

Recent Shellcode Developments

spoonm

ReCon, 2005

Part I

Introduction

Who am I?

- ▶ spoonm
- ▶ Metasploit developer since 2003
- ▶ University student
- ▶ Independent security researcher

What is this talk about?

- ▶ Recent shellcode research
- ▶ Older but lesser known tricks
- ▶ New tricks and techniques

Overview of shellcode structure

[payload]

- ▶ The payload does stuff

Part II

Shellcode

Making a connection

- ▶ Connect back to attacker (reverse)
- ▶ Connect to victim (bind)
- ▶ Reuse existing connection (find)
- ▶ Findsock only worth talking about

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- ▶ Embed connection source port in shellcode
- ▶ Call getpeername in a loop
- ▶ Found the socket when matching source port

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 - ▶ Simple idea and implementation
 - ▶ Works well when it works
- ▶ Cons
 - ▶ Doesn't work through a proxy
 - ▶ Doesn't work through a NAT
 - ▶ You need to embed source port

Find tag findsock

- ▶ Metasploit moved to using this
- ▶ Embed tag in shellcode
- ▶ Call recv in a loop
- ▶ Use MSG_DONTWAIT on linux
- ▶ Found the socket when you find the tag

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 - ▶ Will work through a proxy or NAT
 - ▶ Still fairly simple and small
- ▶ Cons
 - ▶ Data may be left on some sockets
 - ▶ Can be timing sensitive
 - ▶ More complicated on windows (ioctlsocket, etc)

Part III

Building a Nop Sled

Nop sleds, what and why?

- ▶ Often (especially on unix) we don't know exactly where our payload is
- ▶ Nop sleds makes the target we are trying to hit bigger
- ▶ The bigger the nop sled, the better the brute force

Improved shellcode structure

[nop sled][payload]

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[nop sled][payload]

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- ▶ The payload does stuff

Multibyte Nop Sled Concept

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- ▶ Generates instructions 1 to 6 bytes long, and uses 0x66 prefix
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- ▶ Prepends to the sled 1 byte at a time
- ▶ Generates a random byte and checks against tables

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- ▶ Builds the sled from back to front
- ▶ Prepends to the sled 1 byte at a time
- ▶ Generates a random byte and checks against tables
 - ▶ Is the instruction length too long?
 - ▶ Is it a valid instruction?
 - ▶ Does it have any bad bytes?
 - ▶ Does it modify restricted registers?

Backwardz

```
bb b0 bf 2c b6 27 67 2F 4A 1b f9 -- shellcode
| | | | | | | | | | | ... stc
| | | | | | | | | | | ^ . sbb edi,ecx
| | | | | | | | | | | ..... dec edx
| | | | | | | | | | | ..... das
| | | | | | | | | | | ^ ..... a16 das
| | | | | | | | | | | ..... daa
| | | | | | | | | | | ^ ..... mov dh, 0x27
| | | | | | | | | | | ^ ..... sub al, 0xb6
| | | | | | | | | | | ^ ..... mov edi, 0x6727b62c
| | | | | | | | | | | ^ ..... mov al, 0xbf
| | | | | | | | | | | ^ ..... mov ebx, 0xb62cbfb0
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 - ▶ Very fast, simple, deterministic
 - ▶ Allows for different scoring systems, recursion...

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 - ▶ Fairly language independent, C version 100 lines
 - ▶ Very fast, simple, deterministic
 - ▶ Allows for different scoring systems, recursion...
 - ▶ Can't support multibyte opcodes, escape groups, etc
 - ▶ Tables are pretty large, about 124k

OptyNop2 Output

