



SHALL HE PLAY A GAME?

Lessons learned while playing CoreWars8086

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#Whois Elad Shapira ('Zest')



- Reverser from the Holy Land.
- Mobile Security Researcher @AVG.
- Highly passionate for RE, Assembly and Low-Level.
- Speaker (ClubHack, Ground Zero Summit..).
- Co-Organizer of CoreWars8086 competition (IL).















A total of 300 tickets will be sold for this year's conference.





Agenda

- Timeline of the CoreWars8086 competition.
- Arena, Engines and rules.
- How to analyze and write survivors.
- Optimization.
- Anti reversing techniques.
- Future / Improvements.
- Share ideas \rightarrow Create new ideas!
- Hangover.



Origin

- Alexander Dewdney / D.G. Jones.
- CoreWars / RedCode
- <u>http://vyznev.net/corewar/guide.html</u>

[Note: This is a reproduction of the Core War Guidelines originally produced by Jones and Dewdney in March of 1984] Red's dead baby. Red's dead.

CORE WAR GUIDELINES

D. G. Jones and A. K. Dewdney

Department of Computer Science The University of Western Ontario

March, 1984



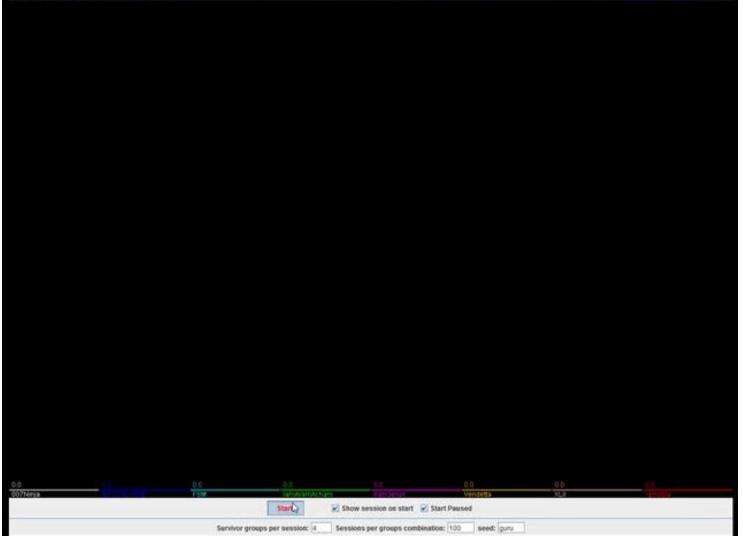








CodeGuru Extreme - Competition Viewer





Cameras usually add 5 kg ..



- Top 4 survivors get to the final.
- Final \rightarrow Winners!





https://code.google.com/p/corewars8086/

corewars8086

CoreWars 8086 game engine, written in Java.

Survivors in general

- Download, Unzip & play (Google Code).
- Survivor's name == file's name (without extension).
- 8086 opcodes, 16bit instructions.
- Not all instructions are supported (Pusha, Popa,..).
- Compiled as `com' file
 - DOS command file format.
- Maximal survivor size **512** bytes.
- Each team can submit two survivors.
 - Rocky1 & Rocky2.







Virtual Arena

- Loaded to the virtual arena each time with random address (copied "as is").
- Distance between two survivors and the sides is at least 1024 bytes.
- All cells initialized to 'CCh' before start.
- End of the battle
 - 200,000 rounds or one survivor left.
- Order of the survivors is determined randomly at the beginning and cannot be changed.



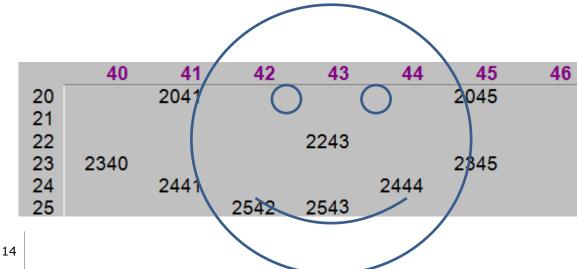
Arena (NOT virtual)





Arena & Addresses

	00	01			FE	FF
00	0000	0001	••	••	00FE	00FF
01	0100	0101	••	••	01FE	01FF
:	:	:			:	:
:	:	:			:	:
FE	FE00	FF01			FEFE	FEFF
FF	FF00	FF01			FFFE	FFFF



mov [2041h], al mov [2045h], al mov [2243h], al mov [2340h], al mov [2441h], al mov [2542h], ax mov [2444h], al mov [2345h], al



Survivor's Registers (before 1st round)

- **BX,CX,DX,SI,DI,BP** = 00s.
- **Flags** = 00s.
- AX, IP Initial location of the survivor, offset.
- CS, DS Segment that was assigned to the survivors.
- ES Segment for survivors from same team (shared memory) – 2048 bytes.
- **SS** Beginning of the personal stack (2048).
 - ss:0x00 ss:0x7ff, initialized to 0x00.
- SP Offset of beginning of personal stack (00s).





