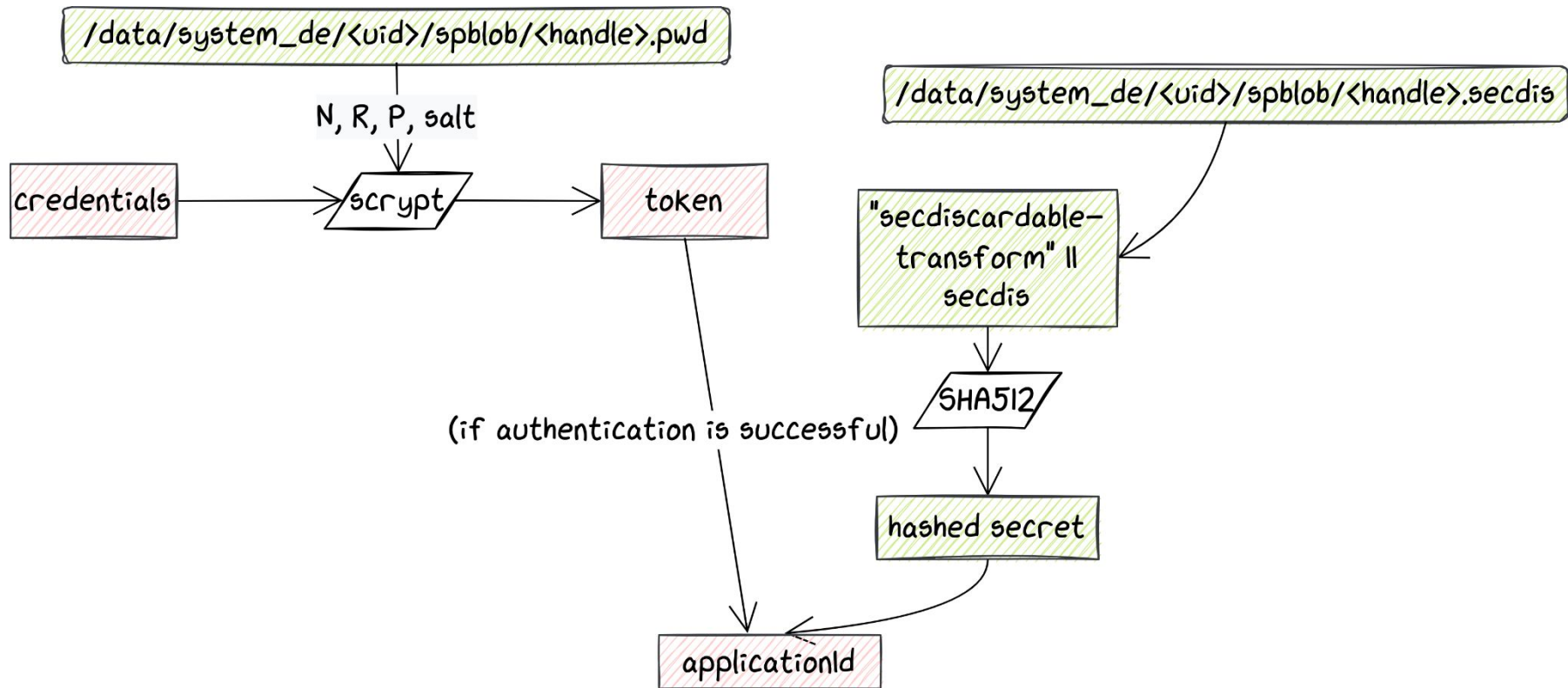
Android Keystore system







Credentials, script, secdis

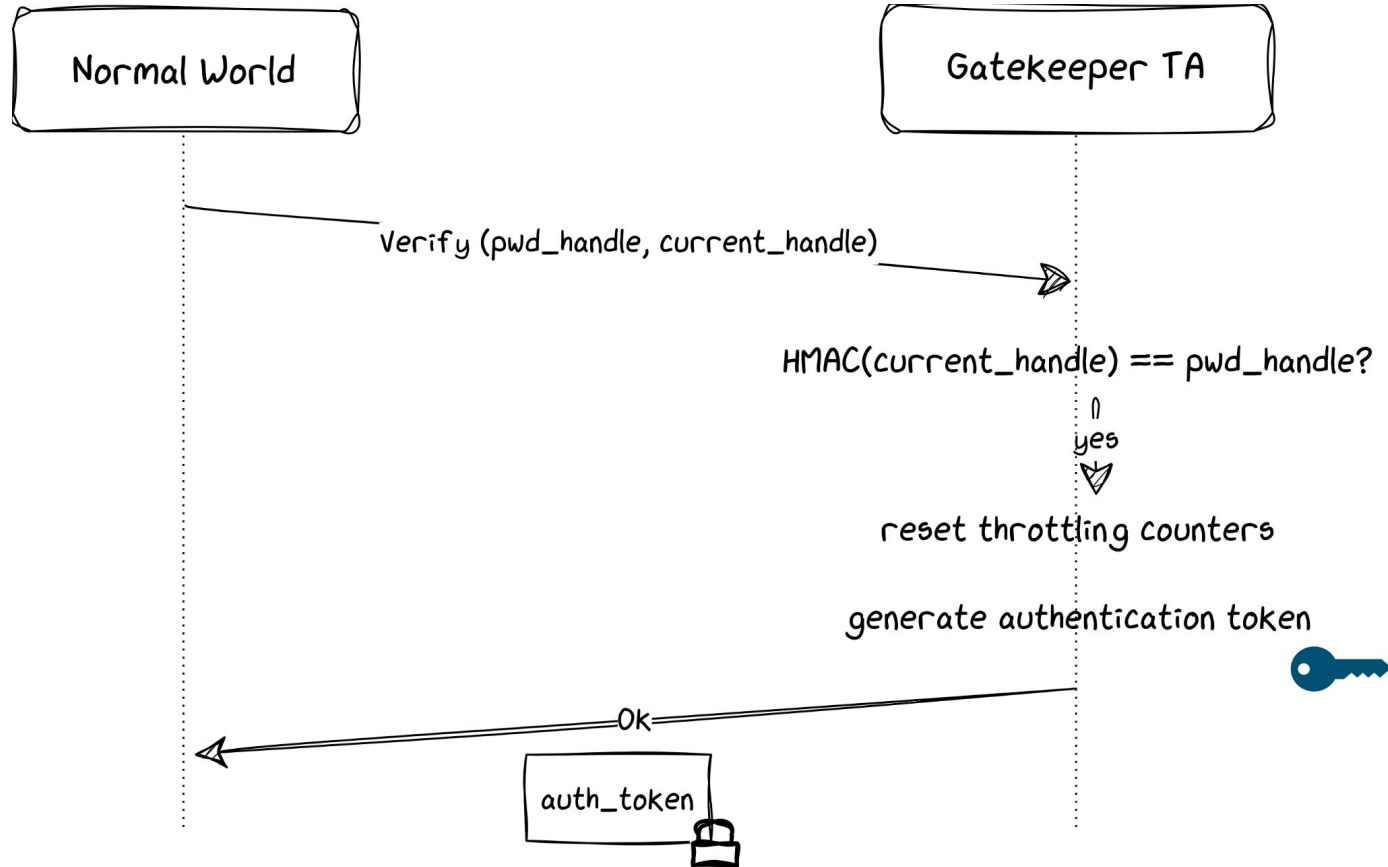


Authentication with Gatekeeper

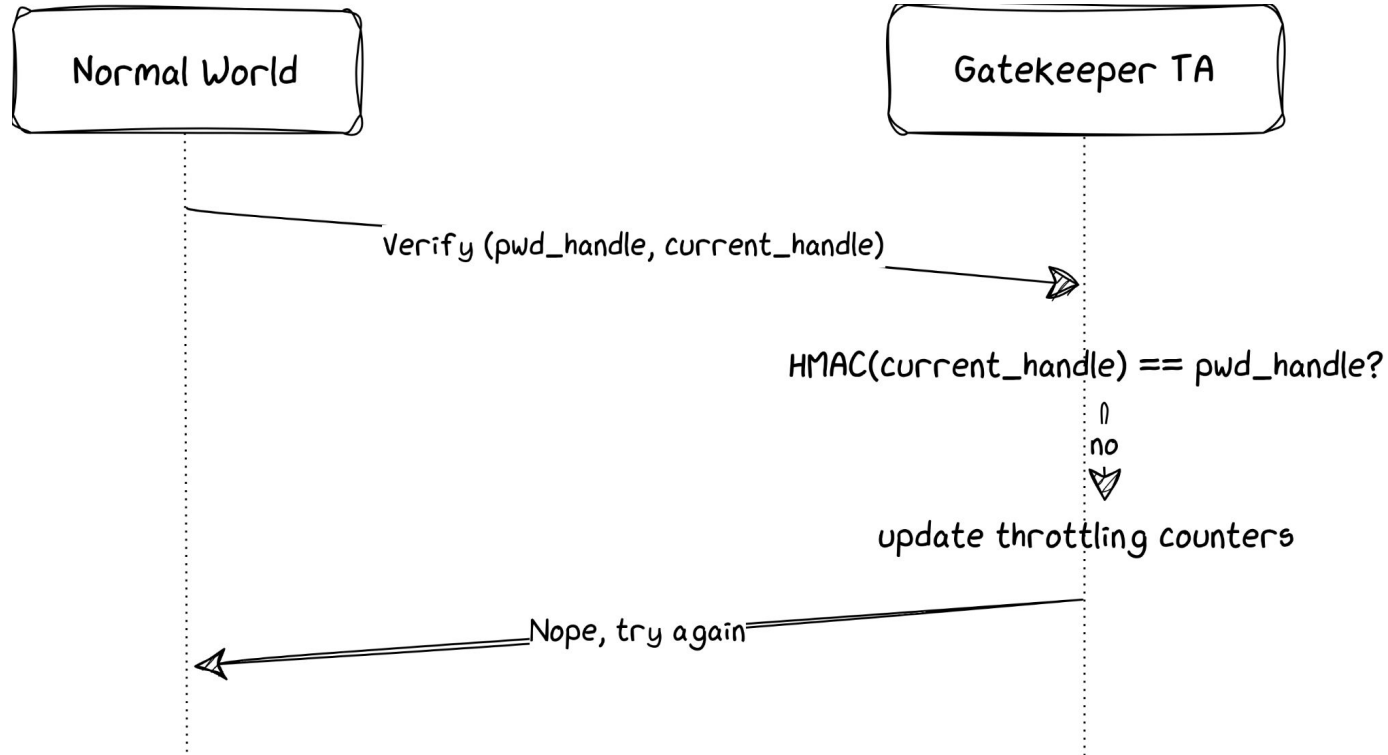
- The Gatekeeper TA verifies credentials from the TEE
- `/data/system_de/<uid>/spblob/<handle>.pwd`
 - script parameters
 - *password handle*, i.e. `HMAC(SHA512("user-gk-authentication" || script(credentials, params)))`
- If successful, Gatekeeper returns an *authentication token*
 - Signed token to be used to prove successful authentication
 - Needed by Keymaster to use authentication-bound keys
 - Standard format, designed not to allow replay attacks³
- Gatekeeper implements throttling to prevent bruteforcing