```
$ ./waka 1000 4 5 | ndisasm -u - | head -700 | tail -20
000003B6 05419F40D4      add eax,0xd4409f41
000003BB 711C            jno 0x3d9
000003BD 9B             wait
000003BE 2C98           sub al,0x98
000003C0 37             aaa
000003C1 24A8           and al,0xa8
000003C3 27             daa
000003C4 E00D           loopne 0x3d3
000003C6 6692           xchg ax,dx
000003C8 2F            das
000003C9 49            dec ecx
000003CA B34A           mov bl,0x4a
000003CC F5            cmc
000003CD BA4B257715     mov edx,0x1577254b
000003D2 700C           jo 0x3e0
000003D4 C0D6B0         rcl dh,0xb0
000003D7 A9FD469342     test eax,0x429346fd
000003DC 67BBB191B23D   a16 mov ebx,0x3db291b1
000003E2 1D9938FCB6     sbb eax,0xb6fc3899
000003E7 43            inc ebx
```

ADMmutate Distribution - 1

total: 6000

uniq: 52

	00	01	02	03	04	05	06	07	08	09	0a	0b	0c	0d	0e	0f
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
10	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
20	00	00	00	00	00	00	00	6e	00	00	00	00	00	00	00	76
30	00	00	00	00	00	00	00	87	00	00	00	00	00	00	00	6a
40	6b	72	6a	68	74	66	77	6f	6d	74	6c	77	70	74	58	72
50	6a	67	71	70	7b	74	76	7c	70	7c	6b	78	00	6e	56	64
60	71	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
70	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
80	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
90	00	89	6c	78	00	74	72	df	7a	79	00	56	82	00	76	77
a0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
b0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
c0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
d0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
e0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
f0	00	00	00	00	00	7c	00	00	71	7f	00	00	69	00	00	00

ADMmutate Distribution - 2

total: 6000

uniq: 52

	00	01	02	03	04	05	06	07	08	09	0a	0b	0c	0d	0e	0f
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
10	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
20	00	00	00	00	00	00	00	64	00	00	00	00	00	00	00	6f
30	00	00	00	00	00	00	00	78	00	00	00	00	00	00	00	74
40	7f	6b	6f	7b	79	72	75	73	76	58	6f	7a	6c	78	7a	7e
50	71	6d	65	75	7f	72	7b	72	71	77	6d	64	00	71	7c	64
60	73	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
70	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
80	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
90	00	6b	79	87	00	74	74	e8	6b	68	00	76	5b	00	6d	72
a0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
b0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
c0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
d0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
e0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
f0	00	00	00	00	00	75	00	00	57	6b	00	00	6f	00	00	00

OptyNop2 Distribution - 1

total: 6000

uniq: 141

	00	01	02	03	04	05	06	07	08	09	0a	0b	0c	0d	0e	0f
00	00	12	12	12	39	39	00	00	12	11	11	11	39	39	00	00
10	12	12	12	11	39	39	00	00	12	12	12	12	39	39	00	00
20	12	11	12	12	39	39	00	39	12	12	11	12	39	39	00	39
30	11	11	12	12	39	39	00	39	11	11	12	11	39	39	00	39
40	39	39	39	3a	00	00	39	39	39	39	39	39	00	00	39	3a
50	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
60	00	00	00	00	00	00	39	39	00	12	00	11	00	00	00	00
70	3a	39	39	39	39	39	39	39	39	39	39	39	3a	39	39	39
80	12	12	00	12	12	11	11	12	12	12	00	00	00	00	00	00
90	39	39	39	3a	00	00	39	39	39	39	00	39	00	00	00	39
a0	00	00	00	00	00	00	00	00	3a	39	00	00	00	00	00	00
b0	3a	39	39	39	39	3a	39	39	39	39	39	39	00	00	3a	39
c0	12	12	00	00	00	00	00	00	00	00	00	00	00	00	00	00
d0	12	12	12	11	39	39	39	00	00	00	00	00	00	00	00	00
e0	39	39	39	39	00	00	00	00	00	00	00	39	00	00	00	00
f0	00	00	00	00	00	39	11	11	3a	39	00	00	39	39	11	11

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	00	01	02	03	04	05	06	07	08	09	0a	0b	0c	0d	0e	0f
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10	11	11	11	11	39	39	00	00	11	12	11	11	39	39	00	00
20	12	12	12	12	39	3a	00	3a	12	11	12	12	39	39	00	39
30	11	12	12	11	39	3a	00	3a	12	12	12	12	39	39	00	39
40	39	3a	3a	39	00	00	39	39	39	39	39	3a	00	00	39	39
50	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
60	00	00	00	00	00	00	39	39	00	12	00	11	00	00	00	00
70	39	39	39	39	3a	39	39	39	39	39	39	39	39	3a	39	39
80	11	12	00	12	11	12	11	12	12	12	00	00	00	00	00	00
90	39	39	39	3a	00	00	39	3a	3a	3a	00	39	00	00	00	39
a0	00	00	00	00	00	00	00	00	39	39	00	00	00	00	00	00
b0	39	39	39	39	39	39	39	39	39	3a	39	39	00	00	39	39
c0	11	11	00	00	00	00	00	00	00	00	00	00	00	00	00	00
d0	12	12	11	11	39	39	3a	00	00	00	00	00	00	00	00	00
e0	3a	39	39	39	00	00	00	00	00	00	00	39	00	00	00	00
f0	00	00	00	00	00	39	11	12	39	39	00	00	39	39	10	10

ADMmutate and optyx-mutate Gzip'd