How survivor gets killed

- Running illegal command
 - The 060h byte does not translated to an assembly command.
 - Engine: "Died due to CPU".
- Running commands that are not supported by the engine
 - For example `int 21h'.
- Access to memory not in the arena or not in the range of the survivor's personal stack.
 - For example ES:0x1234.
 - Engine: "Died to memory exception".

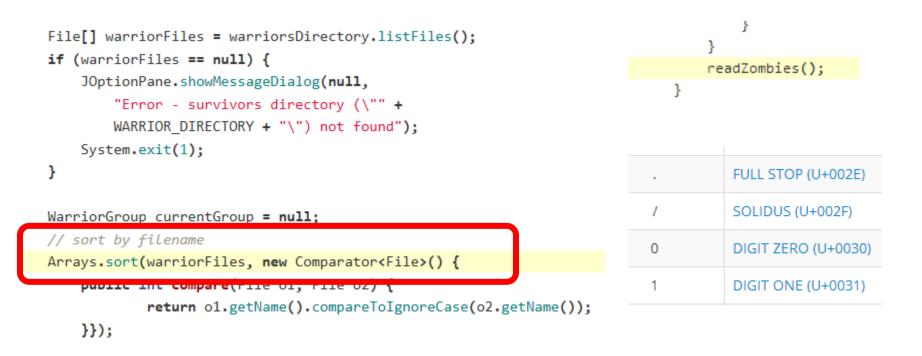


Zombies

- Sent by organizers before competition begins.
- Regular survivors that do not get points.
- Different CPU states problem.
 - Direction flag (MOVSW will kill master).
- Zombies can still win the battle
 - less points for us.
 - We need to encourage them to commit suicide.
- Contain Math Riddles (That you need to solve).



Pwning bugs in the engine



How to make your survivors be the firsts to run? OSurvivorName What is the advantage?

