

Creating Code Obfuscation Virtual Machines

VM Creation 101

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What is a VM?

Not VMWare, VirtualPC, etc.

Our own Custom Emulator

- Our own CPU
- Our own Language
- Our own Compiler (P-Code)

Why go through all this trouble?

- Code Obfuscation
- Hide Functionality and Intellectual Property
- Increases Analysis and Reversing Time
- Anti-Dumping Method

VM Benefits and Tricks

- The Virtual CPU is specialized for your tasks
- Built in Encryption
- Hidden Anti-Debugging Techniques
- VM Self-Modifying Code (SMC)
- Library or System Call obfuscation
- VM Junk Code

Who uses this?

DRM Does

- Themedia
- VMProtect
- BD+ (Blue-Ray)

Virus, Spyware, Exploits can use this as well

- Made popular by HoneyNet's SOTM32 (Nicolas Brulez)

Core Concepts

P-Code →

C8	D2	A1	FF	00	00	00	D9	CC	09
----	----	----	----	----	----	----	----	----	----

A0	POP VALUE INTO R1
A1	PUSH VALUE INTO R1
A2	POP VALUE INTO R2
A3	PUSH VALUE INTO R2
A4	POP VALUE INTO R3
A5	PUSH VALUE INTO R3
B0	POP VALUE INTO R4
B1	PUSH VALUE INTO R4

Handler Routine

Virtual Registers

r1, r2, r3, etc...

Our VM Crackme

Design for our VM Crackme:

- Core App runs and prompts user for password
- Pass password to our VM
- VM Does math on the password to make a “Key”
- “Key” is returned and used to Decrypt JMP to GoodBoy message
- Bonus Nugget: Null key is returned if password is wrong

What do we need to get started?

How many registers you want? **9**

How you want to control the program flow? **EIP**

How you going to handle memory? **ESP**

Macros or own custom language? **Custom**

What language will you write your compiler in?
Ruby

Virtual CPU Register Layout

Four General Purpose Registers:

r1, r2, r3, r4

Instruction Registers

IP, baseip

Stack Registers

SP, basesp

Flag Registers

flags

Our Virtual CPU Instruction Set

MOV r32, r32

MOV [r1], r32

MOV r1, [r1]

MOV r32, value

CMP r32, value

INC/DEC r32

AND/OR r1, value

XOR r32,r32

PUSH/POP r32

JMP (Relative / Direct) JE, JL, JG

CALL (r1 / value)

EXITVM

Virtual CPU Initialization

miniVm proc

```
pop ebx
pop eax                ; Stack Argument
mov [stack],eax
pop eax                ; P-Code
mov [ip],eax
mov [baseip],eax
pusha                  ; Save Registers
mov [flags],0         ; Init some regs
mov [stackp],0
call _core             ; State Machine
popa                   ; Restore
push ebx              ; Jump back to code
ret
```

Calling our CPU

- `invoke SendMessage, PasswordHandle, WM_GETTEXT, 20, addr hPassword`
- `mov eax, offset mystack`
- `mov [eax], offset hPassword`
- `push offset vmcode`
- `push offset mystack`
- `call miniVm`

Our Opcode Processor (State Machine)

_next_ip:

```

mov ebx,[ip]
xor eax,eax
xor ecx,ecx
xor edx,edx
mov al,byte ptr [ebx] ; al = instructional opcode
mov dl,al
and dl,0Fh ; Major opcode command
and al,0F0h ; Minor opcode command
cmp al,0C0h ; MOV r32
je _call_mov
...
call _inc_ip
jmp _next_ip

```

Compiler Design

Can be as complex or as simple as you want. Don't forget if your VM is small you can always use Macros instead.