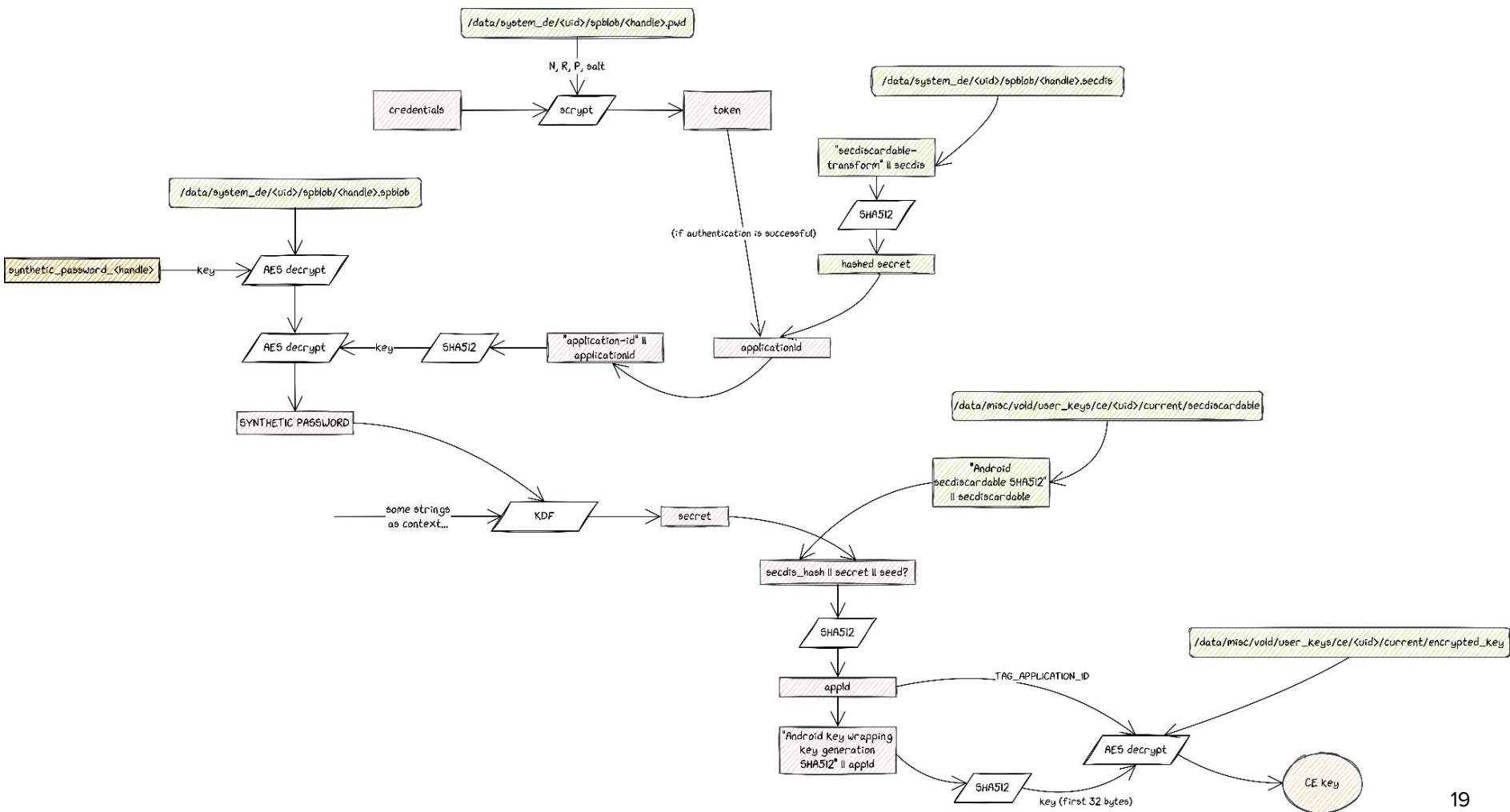
[3]: https://android.googlesource.com/platform/hardware/libhardware/+master/include/hardware/hw_auth_token.h

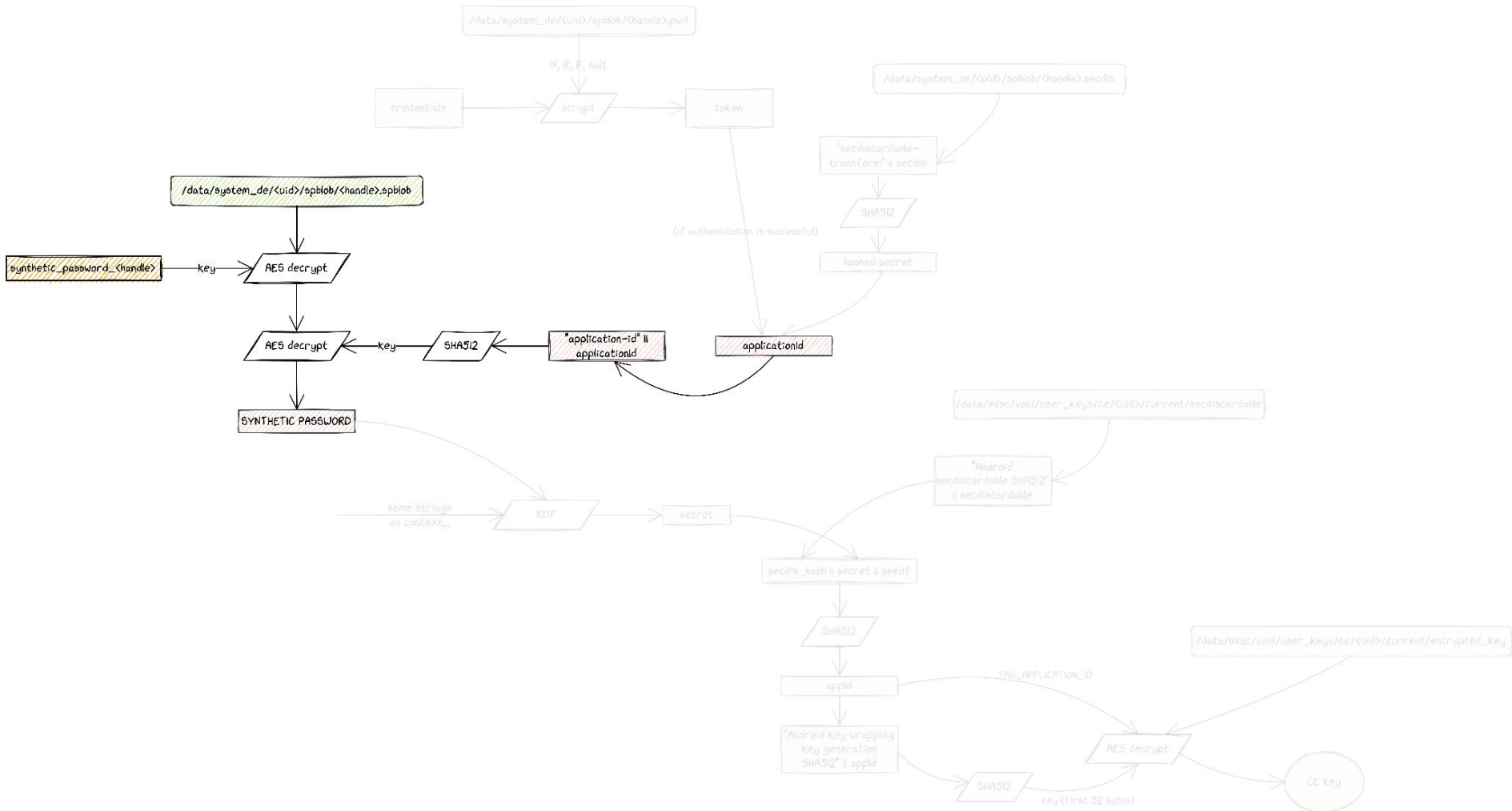
Successful authentication



Failed authentication

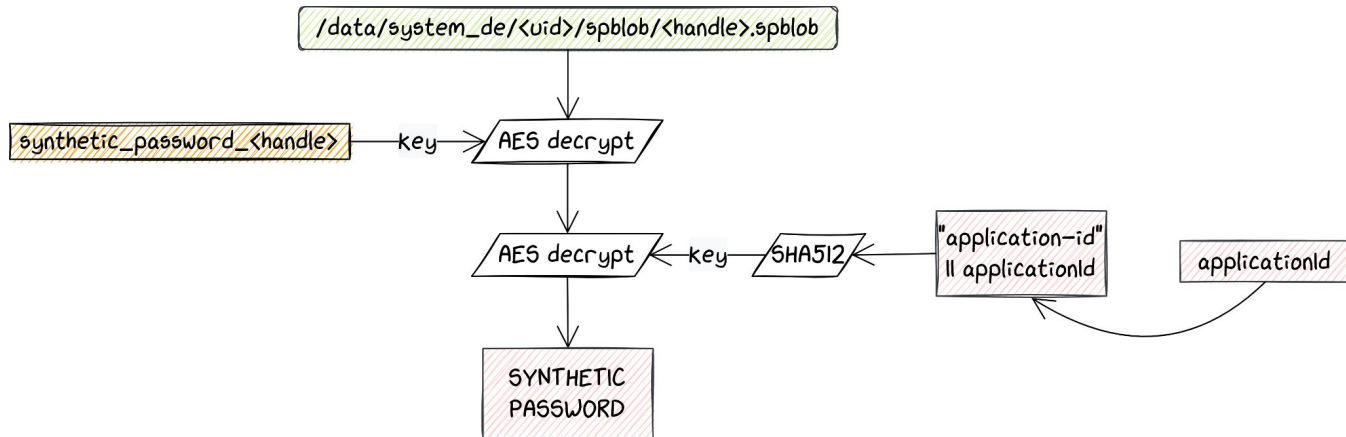






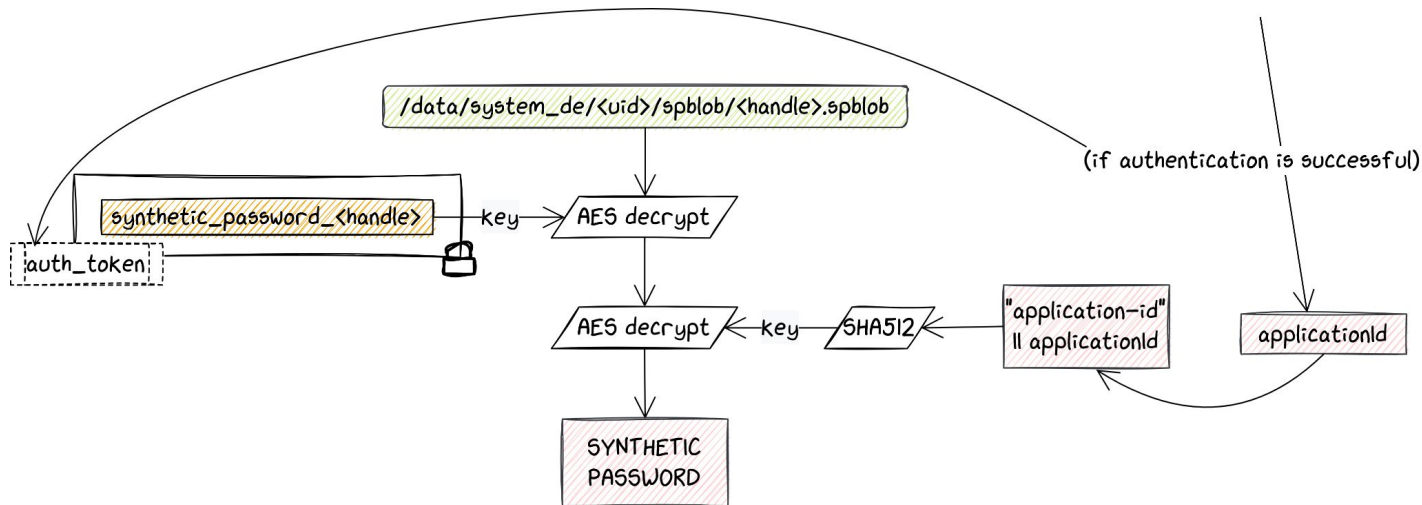
Synthetic Password

- Problem: credentials shouldn't be linked to the CE key
 - What if the user changes them?
- Solution: Synthetic Password
 - Key blob stored in `/data/system_de/<uid>/spblob/<handle>.spblob`
 - First, decrypted with an authentication-bound, TEE-protected key
 - Then, decrypted with the (hashed) applicationId



Attacking SP derivation

- Need to target the TEE
- Two alternatives
 - Keymaster TA (accessing the first AES key)
 - **Gatekeeper TA** (validating credentials and minting auth tokens)



Global strategy

- Our goal
 - Root the device and access all the device encrypted files
 - Patch the Gatekeeper trustlet to accept any credentials
- For that we need
 - Either multiple bugs (code exec, priv esc, etc)
 - Or one critical bug early in the boot process

PoC on Samsung Device

- Samsung A225f and A226b
 - Cheap (~250€)
 - Mediatek SoC MT6769V and MT6833V
 - No security chip
 - Mix of Mediatek and Samsung code
 - Trustzone OS: TEEGRIS
 - Known critical Boot ROM vulnerability