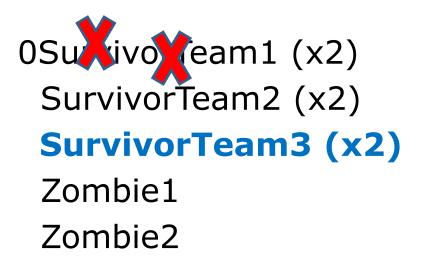


Zombies can fix your survivor's code

0Survivor eam1 (x2) Survivor Team2 (x2) Survivor Team3 (x2) Zombie1 Zombie2

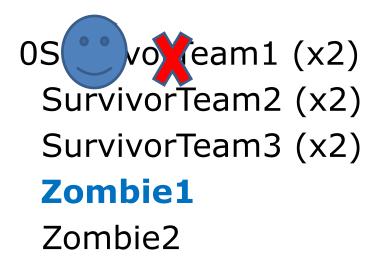


Zombies can fix your survivors code



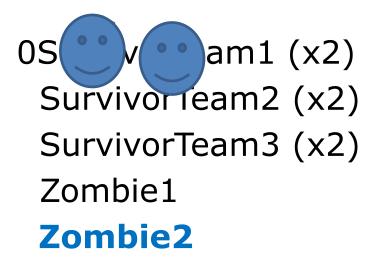


Zombies can fix your survivors code



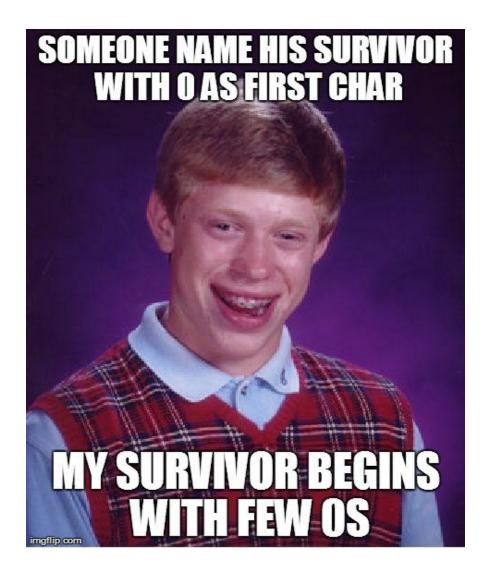


Zombies can fix your survivors code



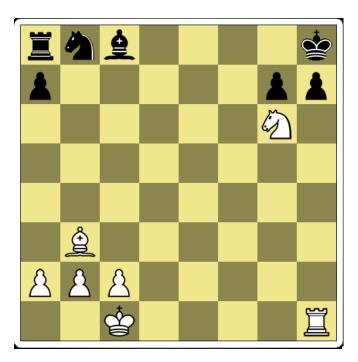


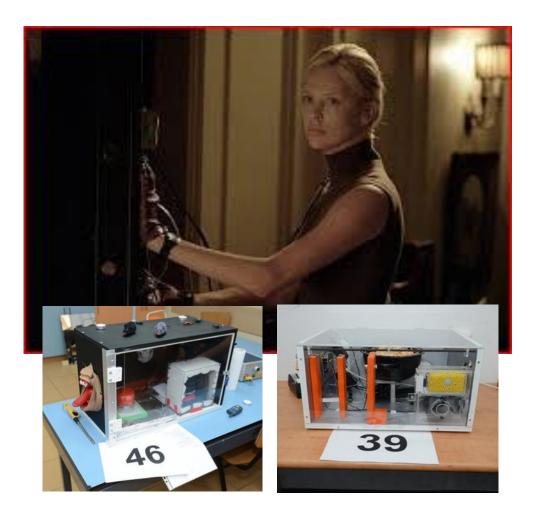
To stay on the safe side..





Safe Cracking







Safe example#1

loop: mov AX,[1234] [1234] = AAAB 3*AX=1 mov BX,3 mul AX BX*AX=1 sub AX,1 **AX=1** loop jnz **ZF=1** killer:

		-2	1845
1010 15	1010	1010	1011 0
			AAAB

1010 15	1010	1010	1011 0

killer: Solution: mov AX, AAAB mov ptr word [1234], AX JMP killer



Safe example#2

```
loop:
mov AL,[111] [111] = 49H
add AL,0A8h 73+168=241(F1)
mov AH, [112] [112] = 42H
xor AH,0ADh ADH xor 42H = EFH (239d)
mul AH
               AX = AH * AL = 239 * 241 = 57599
cmp AX,0xeOff AX=57599d
               ZF=1
    loop
ine
```





Safe example#2

loop: mov AL,[111] add AL,0A8h mov AH, [112] xor AH,0ADh mul AH cmp AX,0xeOff loop ine



Solution: killer: mov AL, 49H mov AH, 42H mov ptr byte [111], AL mov ptr byte [112], AH killer jmp



Important factors

- Survivors usually contain
 - Initialization.
 - Bombing loop.
 - Write -> Update address for next writing -> Jumping to beginning of loop
- We usually measure survivors by
 - 'Area of vulnerability'
 - 'Attack rate'.
- We can cause unexpected phenomenon
 - mov AX, 0000 -> mov ax, 0cccch (2,3 bytes).



Looper

Smallest functional survivor (EBFE, jmp \$): Loop: Jmp loop

Good to test other survivors.

Attack	Vulnerability		
sequence	profile		
3/1	5		

Bomber Demo

AVG

CodeGuru Extreme - Session Viewer			-	
<u></u>			i-	Bomber nothinga nothingb
Bomber enters the arena. nothing enters the arena. nothing enters the arena. nothingb enters the arena.				
Round:	ise Speed:	O View CPU	Resultye Single	Round