My First VM Compiler was in Perl (Back in 2004)

This one is in Ruby

- **Object Oriented Core**
- **Simple method for adding Opcodes**
- **Easily expandable**
- **Portable**

Very Simple VM Password Demo

; Sample MiniVM Code

POP r1 ; Get String off of stack

MOV r1,[r1] ; Get DWORD

CMP r1,0x34333231 ; Cmp to "1234"

JE GoodPassword

MOV r1,0 ; Set Stack to NULL

PUSH r1

JMP ExitCode

GoodPassword:

**MOV r1,1 ; Set Stack to 1 to show
; password was valid**

PUSH r1

ExitCode:

EXIT ; Quit VM

Ruby Compiler Object

MiniVMParser.rb:

```
class MiniVMParser < VMParser
```

```
...
```

```
  def define_opcodes
```

```
    ops = VMOpCodes.new
```

```
    ops.add("PUSH","r1",nil,"\x30")
```

```
    ops.add("PUSH","r2",nil,"\x31")
```

```
    ops.add("PUSH","r3",nil,"\x32")
```

```
    ops.add("CMP","r1","r2","\xd0")
```

```
    ops.add("CMP","r1",:value,"\xd8")
```

```
    ops.add("CMP","r2",:value,"\xd9")
```

```
    ops.add("MOV","eip",:value,"\xcc")
```

```
...etc...
```


Compiler Usage Output

miniVM Compiler

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Usage: minivmc [options]

Suggested Options:

-s, --source src	Source file to compile
-d, --destination dst	Destination file
-v, --verbose	Show opcodes per line
-o, --output style	Output style. [Bin, C, MASM]
-h, --help	Show this message

Compiler Directive XOR Example

```
; dbx directive example  
MOV r2, msg  
MOV r3, 76 ; r3 holds the xorkey  
JMP code  
  
msg:  
.dbx 76, '/etc/passwd',0  
code:  
MOV r1, r2  
MOV r1, [r1]  
AND r1, 0x000000FFh  
XOR r1, r3  
CMP r1, 0  
JE done  
  
...
```

Output of DBX directive

```
xxd -c 8 minivm.bin
```

```
00000000: c900 0000 0aca 0000 .....
```

```
00000008: 004c 6329 382f 633c .Lc)8/c<
```

```
00000010: 2d3f 3f3b 284c b042 -??;(L.B
```

```
00000018: 4000 0000 0002 d800 @.....
```

```
00000020: 0000 0020 0000 0028 ... ... (
```

XOR KEY

String

How to add your own directives

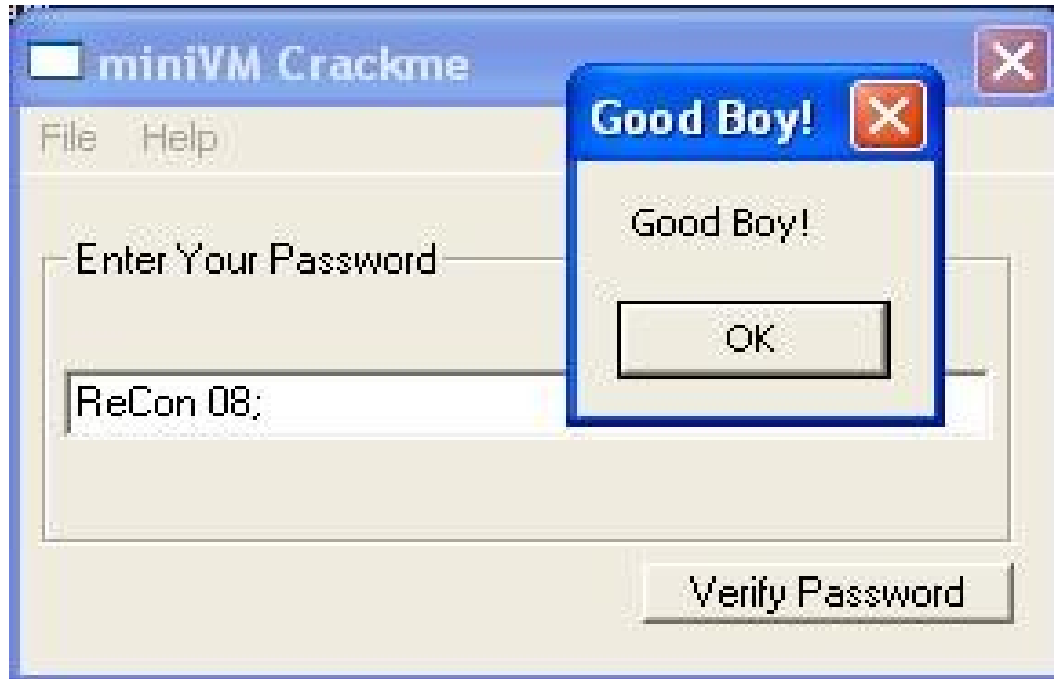
3 Steps:

- `@directives.add("mydirective")`
- `Def get_directive_size(tok)` (optional)
- `Def process_directive(tok, tokens)`

`tok.directive.cmd`

`tok.directive.line`

Our Crackme



Valid: **ReCon 08;**

Goal: To find more valid passwords

Tips for Debugging your VM

Debugging Techniques:

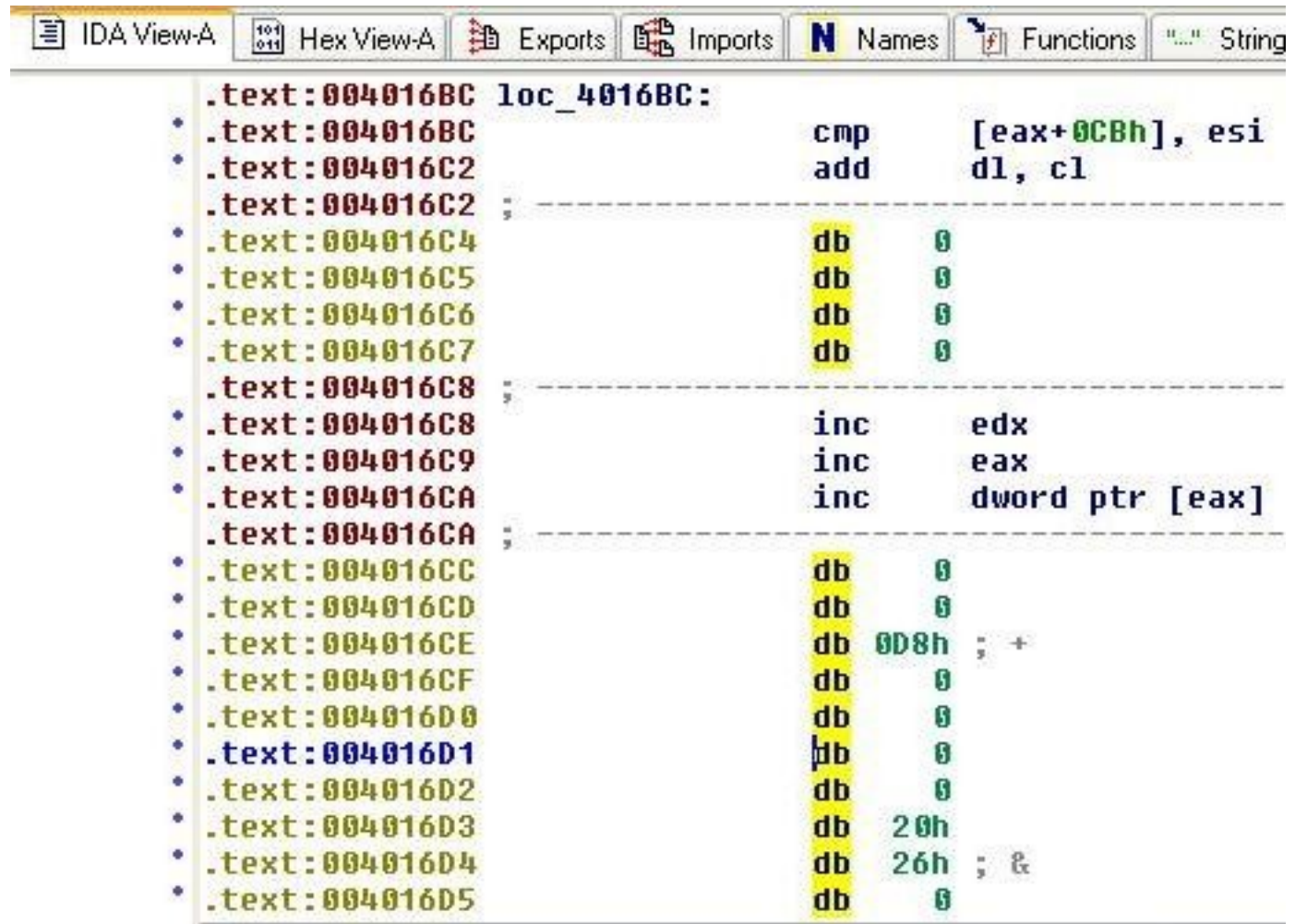
- Add INT 3 Breakpoints to your VM
- Break on the call handler table
- Minivmcs -v
- View your Virtual Registers while you are debugging