The Boot ROM Known Vulnerability

We use the project MTKClient⁴ (by Bjoern Kerler – [@viperbjk](#))

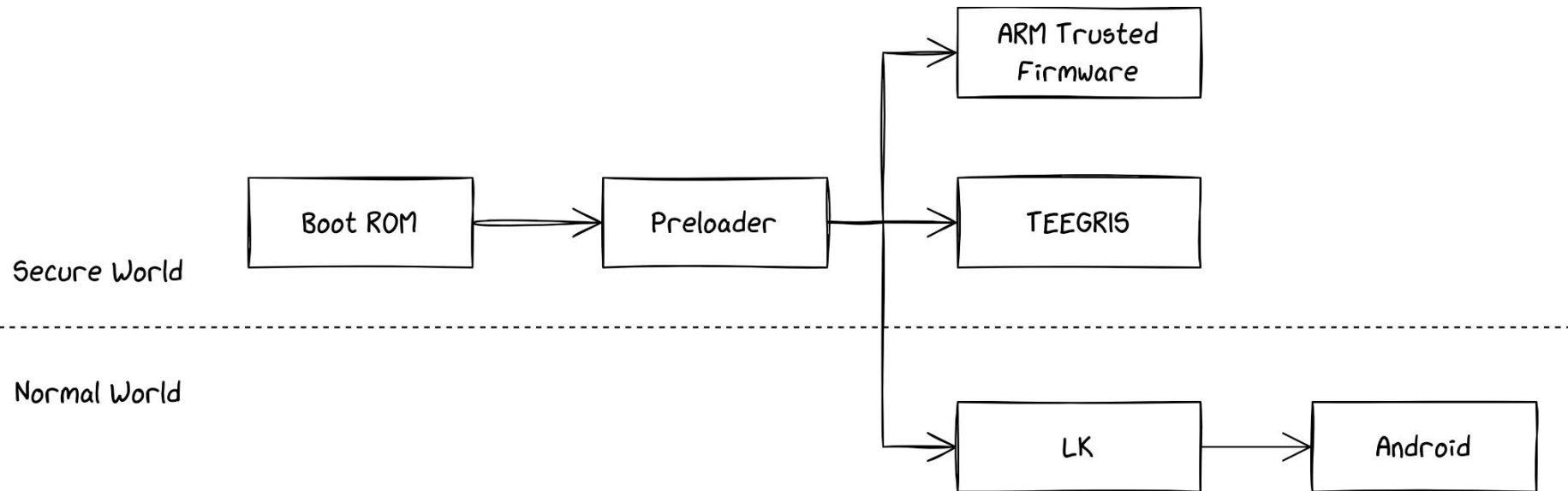
- Exploit boot ROM bugs impacting plenty of Mediatek SoC

In short, we use it to

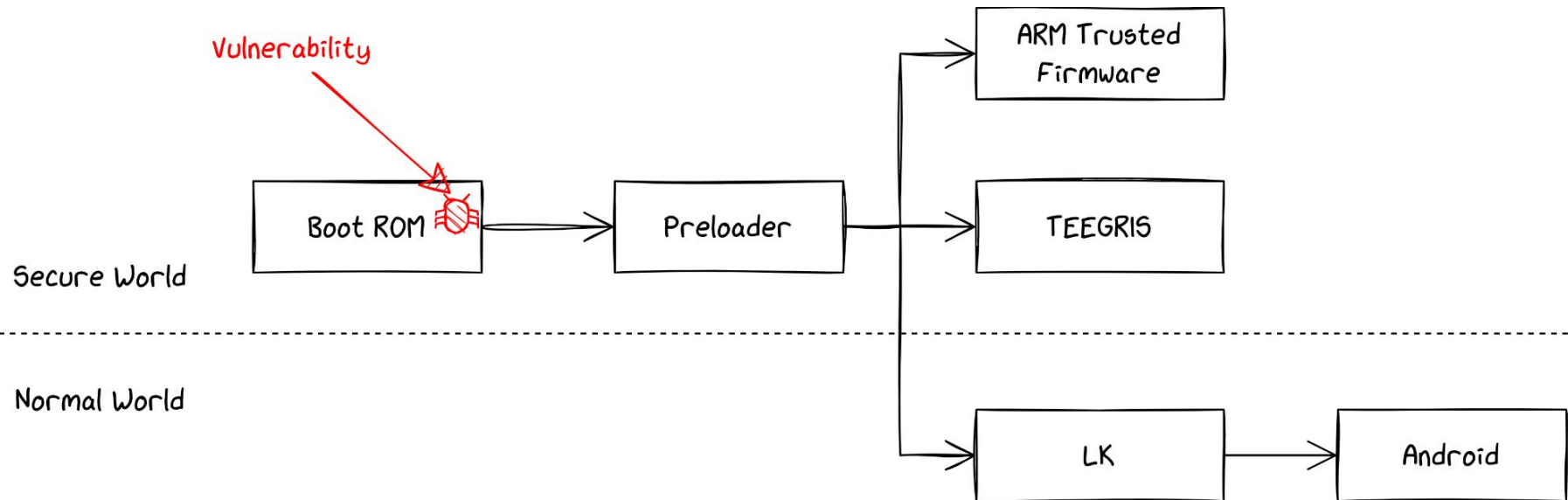
- Read/write all the partitions we need to patch
- Boot a patched preloader (BL2) image
- Bypass the secure boot checks done in boot ROM and preloader
- It just works :)