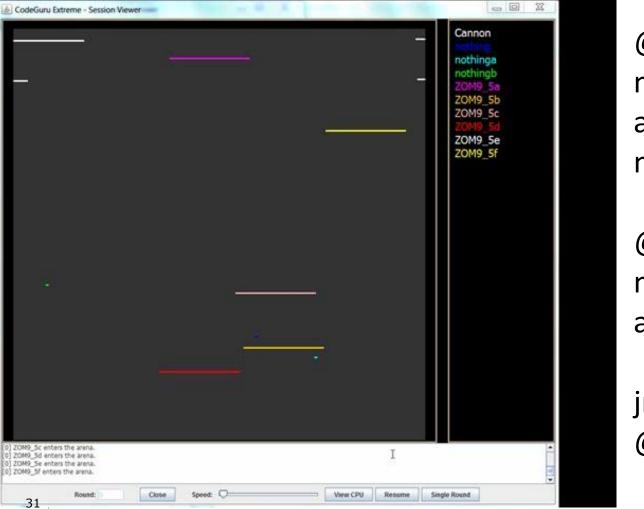
mov al, 0CCh mov bx, 0 @loop: mov [bx], al inc bx jmp @loop



Attack	Vulnerability
sequence	profile
3/1	7



Cannon Demo



@start: mov bx, ax add bx, (@end - @start) mov al, 0CCh

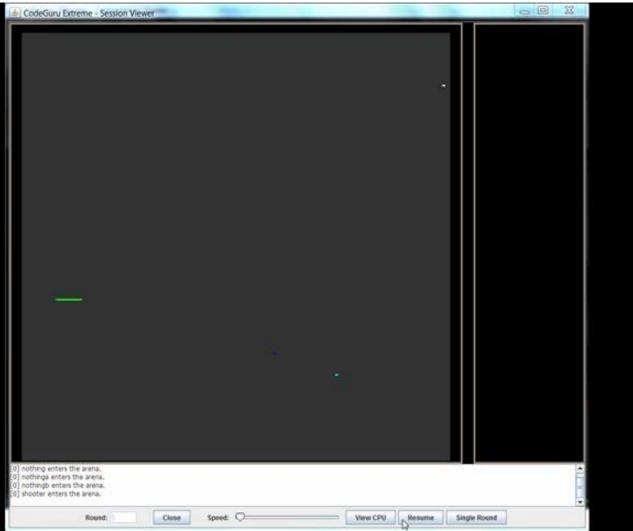
@loop: mov [bx], al add bx, 8

jmp @loop @end:



Attack	Vulnerability
sequence	profile
3 / 2	6

Shooter Demo



MOV DI,AX MOV AX,0CCCCh

@loop:STOSWADD DI,9JMP @loop



Heavy Bombing



dh

di

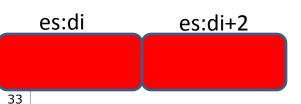
0000

dl

es

0000

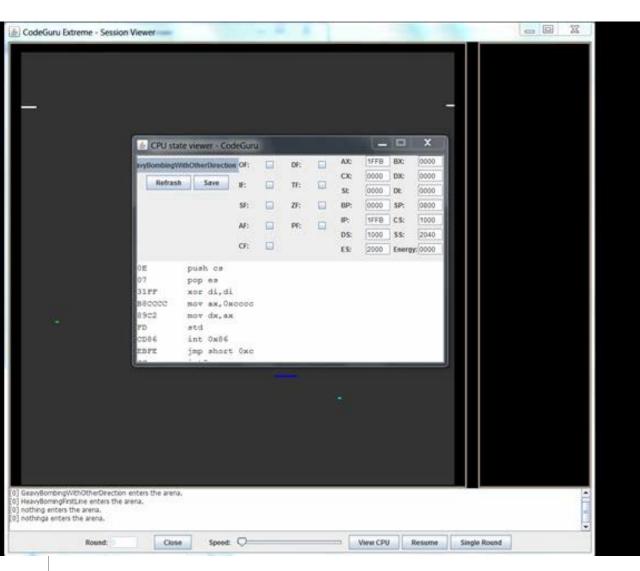
- Writes on 256 bytes (es:di -> 255 addresses)
- es same value as cs -> if not memory exception after the interrupt
- CLD/STD -> change direction
- 2 Heavy Bombing each battle
- We can bomb shared segment
- INT 86h







Heavy Bombing Demo (Opposite direction)

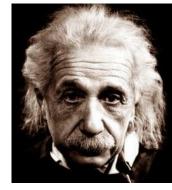


push cs pop es xor di,di mov ax, Occcch mov dx, ax std int 86h jmp \$



35

Smart Bombing



- Bombing the first occurrence of AX:DX in memory.
- Replacing it with data we want
 - Illegal commands or jmp to our code.
- We can attack ourselves...
- 1 Smart Bombing each battle.
- INT 87





Protection from Smart Bombing

- Change functionality of registers (BX <-> BP).
 - Usually does not matter.
- Change order of independent commands
 - Put 3 values to 3 registers = Few different ways.
- copy parts of the code
 - To the beginning and the end.
- Variable that changed during runtime near main loop/code part (SP).
- Encoding with random numbers.
- **XORing** (will be discussed later).