```
# ADMmutate
```

```
$ time ./nops 1000000 | gzip -v >/dev/null  
 27.3%  
real    0m0.241s
```

```
# optyx's interzone mutate
```

```
$ time ./driver nop 1000000 | gzip -v >/dev/null  
 29.7%  
real    0m0.467s
```

OptyNop2 Gzip'd

```
# C version, save ESP and EBP
```

```
$ time ./waka 1000000 4 5 | gzip -v >/dev/null  
12.2%  
real    0m11.900s
```

```
# save just ESP
```

```
$ time ./waka 1000000 4 | gzip -v >/dev/null  
11.7%  
real    0m11.277s
```

```
# save nothing (good way to crash process)
```

```
$ time ./waka 1000000 | gzip -v >/dev/null  
8.3%  
real    0m12.404s
```

Conclusion

- ▶ Benefits
 - ▶ Handles restricted bytes and registers
 - ▶ More versatile sled generation (nop stuffing, etc)
 - ▶ Implementation and theory are simple

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- ▶ Possible Improvements
 - ▶ Support processor flags (nop stuffing)
 - ▶ Support 2-byte opcodes and escape groups
 - ▶ Improved byte scoring systems and look-ahead
 - ▶ Output according to a given byte distribution
 - ▶ Reduce the table sizes, memory usage

Part IV

Encoders

Encoder, what and why?

- ▶ We need to avoid bad characters
- ▶ Now we don't need to worry about this in the payload

Encoder, what and why?

- ▶ We need to avoid bad characters
- ▶ Now we don't need to worry about this in the payload

- ▶ Encodes the payload to a different set of characters
- ▶ Common methods: byte/word/dword xor, add, etc
- ▶ Prepends a decoder before the decoded data
- ▶ Decoder loops and decodes the payload

Improved shellcode structure

[nop sled][encoder (payload)]

- ▶ The nop sled slides into the decoder
- ▶ The decoder decodes the payload
- ▶ The payload does stuff

The simplest, call, 6 bytes

```
00000000  E800000000      call 0x5
00000005  58              pop  eax
```

The simplest, call, 6 bytes

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```

- ▶ Call pushes EIP of the pop instruction on the stack
- ▶ pop puts it in a register

jmp / call, 8 bytes

- ▶ You often need to avoid having 0x00

jmp / call, 8 bytes

- ▶ You often need to avoid having 0x00

```
00000000 EB02          jmp short 0x4
00000002 58             pop eax
00000003 90             nop
00000004 E8F9FFFFFF    call 0x2
```

jmp / call, 8 bytes

- ▶ You often need to avoid having 0x00

```
00000000 EB02          jmp short 0x4
00000002 58           pop eax
00000003 90           nop
00000004 E8F9FFFFFF   call 0x2
```

- ▶ Jmp to a call instruction
- ▶ The call is now backwards (negative)

FPU Get EIP, 7 bytes

- ▶ Sometimes you want to avoid 0xff
- ▶ Noir's Get EIP

FPU Get EIP, 7 bytes

- ▶ Sometimes you want to avoid 0xff
- ▶ Noir's Get EIP

```
00000000  D9EE                fldz
00000002  D97424F4           fnstenv [esp-0xc]
00000006  58                pop eax
```


FPU Get EIP, 7 bytes

- ▶ Sometimes you want to avoid 0xff
- ▶ Noir's Get EIP