[4]: <https://github.com/bkerler/mtkclient>

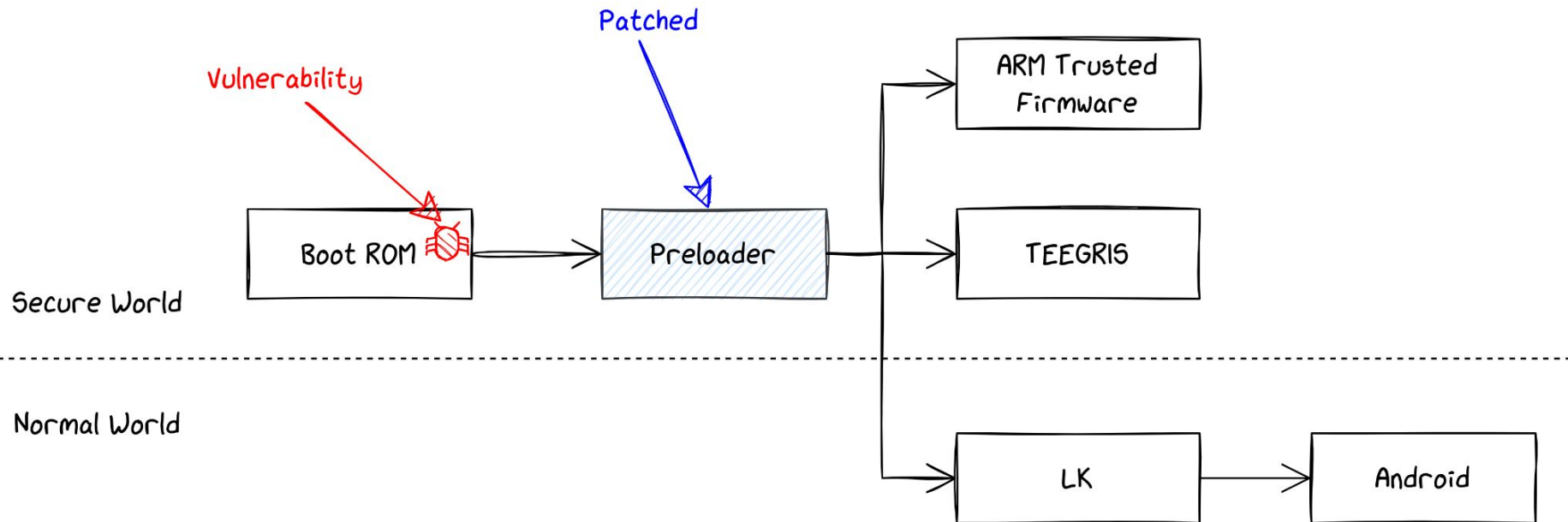
Boot Process



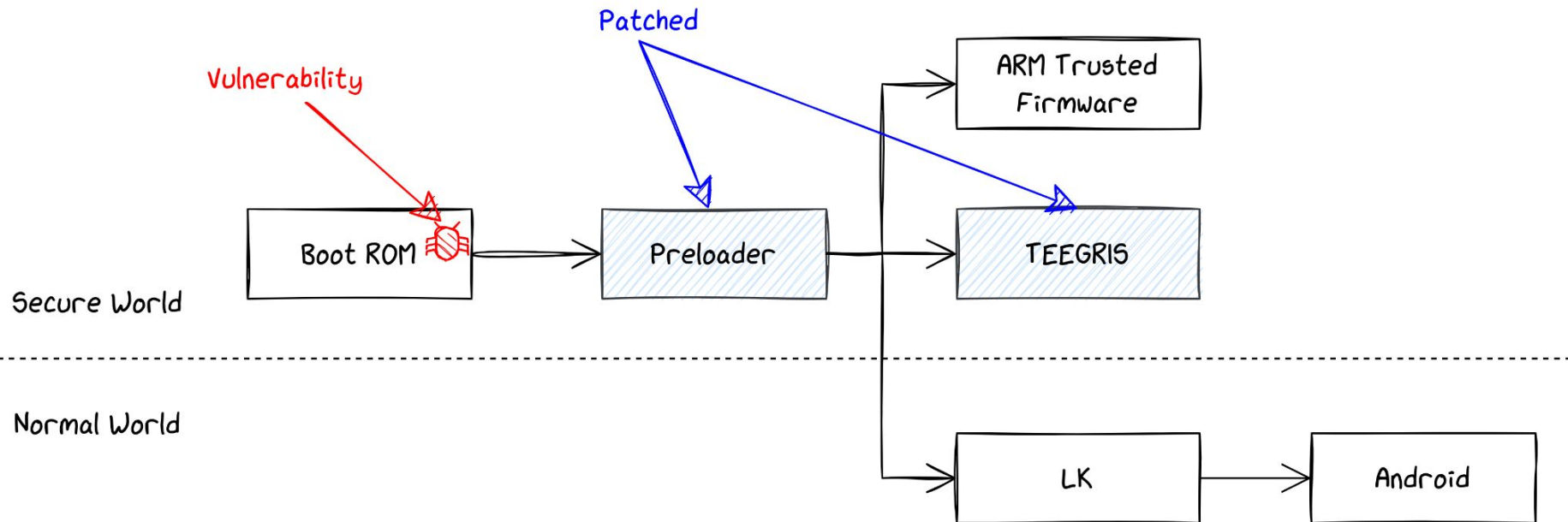
Boot Process



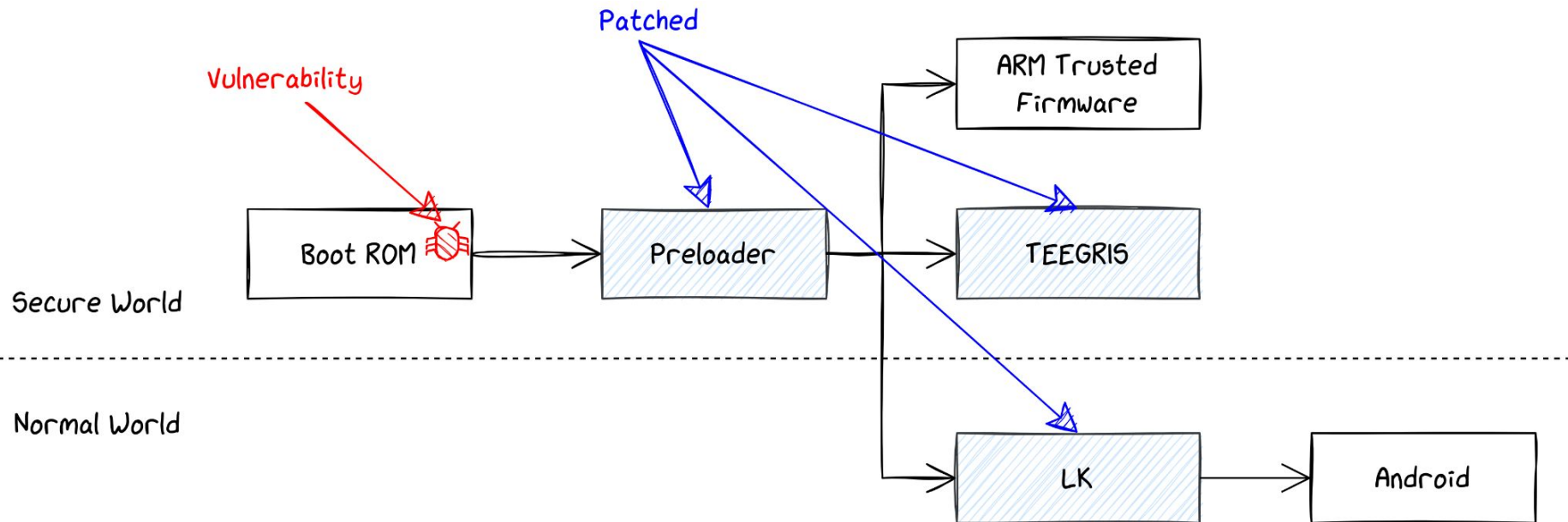
Boot Process



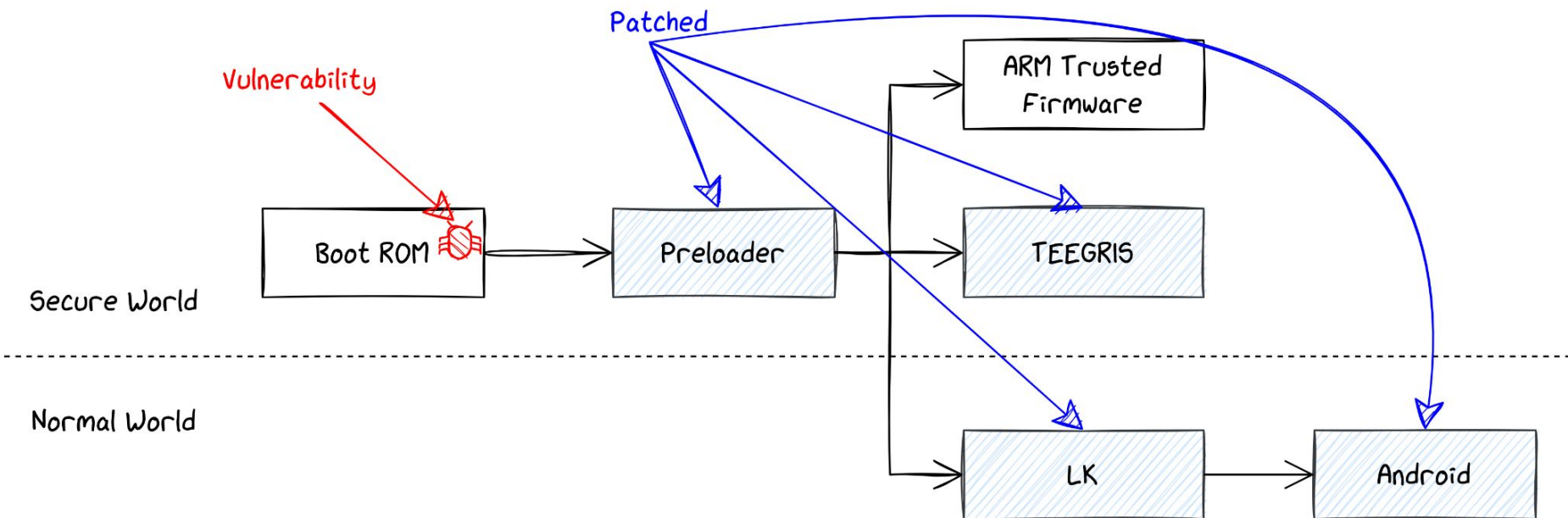
Boot Process



Boot Process



Boot Process



Little Kernel Patching

- Patching strategy: empirical approach
 - Reverse engineering and identify checks
 - Patch, test and repeat
- In the end we patch AVB to launch a modified boot image



Security Error 系统错误

This phone has been flashed with unauthorized software & is locked. Call your mobile operator for additional support. Please note that repair/return for this issue may have additional cost.

本机由于安装了未授权的软件而被锁定,请前往就近的售后服务中心寻求帮助,届时所发生的维修费用有可能需要自行承担,请知悉

Little Kernel Patching

```
26  iVar1 = do_hash(param_1,param_2,DAT_4c6463e0 - param_2,&hash,0x20);
27  if (iVar1 == 0) {
28      iVar2 = memcmp(&STORED_HASH,&hash,0x20);
29      if (iVar2 == 0) {
30          print("[%s][oem] img auth pass\n",&s_SBC_030151a8);
31          goto LAB_02ff82e0;
32      }
33      iVar1 = 0x7021;
34  }
35  print("[%s][oem] img auth fail (0x%x)\n",&s_SBC_030151a8,iVar1);
```

Little Kernel Patching

```
28  iVar1 = do_hash(param_1,param_2,_DAT_4c6463e0 - param_2,&hash,0x20);
29  if (iVar1 == 0) {
30      iVar2 = memcmp(&hash,&hash,0x20);
31      if (iVar2 == 0) {
32          print("[%s][oem] img auth pass\n",&DAT_030151a8);
33          goto LAB_02ff82e0;
34      }
35      iVar1 = 0x7021;
36  }
37  print("[%s][oem] img auth fail (0x%x)\n",&DAT_030151a8,iVar1);
```

Rooting Android

Main partitions used by Android: **boot** and **super**

- Boot contains the kernel and a ramdisk (only used for first boot stage)
- Super is a Dynamic Partition that contains 4 logical partitions
 - system, vendor, product, odm

To root it

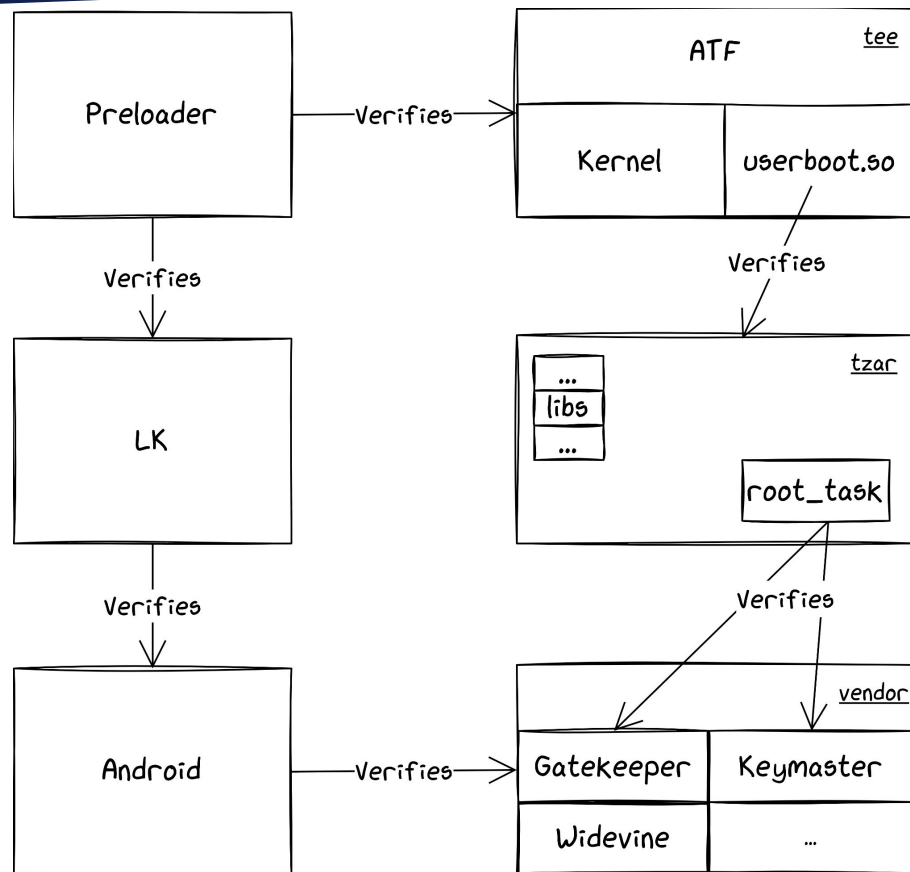
- Magisk⁵ to patch the boot image
- We made few modifications to su
- Plus other little tricks to patch the super partition

[5]: <https://github.com/topjohnwu/Magisk>

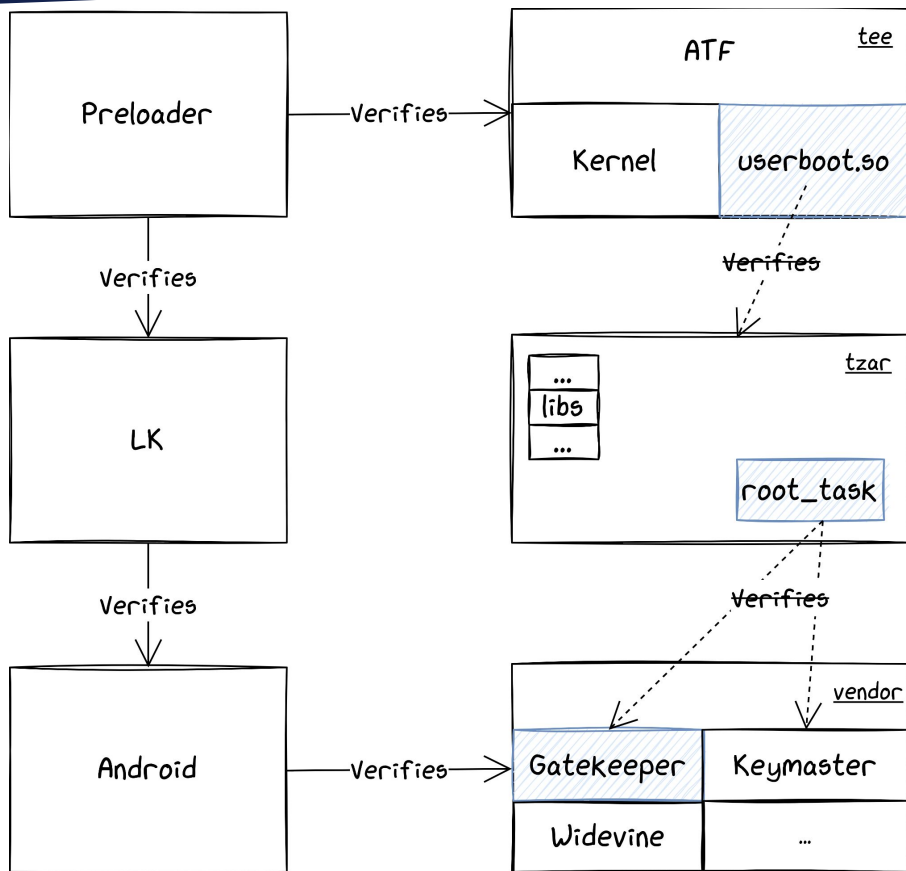
- Trustzone OS designed by Samsung
- For Mediatek and Exynos SoCs
- ROM images:
 - tee1.img: ATF, TEEGRIS kernel, userboot.so
 - tzar.img: TEE root filesystem
 - super.img: Android system, Trusted Applications and Drivers
- Excellent references online⁶

[6]: <https://www.riscure.com/tee-security-samsung-teegris-part-1/>

TEEGRIS Images Verification



Patching TEEGRIS



Reversing Gatekeeper

- TAs come in a slightly modified ELF format
 - 8-bytes header and footer with signature
 - Removing them allows to load a nice ELF in your favourite disassembler
- GlobalPlatform API
 - Standard API for TEEs (memory allocation, crypto operations, etc.)
 - Makes reversing easier
- Trusty reference implementation⁷
 - Suggests what to expect from a TA

[7]: <https://source.android.com/docs/security/features/trusty>

Gatekeeper Reference Implementation

- 2 Gatekeeper commands: Enroll and Verify
- Verify does two things:
 - `HMAC(pwd_handle) == expected?`
 - If so, create new authentication token
- What if we can leak the key used by HMAC?
 1. `pwd = generate new password`
 2. `Value = HMAC(pwd_handle)`
 3. `Value == expected`