Smart bombing FAIL protection (CGX#9.5)

jmp short 0x12 mov si,0x95a0 xchg ax,bx cld lodsw std cmp ax,bx jnc Oxc or al,0x90 lodsw loop 0x6 E2F4 mov si,0x95a0 **BEAO** xchg ax,bx Cld lodsw

std cmp ax,bx jnc 0x1c or al,0x90 lodsw loop 0x16 mov si,0x95a0 xchg ax,bx cld lodsw std cmp ax,bx jnc 0x2c or al,0x90 lodsw loop 0x26

push cs pop es mov ax, 0F4E2h mov dx, 0A0BEh mov cx, 0cccch mov bx,cx STD Int 87h Jmp \$ Zombie ==?



Binary search ("Lion in the desert")

Jumping to body jmp short 0×12 The "talking location" that the . . mov si,0x95a0 survivors and the zombie talk in xchg ax,bx Keep loading address on the side cld • lodsw (LODSW will change AX) std Clears the direction flag (DF=0) cmp axbx jnc 0x1c LODSW === MOV AX,[SI++ or SI--] or al,0×90 lodsw AX will hold the 'talking location' loop 0×16 DF=1 (later SUB SI, 2 to change back) 38



Binary search ("Lion in the desert")

jmp short 0×12 Compare his address (BX) to talking location (AX) - change only flags. . . mov si,0x95a0 AX >= BXxchg ax,bx cld 👉 jumps into itself (IP increased by 1) lodsw std Dec [si] 🗲 Next cell 73 FF 73 **FF** cmp ax,bx OC 90 **OC** 90 nop jnc 0x1c changes AL + AX changed again? or al,0×90 lodsw hidden Dec[Si] command 😳 loop 0×16 DF=1 (sub si, 2 to change back)



6 Zombies

push cs pop es int 0x87 and ax,0x7fff push ax mov bl,[0xc0de] test bl,bl jns 0x16 div bl mov [0xc0dd],ah pop ax jmp short 0x7

Zombie ==?

mov bl,[0xc0de] mov bl,[0xc1de] mov bl,[0xc2de] mov bl,[0xc3de] mov bl,[0xc4de]

mov [0xc0dd],ah mov [0xc1dd],ah mov [0xc2dd],ah mov [0xc3dd],ah mov [0xc4dd],ah

public start start proc near push cs pop es assume es:seg000 int 87h and ax, 7FFFh
ý
<pre>loc_10107: push ax mov bl, ds:0C0DEh test bl, bl jns short loc_10116</pre>
div bl mov ds:0C0DDh, ah
loc_10116: pop ax jmp shortloc_10107 start endp seg000 ends
end start



Chinese Remainder Theorem

$$ZX \equiv 1 \pmod{3}$$
 $X \equiv 2 \pmod{3}$
 $3x \equiv 2 \pmod{4} \iff X \equiv 2 \pmod{4}$
 $4x \equiv 3 \pmod{5}$ $X \equiv 2 \pmod{5}$
 $X \equiv 2 \pmod{5}$

Formula used to find all the zombies: input = ? a1 = (input%254); a2 = (input%255); input = (a1*255*1 + a2*254*254)%(255*254);





Sometime, the organizers send invalid zombies...