```
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00000002  D97424F4           fnstenv [esp-0xc]
00000006  58                pop eax
```

- ▶ fnsetenv will get EIP of last fpu instruction
- ▶ It allows for much more permutations
- ▶ Also smaller than jmp/call

Call \$+4, 7 bytes

- ▶ The coolest of them all
- ▶ Gera & CoreST

Call \$+4, 7 bytes

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00000000	E8FFFFFFF	call 0x4
00000005	C3	ret
00000006	58	pop eax
00000004	FFC3	inc ebx
00000006	58	pop eax

Call \$+4, 7 bytes

- ▶ The coolest of them all
- ▶ Gera & CoreST

00000000	E8FFFFFFF	call 0x4
00000005	C3	ret
00000006	58	pop eax
00000004	FFC3	inc ebx
00000006	58	pop eax

- ▶ Call is relative to the end of it's instruction
- ▶ Call jmps into itself, decodes an 0xff instruction
- ▶ can inc, dec, or push reg

SEH, bigger

- ▶ People really want a alpha numeric get eip

SEH, bigger

- ▶ People really want a alpha numeric get eip
- ▶ Well, so far, only found a way for windows
- ▶ Link in exception handler, cause exception
- ▶ Get EIP from the CONTEXT record

CLET

- ▶ Generates permutations of decoder stubs
- ▶ Inserts reversing instructions, nop equivalents
- ▶ All decoders are C code to generate themselves

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- ▶ Generates permutations of decoder stubs
- ▶ Inserts reversing instructions, nop equivalents
- ▶ All decoders are C code to generate themselves
- ▶ Pros:
 - ▶ Well thought out - analyzed attacks against NIDS
 - ▶ Mathematica files output, mathy backing
 - ▶ Specturm analysis - push sled to byte distribution

CLET

- ▶ Generates permutations of decoder stubs
- ▶ Inserts reversing instructions, nop equivalents
- ▶ All decoders are C code to generate themselves
- ▶ Pros:
 - ▶ Well thought out - analyzed attacks against NIDS