Reversing & patching Gatekeeper

- 2 Gatekeeper commands: Enroll and Verify
- Verify does two things:
 - `HMAC(pwd_handle) == expected?`
 - If so, create new authentication token



Reversing & patching Gatekeeper

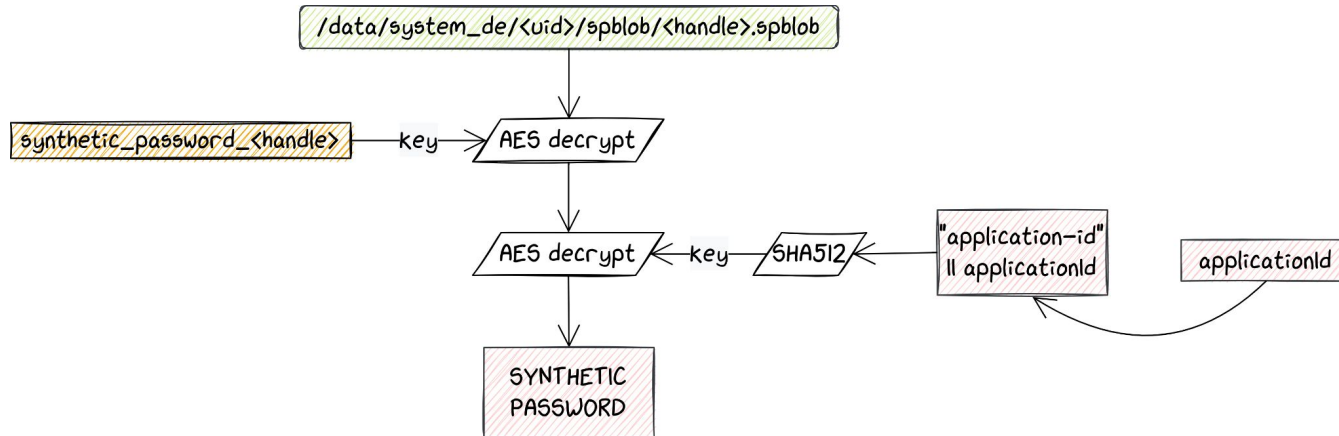
- This Gatekeeper implementation uses a KDF instead of a plain HMAC
 - KDF implemented in a library
 - which calls /dev/crypto
 - many steps to leak the key
- Simpler strategy: patch to accept any credentials
- Always return valid auth token to continue the process
 1. ~~KDF(pwd_handle) == expected?~~
 2. ~~If so,~~ create new auth_token

Reversing Gatekeeper

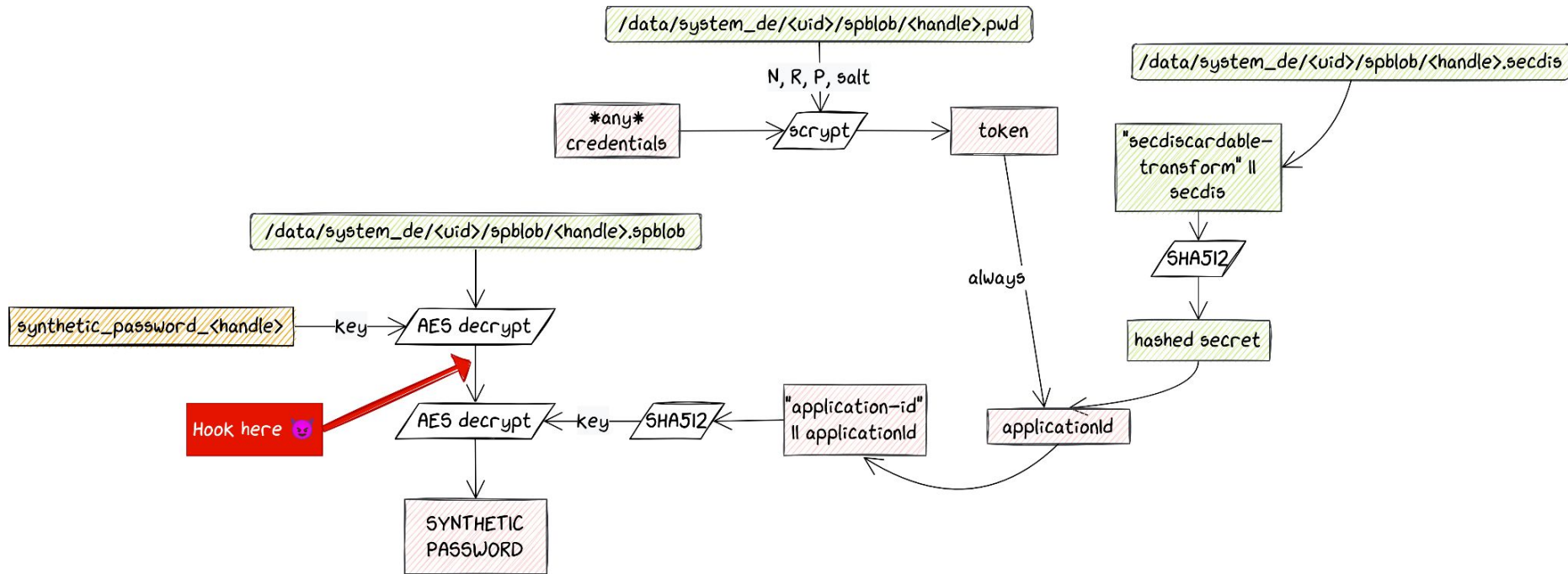
```
22  iVar1 = TEE_AllocateOperation(&local_30,0x50000004,5,0);
23  if (iVar1 == 0) {
24      iVar1 = TEE_DigestDoFinal(local_30,param_1,param_2,auStack_28,&local_38);
25      TEE_FreeOperation(local_30);
26      if (iVar1 == 0) {
27          uVar2 = TEE_AllocateTransientObject(0xa0000000,param_4 << 3,&local_30);
28          if (uVar2 == 0) {
29              uVar2 = TEES_DeriveKeyKDF(auStack_28,local_38,local_48,8,param_4,local_30);
30              if (uVar2 == 0) {
31                  uVar3 = 1;
32                  uVar2 = TEE_GetObjectBufferAttribute(local_30,0xc0000000,param_5,&iStack_34);
33                  if (uVar2 != 0) {
34                      uVar3 = 0;
35                      printf("gatekeeper [ERR] (%s:%u) failed to get object attribute: %x","hal_pwd_hmac",
36                          0x12a,(ulong)uVar2);
37              }
```

Attack strategy

- Read the output of the first AES decrypt
- Bruteforce credentials to generate applicationId
- Thanks to GCM mode, AES decrypt complains if the key is wrong