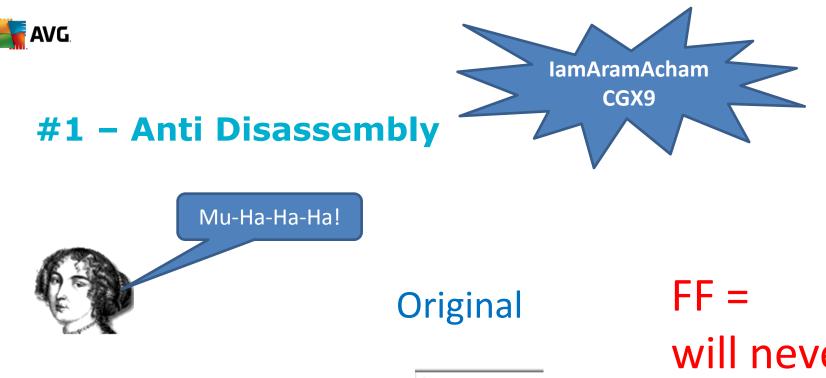


	#	bytes	opcodes	rounds	Not optimized	Optimized	bytes	opcodes	rounds	
	1	3	0x83	1	add ax,1	inc ax	1	0x40	1	
			0xc0							
			0x01							
	2	3	0x83	1	sub ax,1	dec ax	1	0x48	1	
			0xe8							
-	2	3	0x01	1	Maria AV O	Cub AV AV	2	0x29	1	
	3	3	0xb8 0x00	1	Mov AX,0	Sub AX,AX	2	0x29 0xc0	1	
			0x00					UXCU		SHELLCODE
	4	5	0xb9	2	mov CX,2	shr ax,1	3	0xc1	1	Z X2F Last Last Last X15 5 (XCE\X05\X3A\X7E\) 2
			0x02					0xe8		S SAOTART ROCARS J
			0x00					0x02		
			0xf7		div CX					
			0xf1							
	5	6	0x89	3	mov dx,ax	XCHG ax,bx	1	0x93	1	
			0xc2							
			0x89		mov ax,bx					
			0xd8		and the design of the second sec					
			0x89 0xd3		mov bx,dx					
	6	7	0xu3 0xa2	2	mov [0],al	mov [0],ax	3	0xa3	1	
	Ŭ		0x00	-		inev [e],ux	, j	0x00	-	
			0x00					0x00	Bit T	widdling Hacks
			0x88		mov[1],ah					in a starting fraction
			0x26						By Se	an Eron Anderson
	12		0x01							er@cs.stanford.edu
2	43		0x00							



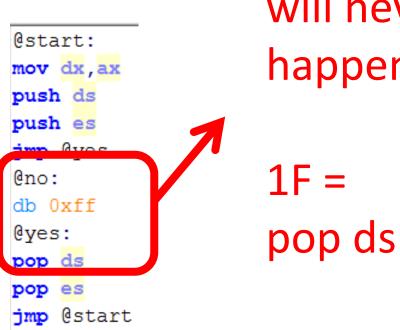
How not to be seen





Disassembly

00000000 00000002 00000003	89C2 1E 06	mov dx,ax push ds push es	
0000000 0000000 0000000	EB01 FF1F	mp short 0x7 all word far op es	[bx]
	EBF5	jmp short 0x0	



FF = will never happen



#2 – Usage of unsupported registers (1/2)

- FS is unsupported by engine.
- difference between opcode interpretation between 8086 and later processors like 80386
- 8086 processor will read it like 'ES'.

C:\Program Files (x86)\nasm>ndisasm.exe CodeToCompile 00000000 8EE3 mov fs,bx

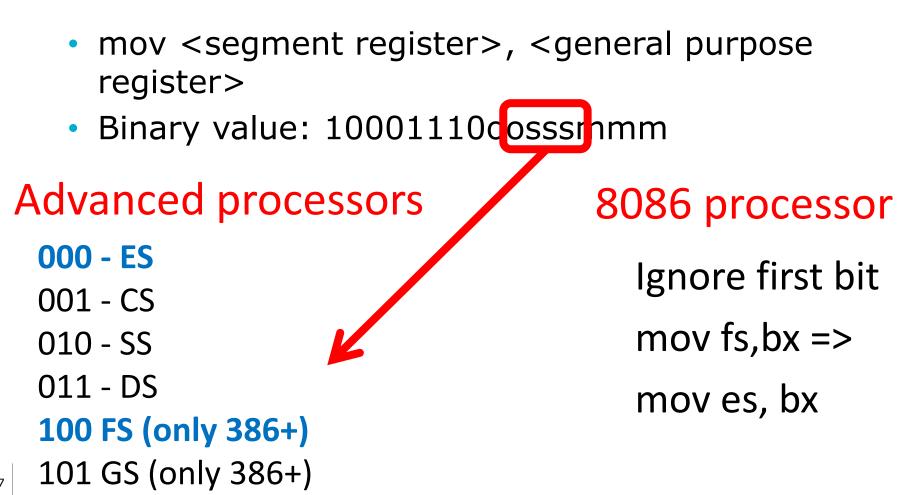
push es, ds move bx, ds mov fs, bx

HutsHuts

CGX3



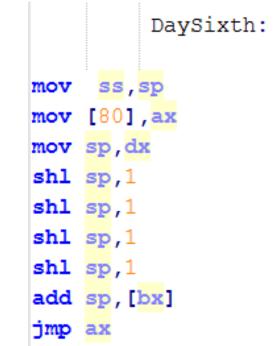
#2 – Usage of unsupported registers (2/2)







- Targeting flaws that can be found in debuggers.
 - Example: debug.exe.
- SP (Stack Pointer) gets really small value like '4' -> debugger crush.