 - ▶ Mathematica files output, mathy backing
 - ▶ Specturm analysis - push sled to byte distribution
- ▶ Cons:
 - ▶ Complicated system, really hard to build upon
 - ▶ Decoder generation isn't that great
 - ▶ Making compromises for size/robustness

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- ▶ "Conservative Polymorphism"
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- ▶ Pros:
 - ▶ Polymorphizing code is pretty easy
 - ▶ No size or functionality compromises
 - ▶ Bad character and register avoidance
- ▶ Cons:
 - ▶ Less thought out, NIDS attacks not deeply analyzed
 - ▶ Hard to push to arbitrary byte distribution
 - ▶ Less "polymorphism", more restrictions

Implementation - Pex::Poly

- ▶ "Blocks" are dependency graph nodes
- ▶ "Blocks" consist of 0 or more possibilities
- ▶ Register pool assignment (mov reg1, reg2)
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 - ▶ Pex::Poly has 3 phases
 - ▶ Dependency iteration and block selection
 - ▶ Instruction offset calculations
 - ▶ Instruction register assignment

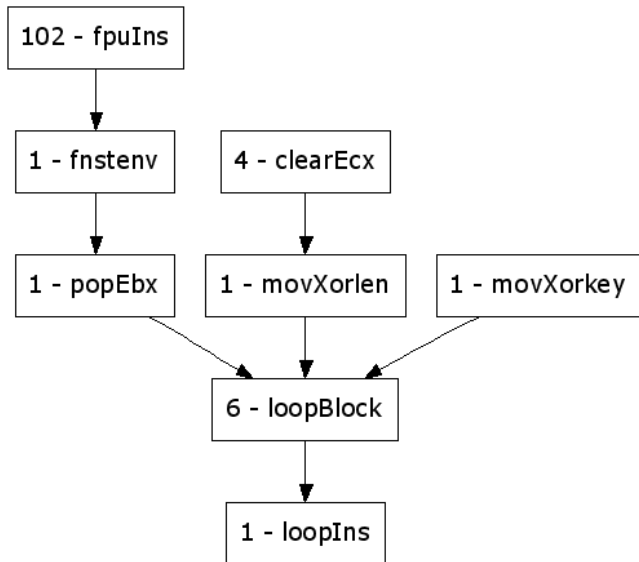
Shikata Ga Nai

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- ▶ Created one decent "polymorphic" encoder

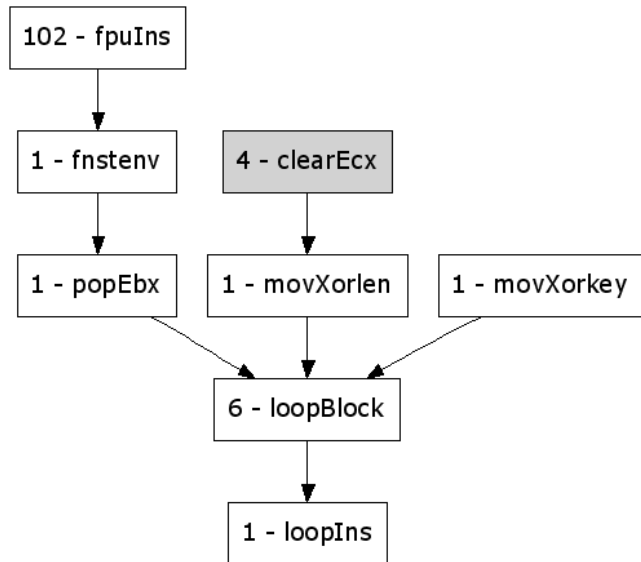
Shikata Ga Nai

- ▶ It's too much work to polyize every payload
- ▶ Created one decent "polymorphic" encoder
- ▶ Noir's FPU geteip technique
- ▶ Approximately 1.3 million permutations
- ▶ Additive feedback xor, encodes it's own end
- ▶ 27 bytes for the stub, 4 key, 4 encoded

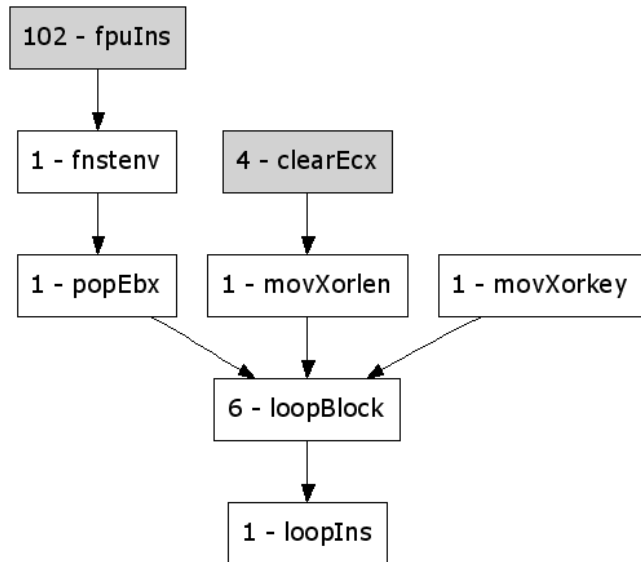
Shikata dependency iteration



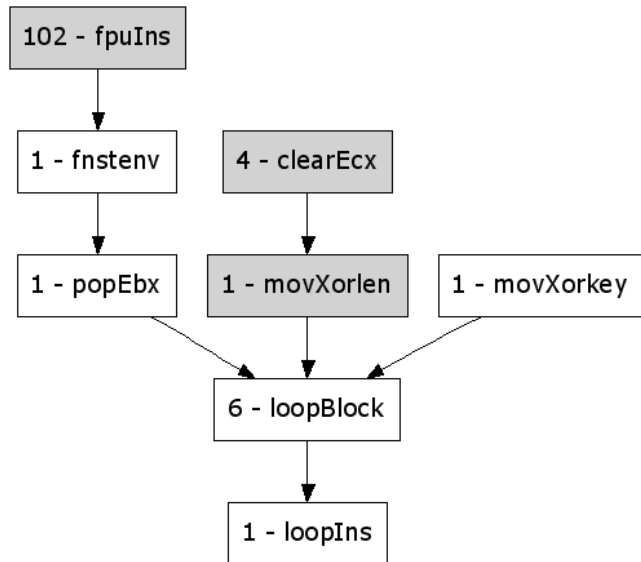
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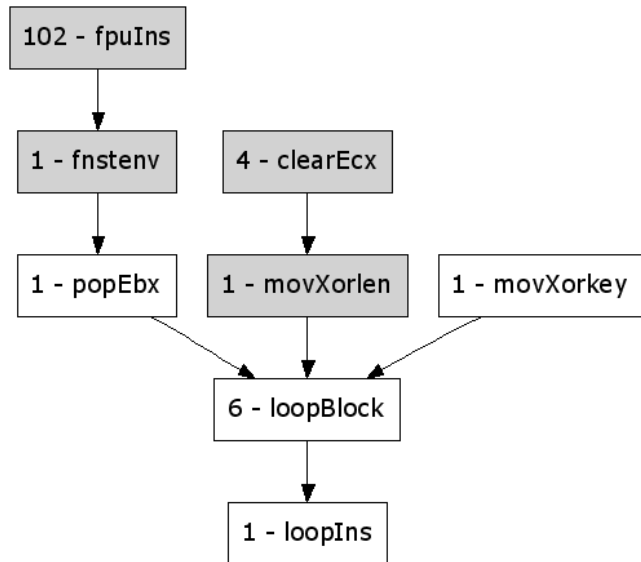