Hooking system_server



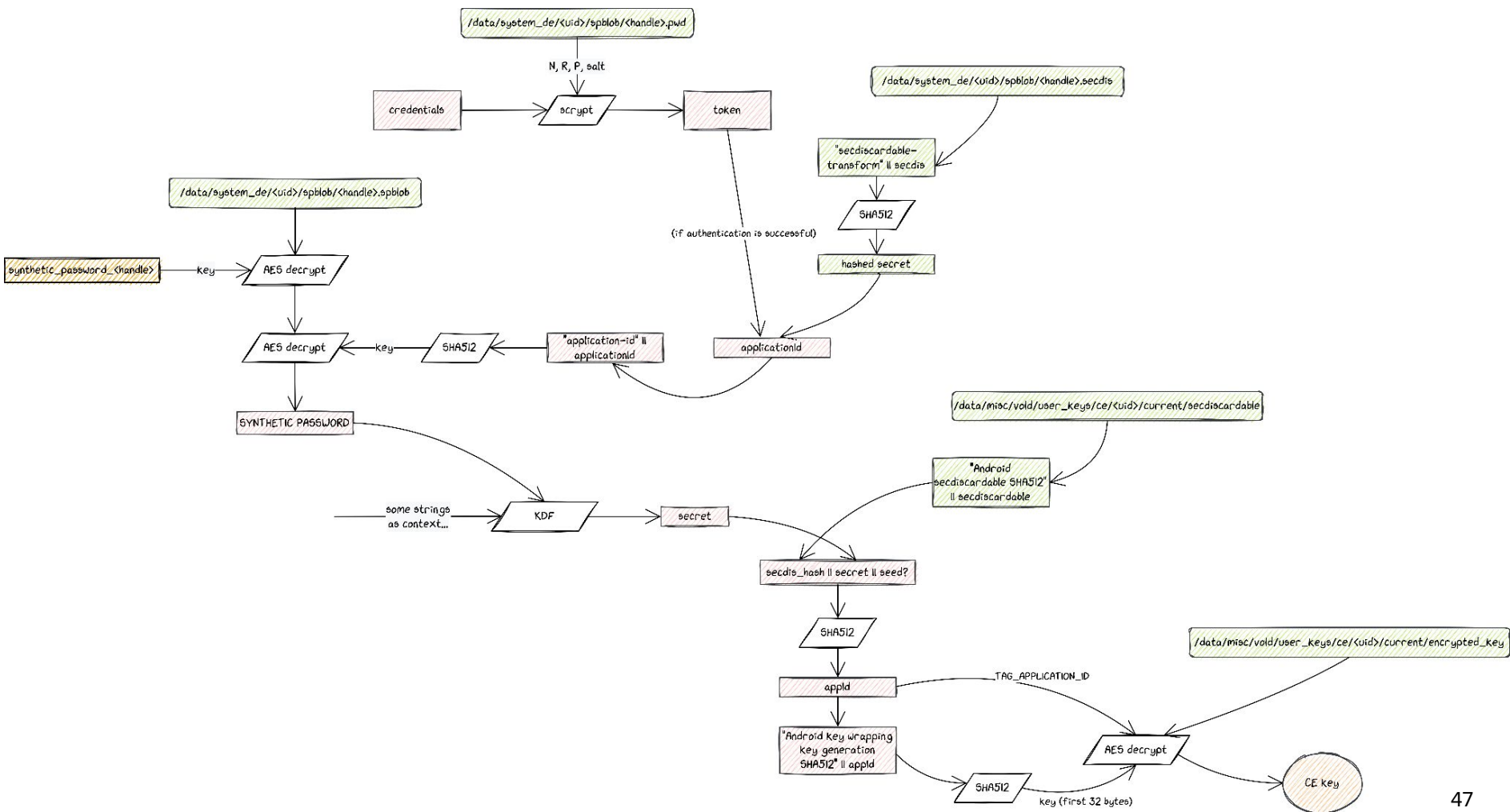
Retrieving intermediate key with Frida

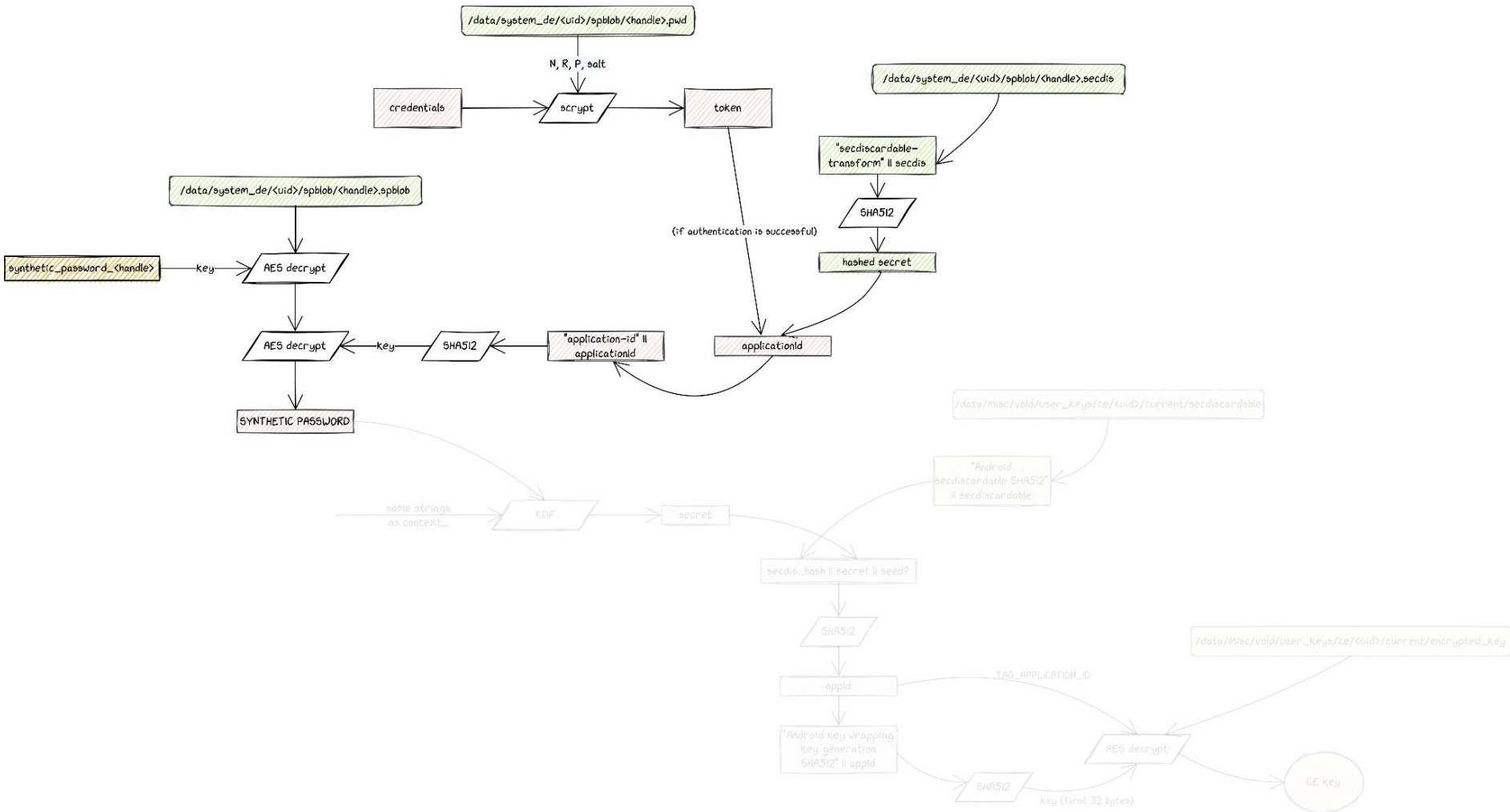
- Use Frida to hook system_server
- Retrieve intermediate buffer decrypted by TEE
 - Possible thanks to the auth token

```
$ frida -U system_server -l system_server.js

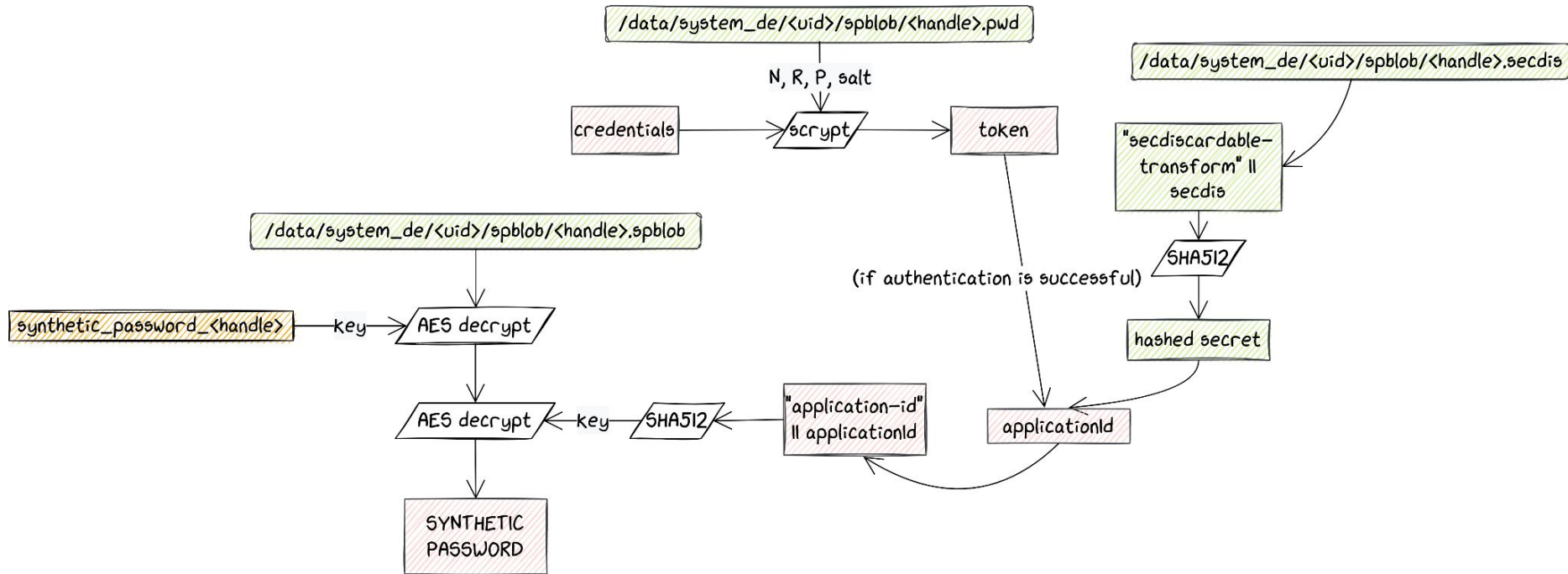
  ____
 /_  _|   Frida 16.0.19 - A world-class dynamic instrumentation toolkit
| (  |
 >_  |   Commands:
/_/_|_|   help      -> Displays the help system
. . . .   object?   -> Display information about 'object'
. . . .   exit/quit -> Exit
. . . .
. . . .   More info at https://frida.re/docs/home/
. . . .
. . . .   Connected to SM A226BR (id=R9WTA0BYDPL)

[SM A226BR::system_server ]-> SyntheticPasswordCrypto.decrypt called!
ciphertext = 641a3ed0a68abdae99976b5aff32f8d5aa4d18127272af6ff638c1e88d57cd157fd6f75b4688465f
470bd4cc81081215e9f2085e4b8ea22e0e8f0ed32a381f641d5cd071d2e177c4a8a1b6e6824f52f251366ff730f66
b7cfd72f11f9761efc5e0cf68bd7bdec00456e07dfb9f1a7f720e97aa262c0507bc87ef46e603a265c821cb1a1dc5
c6f6be6fd43ac3431d0d013de8c9
[SM A226BR::system_server ]->
[SM A226BR::system_server ]-> |
```

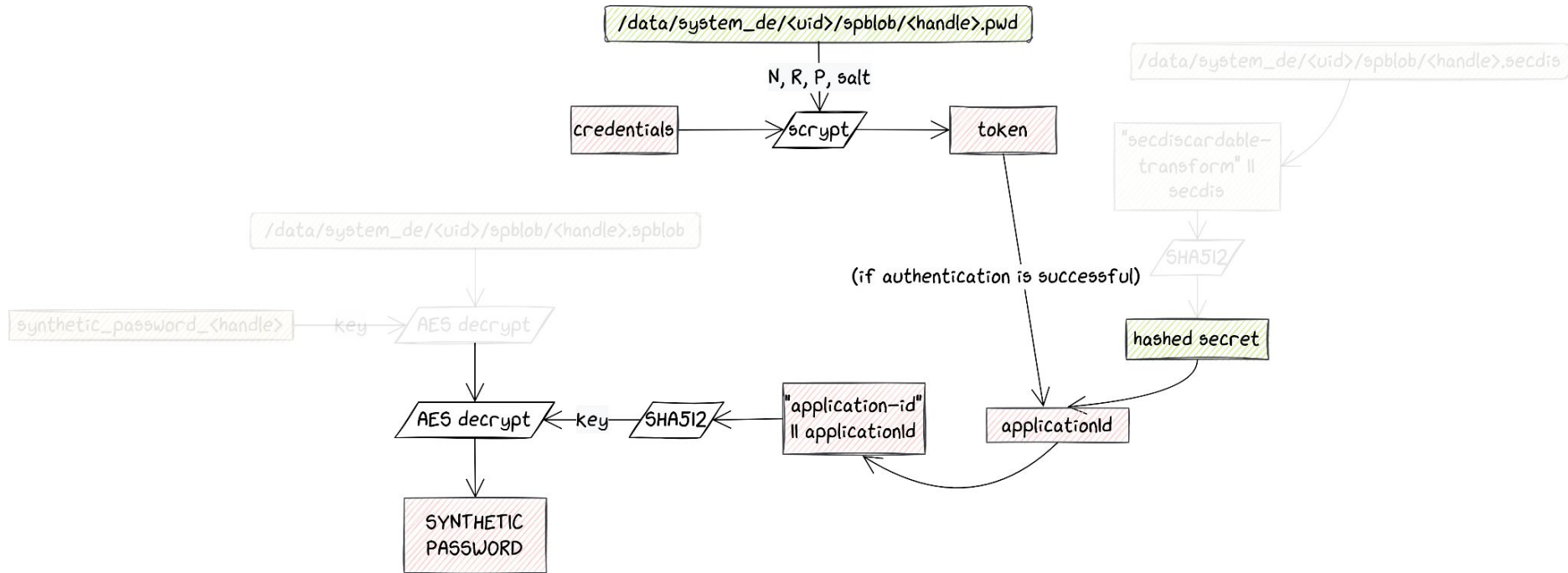




Bruteforce of the password



Bruteforce of the password



Bruteforce of the password

1. `pwd = generate new password`
2. `token = scrypt(pwd, R, N, P, Salt)`
3. `Application_id = token || Prehashed value`
4. `Key = SHA512("application_id" || application_id)`
5. `AES_Decrypt(value_from_keymaster, key)`

Bruteforce of the password

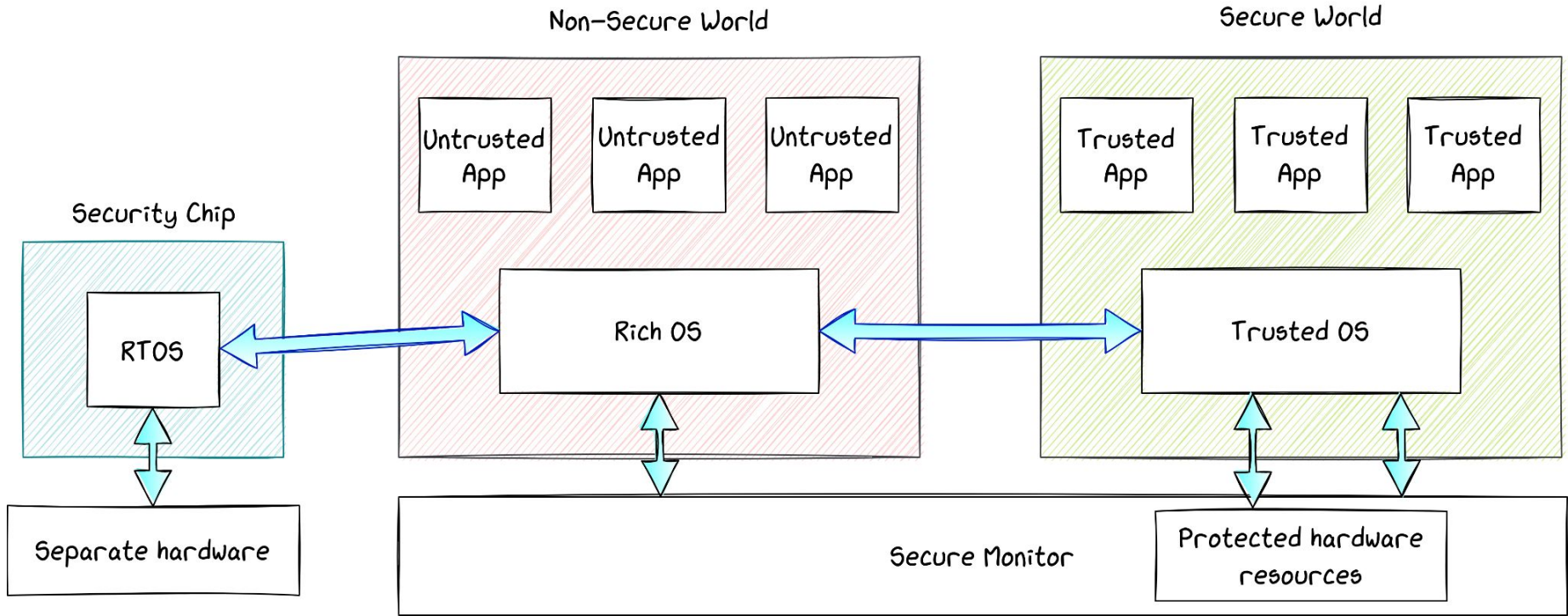
1. `pwd = generate new password`
2. `token = scrypt(pwd, R, N, P, Salt)`
3. `Application_id = token || Prehashed value`
4. `Key = SHA512("application_id" || application_id)`
5. `AES_Decrypt(value_from_keymaster, key)`

```
$ python3 bruteforce-tee.py
workers will cycle through the last 5 chars
Found it: 1234
the plaintext is '1234'
Done in 18.031058311462402s
Throughput: 1478.448992816657 tries/s
```