#4 - Random bits



- Write multiple INT3 ('CC', unsupported opcode) in places that are not part of the code flow.
- After compilation replace all occurrences of 'CC' to random bits (Hex Editor/script).
- For example F1, D6 etc.

case (byte)0xCC: // INT3
 throw new IntOpcodeException();

```
mov es, sp
int 87h
add si,cx
call far [si]
                     O}-<Q<=10018 MTWThF
                   02.28.08 7:55pm |* *| Q<=
int 3
                   O}-< 1/2?...1/256? omg! <3
int 3
                   lol!{u}()3=<">=H207xNaCl
int 3
                   yllambywlaw Q<= "?" O}-<
int 3
                   ..-.etc jv*\o/* #*@%! "--- ---"
                   aabb O}-<:(|)xXx })i({ SW .+
int 3
                   5th/42nd ° No.2 ---> ]?
int 3
int 3
```



#5 – XORing the code



- Taking survivor's body and generate two binary strings - XOR of them will be the body of the original survivor.
- During runtime every survivor copy his part to the shared memory and they calculate XOR of the two parts before it is run by the survivor.
- Also Smart bombing protection..

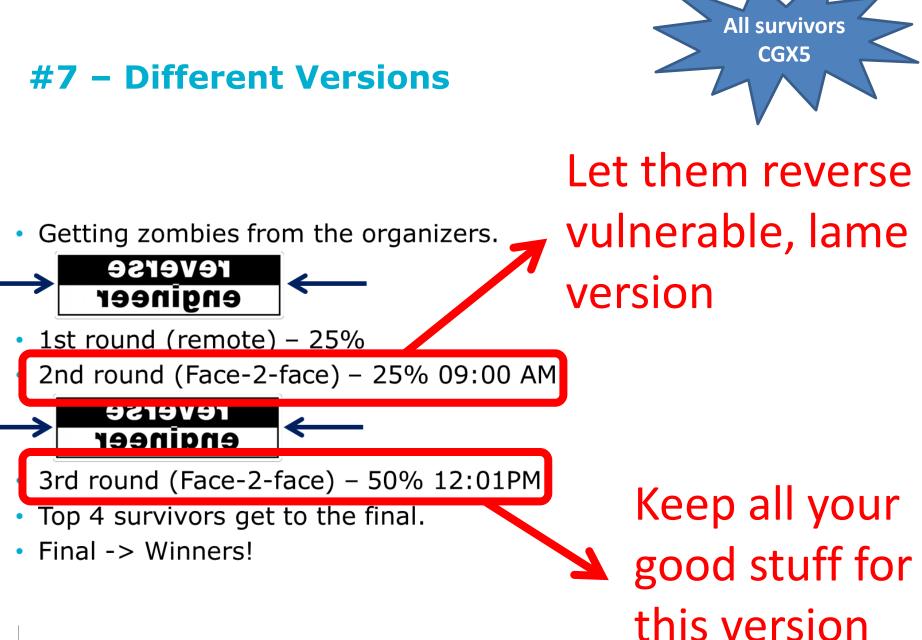




#6 – Copy of a zombie

- Copy zombie into our survivors so others will pwn a fake zombie instead of the real one.
- Cons: valuable space is wasted.







That's what happens to a team that achieves 1st place before the final round..