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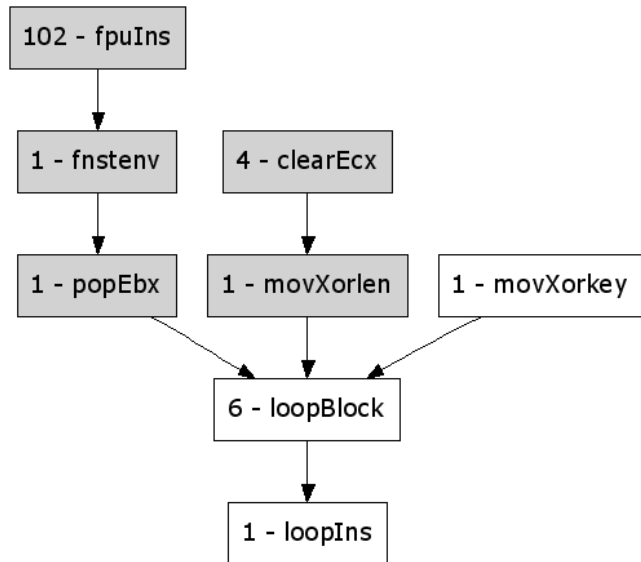
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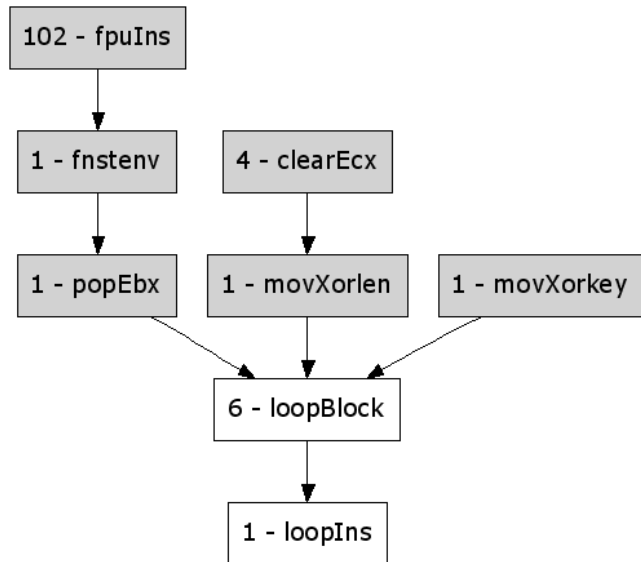
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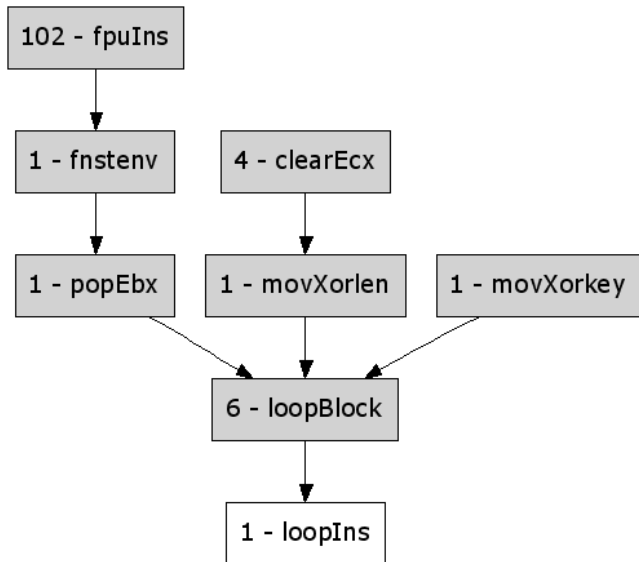
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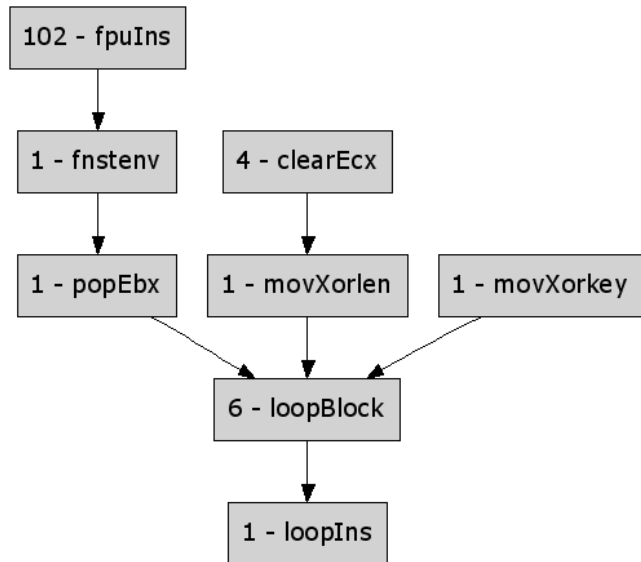
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Shikata dependency iteration



Example output

00000000	BB6E887A69	mov ebx,0x697a886e
00000005	DDC4	ffree st4
00000007	D97424F4	fnstenv [esp-0xc]
0000000B	58	pop eax
0000000C	29C9	sub ecx,ecx
0000000E	B101	mov cl,0x1
00000010	83E8FC	sub eax,byte -0x4
00000013	31580E	xor [eax+0xe],ebx
00000016	03580E	add ebx,[eax+0xe]
00000019	E2F5	loop 0x10

Example output

00000000	DBC1	fcmovnb st1
00000002	31C9	xor ecx,ecx
00000004	B101	mov cl,0x1
00000006	D97424F4	fnstenv [esp-0xc]
0000000A	5B	pop ebx
0000000B	BAC8E2C8F8	mov edx,0xf8c8e2c8
00000010	83C304	add ebx,byte +0x4
00000013	315313	xor [ebx+0x13],edx
00000016	035313	add edx,[ebx+0x13]
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Example output

00000000	BB7B833BB9	mov ebx,0xb93b837b
00000005	DAC0	fcmovb st0
00000007	D97424F4	fntenv [esp-0xc]
0000000B	2BC9	sub ecx,ecx
0000000D	5E	pop esi
0000000E	B101	mov cl,0x1
00000010	315E12	xor [esi+0x12],ebx
00000013	83C604	add esi,byte +0x4
00000016	03	db 0x03
00000017	25	db 0x25
00000018	8D	db 0x8D
00000019	D9	db 0xD9
0000001A	4C	dec esp

Part V

Egg Hunters

Egg Hunting, what and why?

- ▶ Sometimes we have very small size constraints
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- ▶ We need to validate a memory region before we search it

Improved shellcode structure

[nop sled][egg hunter] - [encoder (payload)]

- ▶ The nop sled slides into the decoder