Demo 1



Architecture w/ Trusted Chip



The Titan M Chip

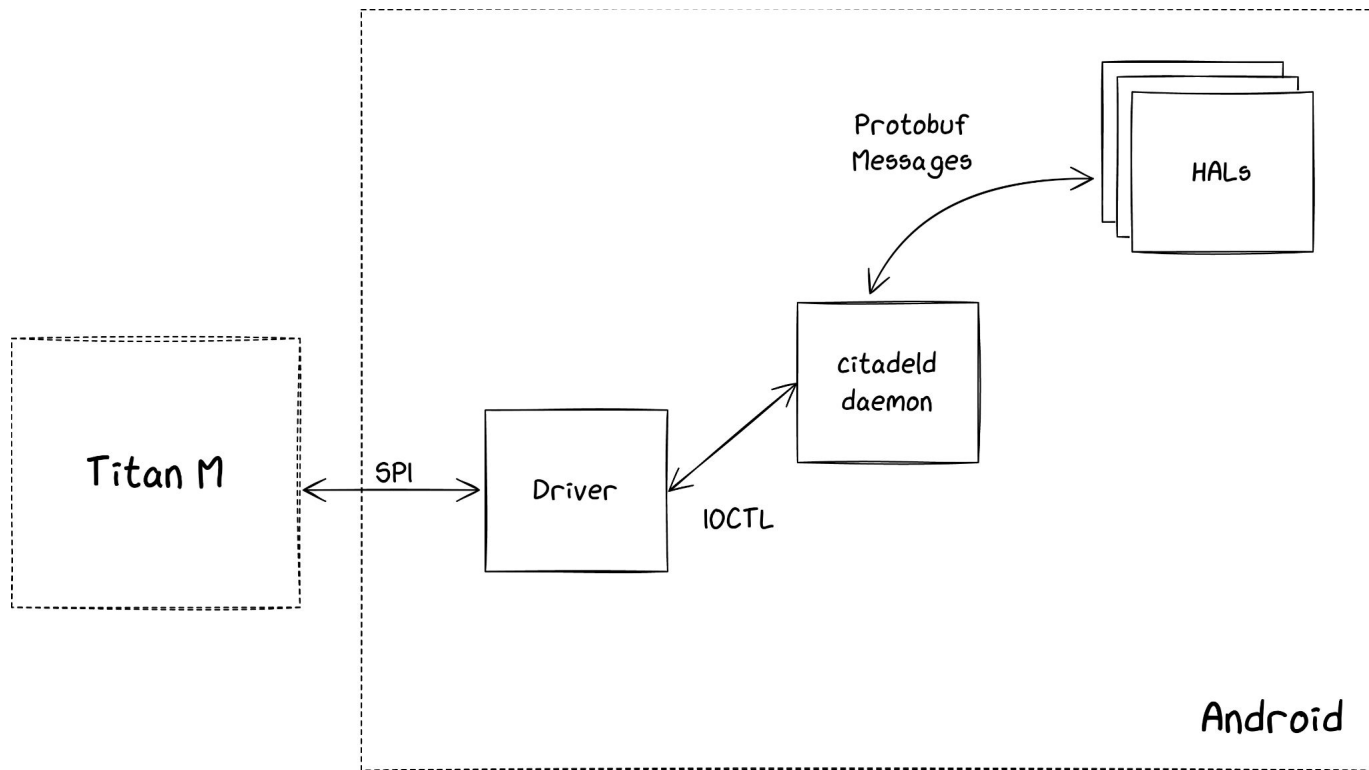
- Security chip made by Google for Pixel phones
- From Pixel 3 to Pixel 5a
 - In this PoC we use a Pixel 3a
 - Titan M2 introduced from Pixel 6
- Based on Arm Cortex-M3
- Most of the code is divided into tasks
 - Keymaster (Strongbox), **Weaver**, AVB, etc
- Separate memory and resources
 - Communicates with Android on SPI bus



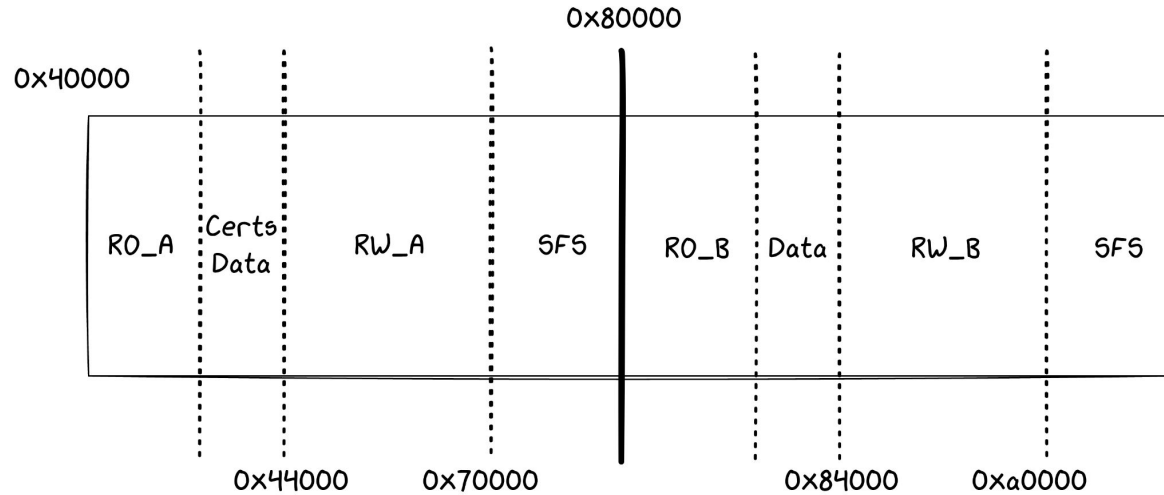
Trusted chip vs TrustZone

- In TrustZone, secure and normal world run on the same CPU
 - Shared hardware (cache, RAM)
 - Side-channel attacks are possible (e.g. Rowhammer)
- Titan M relies on tamper-resistant hardware
- Separate firmware
 - Limited in size
 - Conceptually simple
 - Isolated from the rest of the system