2 –ASM counts

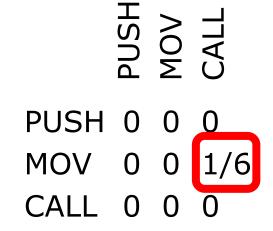
3 – Probabilities



1 – Generating ASM instruction trace

PUSH
 MOV



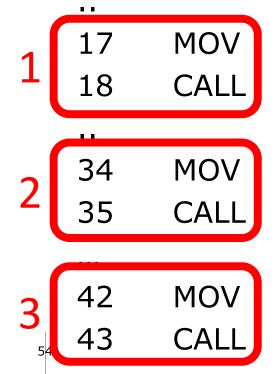


4 – Weighted directed graph for code



5 – Weighted directed graph for code

$$score(A, B) = \frac{1}{N^2} \left(\sum_{i,j=0}^{N-1} |a_{ij} - b_{ij}| \right)^2$$





Genetic Programming for Reverse Engineering

Mark Harman^{*}, William B. Langdon^{*} and Westley Weimer[†] ^{*}University College London, CREST centre, UK [†]University of Virginia, Virginia, USA

Genetic Programming



A lot of work was done on RedCode

 John Perry , Jason Boar, Ryan Colman, Wilkies Benchmark, Dave Hillis and others.

One effort was done on CoreWars8086.

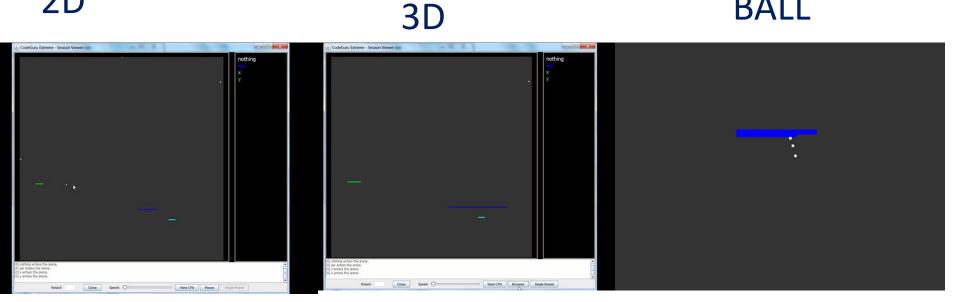
- Darwin8086.
- Gen = Warrior = String 1-512 bytes.
- Chromosome
 - Bit, Command, Meta-command, Combination,...
- Fitness function Endogeny, Exogeny.

Graphical Survivors (Make Love Not War)

2D

AVG.

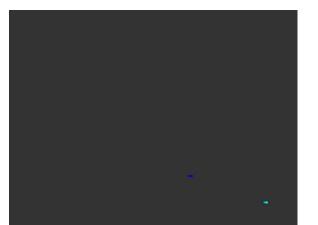
BALL



SIR









Graphical Survivors (Make Love Not War)

Come on you

Hexy Boy..

3D





AVG

2D





This is how can we add 'hardware hacking'..

#@&%*^@& !!!!!!! I knew I should stick to PHP !!!

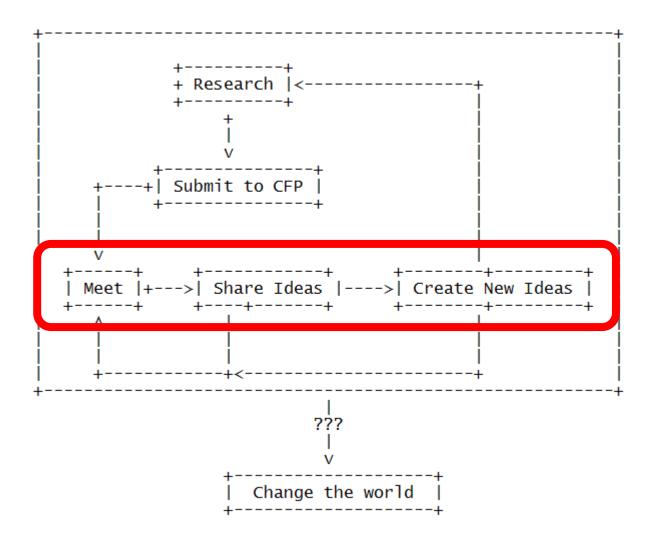




Scrum? Agile?



Q & A / Feedback

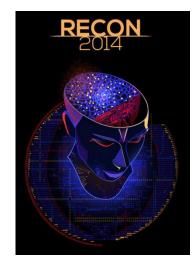


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Thank you! Merci!

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T TC TCP

ACK

- Hugo, Sam, Elizabeth and the ReCon team!!!
- Dr. Oded Margalit, Assaf Nativ, Ange Albertini, ShiftReduce, SonOfLilit, Danny Leshem, DualCore and Others..
- AVG, Oren Barad & The team.
- My (brave) Wife & kids.

300 Ninjas & Reversers..