Attacks against your VM

- Your VM Core must be decrypted in order to process your p-code
- It is very simply to use a signature to identify a VM processor
- Use traditional methods to try and protect your VM core.

Remember this is just obfuscation, not security.
The goal is to quickly write code that takes a reverser much longer to analyze.

Example of our VM in IDA



The screenshot shows the IDA Pro interface with the assembly view of a function named `loc_4016BC`. The code is as follows:

```

.text:004016BC loc_4016BC:
* .text:004016BC          cmp     [eax+0CBh], esi
* .text:004016C2          add     dl, cl
; -----
* .text:004016C4          db     0
* .text:004016C5          db     0
* .text:004016C6          db     0
* .text:004016C7          db     0
; -----
* .text:004016C8          inc     edx
* .text:004016C9          inc     eax
* .text:004016CA          inc     dword ptr [eax]
; -----
* .text:004016CC          db     0
* .text:004016CD          db     0
* .text:004016CE          db     0D8h ; +
* .text:004016CF          db     0
* .text:004016D0          db     0
* .text:004016D1          db     0
* .text:004016D2          db     0
* .text:004016D3          db     20h
* .text:004016D4          db     26h ; &
* .text:004016D5          db     0
  
```


Look Ma, I can mutate!

Self Modifying Code (SMC) example:

```
MOV r1, mutate
ADD r1, eip    ; Adjust for relative offset
MOV [r1], 0x21000000h ; 0x21h == JL opcode
MOV r1, 6
CMP r1, 5     ; 6 > 5
```

mutate:

```
JG fakecode   ; Appears to always goto fakecode
               ; After mutation becomes JL <some addr>
JMP realcode
```

Advanced Extensions

XOR Register coupled with .xorkey directive

All Register Operations first pass through the XOR register

Example:

MOV xorkey, 76 ; where xorkey is a register

MOV r1, 1 ; 1 becomes 0x4Dh

MOV xorkey, r1 ; xorkey becomes 1

.xorkey 1

MOV r1, 1 ; 1 becomes 0

* Note included in this version

Shifting Operands

Shifting Operands

Similar to the xorkey register but used on the operand as follows:

- The CPU can be “seeded” on init with a value
- This value is used when parsing any operand byte (Example XOR)
- The compiler **MUST** know what the seed value is so it can write the appropriate opcode. Example: `.seed 0x4c`
- This seed value can change mid program

* Note included in this version

Where to get the Code

Neohapsis Labs (Blog)

VM, and Compiler

<http://labs.neohapsis.com/>

Crackme is here:

<http://crackmes.de/>

Email: craig.smith (\) neohapsis.com

Questions?