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 - ▶ Windows system calls are not meant to be static
 - ▶ Parse and build tables of all windows syscalls
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- ▶ Windows system calls are not meant to be static
- ▶ Parse and build tables of all windows syscalls
- ▶ Find a static system call, or collection that will work
- ▶ `NtAccessCheckAndAuditAlarm` works
- ▶ Need special privileges for to work normally
- ▶ It validates the pointer before it validates rights

Windows egghunt example

```
// Skape's syscall egghunter

                                // Address to check in edx
push  0x2                        // Push NtAccessCheckAndAuditAlarm
pop   eax                        // Pop into eax
int   0x2e                       // Perform the syscall
cmp   al, 0x05                   // Did we get 0xc0000005 (ACCESS_VIOLATION)
```

Part VI

Staging

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- ▶ We often have size constraints
- ▶ Staging abstracts the connection mechanism from the payload
- ▶ A stager establishes a connection to the attacker
- ▶ The stager reads in more code from the connection
- ▶ The stager executes the stage passing the connection

Improved shellcode structure

[nop sled][encoder (stager)] - [stage]

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Size issues

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 - ▶ As you saw, it would be hard to use syscalls
 - ▶ We need to use the Windows APIs (ws2_32.dll...)
 - ▶ But function resolving takes a ton of code!
 - ▶ What in the world can we do!?

ws2_32.dll static ordinals

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- ▶ Find ws2_32.dll base
- ▶ Resolve our functions by static ordinals
- ▶ 93 byte reverse connect shellcodez y0!

Part VII

Questions?