Communication with the chip

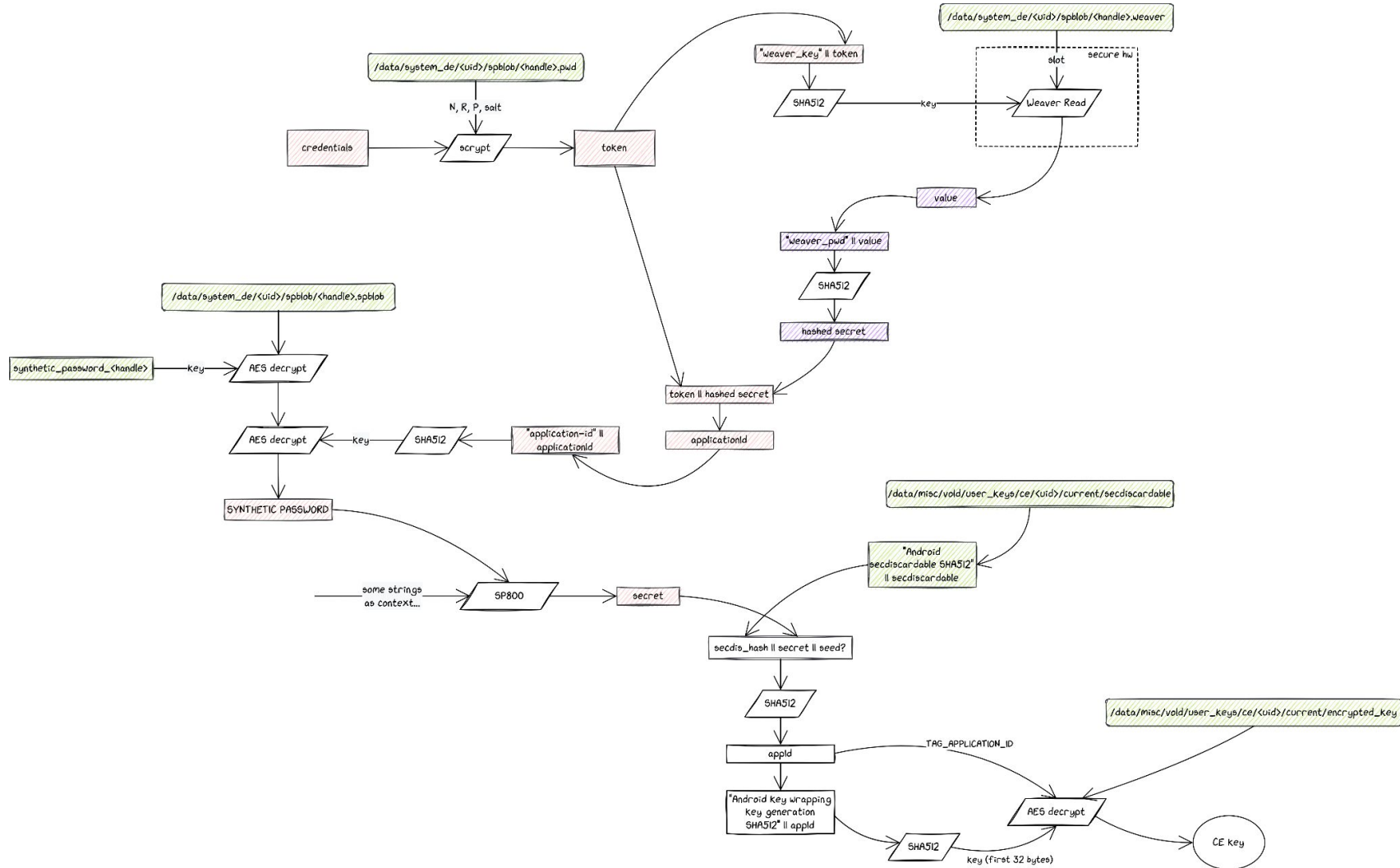


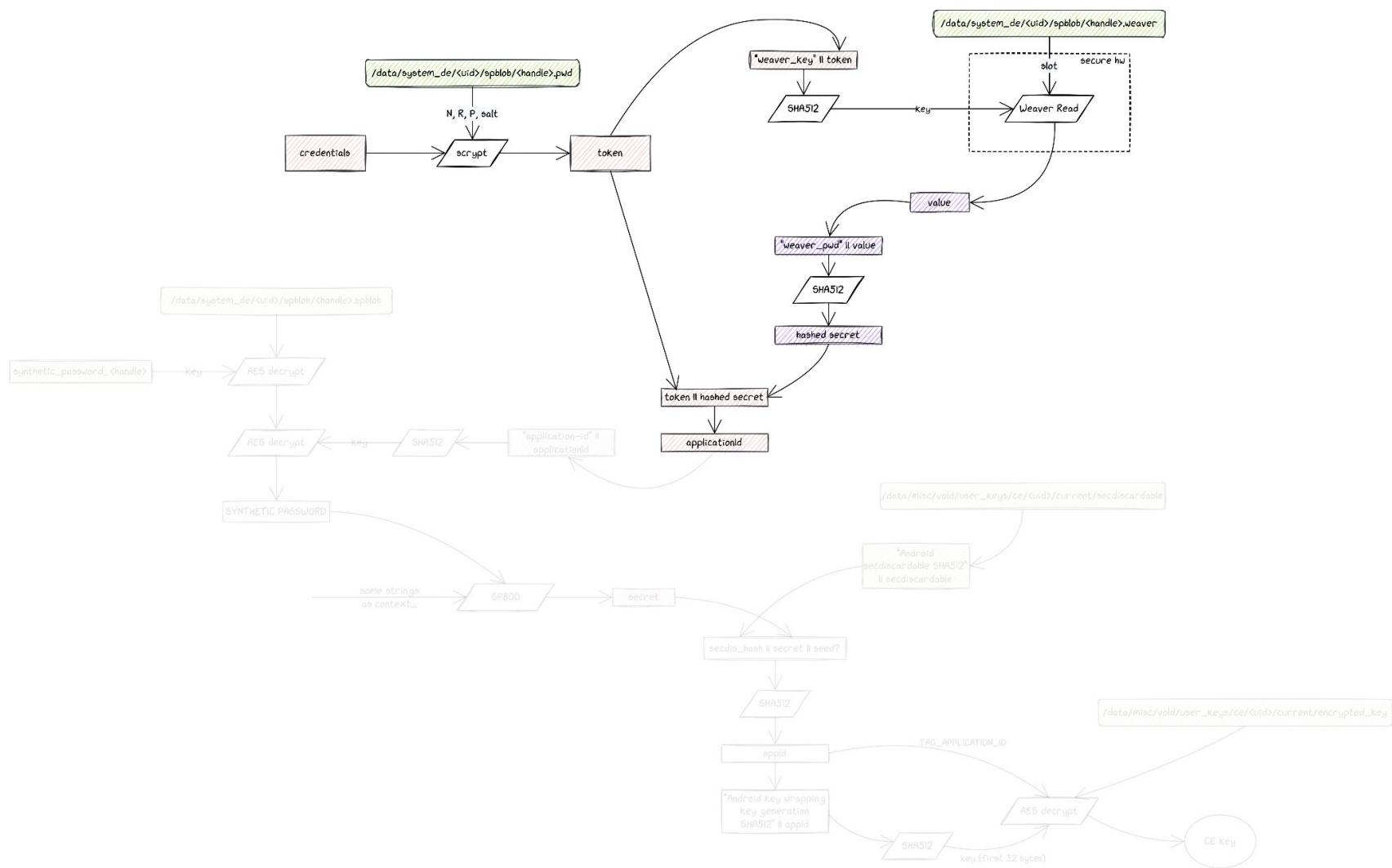
Memory Layout



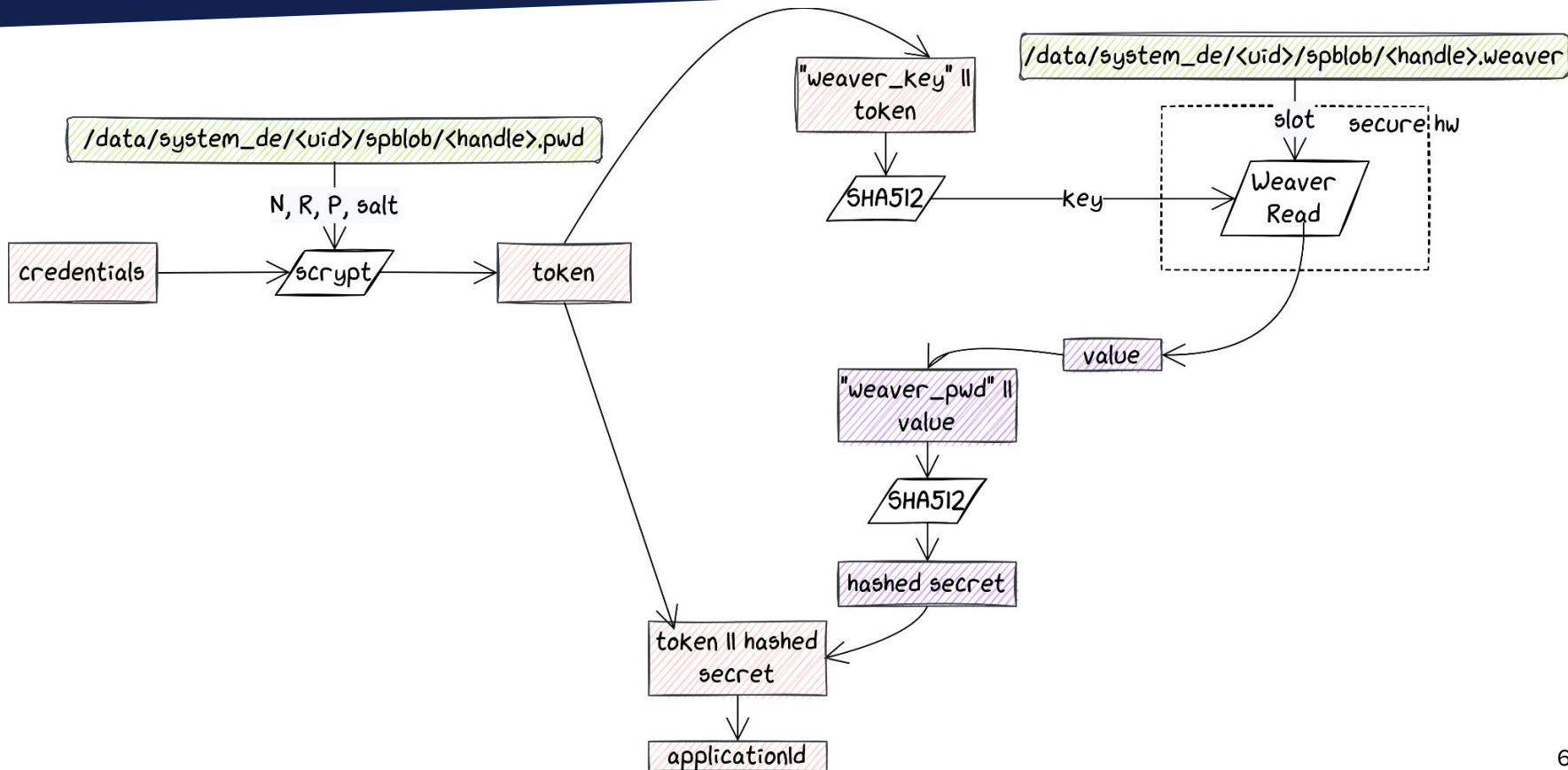
- Key/Value storage
 - Stored in slots
 - In two different places in the flash memory
- 4 commands: GetConfig, Read, Write, Erase
- Implements throttling as well

```
// Read
message ReadRequest {
    uint32 slot = 1;
    bytes key = 2;
}
message ReadResponse {
    Error error = 1;
    uint32 throttle_msec = 2;
    bytes value = 3;
}
```





CE key derivation with Weaver



PoC on Google Pixel

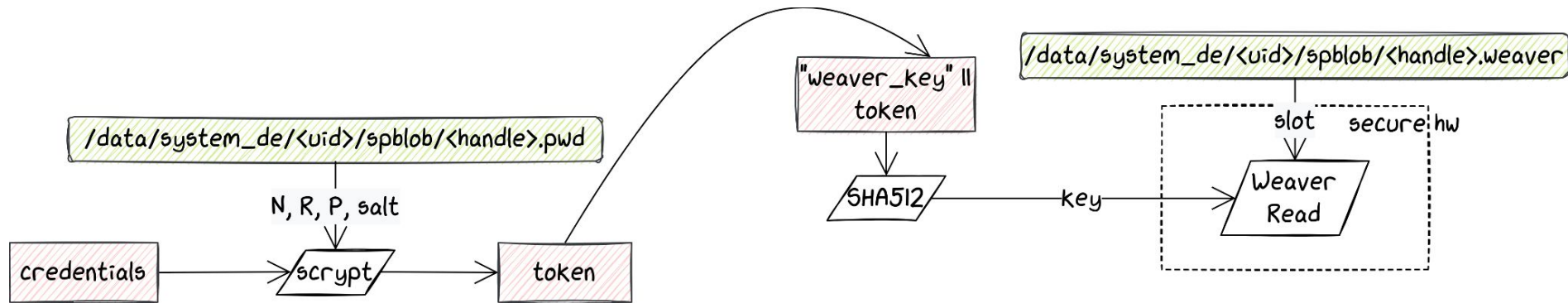
- We consider the device being already rooted
- Weaver relies on the security chip Titan M
- Here we exploit CVE-2022-20233 to execute code on the chip
- Out-of-bounds write of 1 byte to 0x1
 - Can be repeated multiple times
 - Huge constraints on the offset
 - We managed to overwrite a global field and cause another corruption
- Full exploit write-up in our blog⁸

[8]: <https://blog.quarkslab.com/attacking-titan-m-with-only-one-byte.html>

Nosclient and the leak command

- We built a client to communicate with Titan M, nosclient
- “Leak” feature:
 - `./nosclient leak <address> <size>`
 - Read <size> bytes from <address>
 - ➔ Arbitrary read in Titan M’s memory
- Weaver slots and values are stored in flash
 - Reverse engineering to understand a memory range
 - Then search for 16 bytes digests
 - Weaver Write and Read help out

Our Strategy



1. Leak the Weaver key
2. Use it to compare our generated credentials

Bruteforce of the password

1. `pwd = generate new password`
2. `token = scrypt(pwd, R, N, P, Salt)`
3. `key = SHA512("weaver_key" || token)`
4. Compare with leaked Weaver key

```
$ python3 bruteforce.py
workers will cycle through the last 5 chars
Found it: 1106
the plaintext is '1106'
Done in 15.063793659210205 s
Throughput: 1491.722504195411 tries/s
```

Demo 2



Conclusion

- FBE is very well designed
- Ingredients from “everywhere” are used to derive the key
 - Files owned by privileged users
 - TEE-protected keys
 - Weaver values (when available)
- Multiple bugs needed to break it
 - Or a very powerful one
- You still need to bruteforce credentials in the end
- “my very secret password example for REcon 2023”
will be hard to guess :)

Thank you!

contact@quarkslab.com

Quarkslab



@DamianoMelotti
@max_r_b

Little Kernel

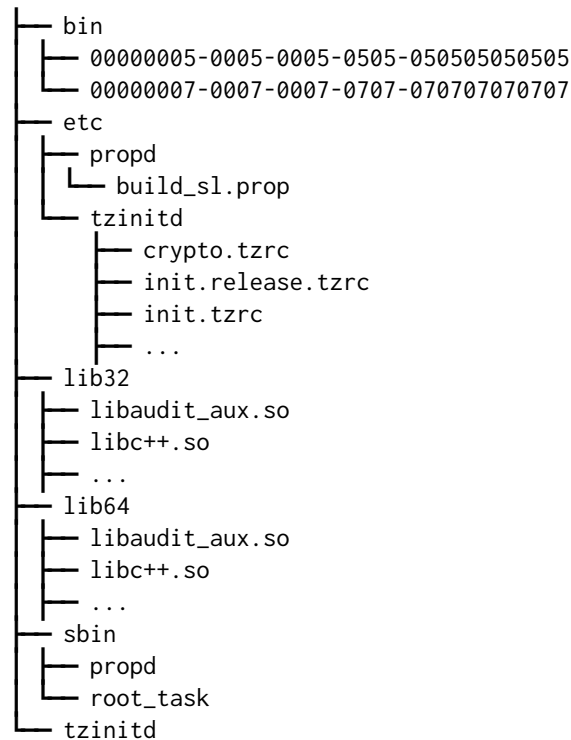
LK: Android bootloader based on Little Kernel

- Allows to boot Android or other modes (Recovery)
- Loads TZAR image in TEEGRIS
- Implements **A**ndroid **V**erified **B**oot v2
 - Verification of Android images
 - Involving boot and vbmeta images
 - Anti-rollback

TZAR image

TrustZone ARchive: contains a root filesystem

- Shared libraries
- Binaries
- tzinitd (init binary)
- root_task



Patching TEEGRIS

Our final goal is to run a modified Gatekeeper TA

- Patch `userboot.so` from the `tee1` partition
 - Disable verification of TZAR image
- Patch `root_task` from the TZAR image
 - Disable verification of TA
- Patch the Gatekeeper TA
 - Accept any credentials and return a valid auth token