

Hardening Registration Number Protection Schemes against Reverse Code Engineering with Multithreaded Petri Nets

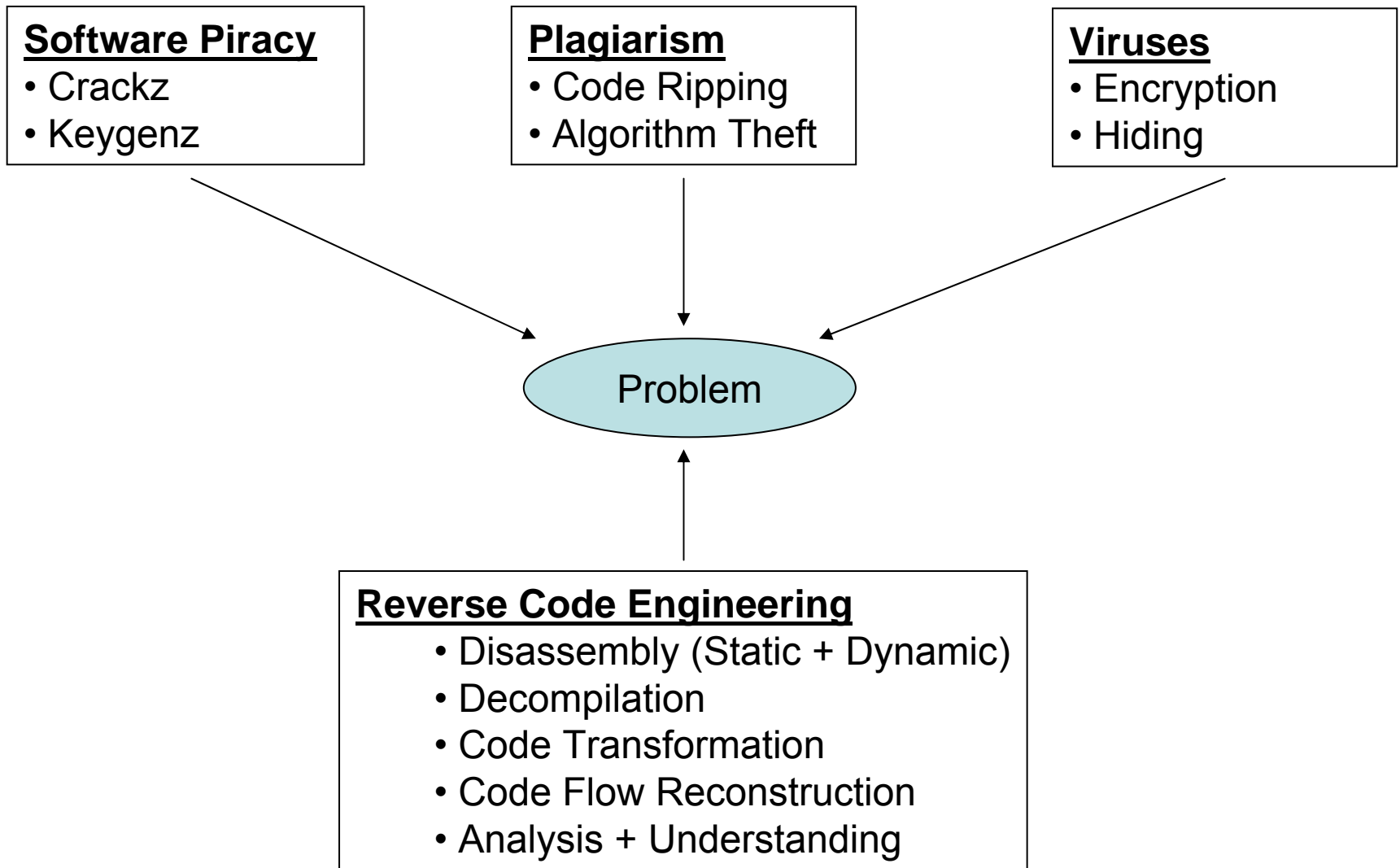
Talk at RECON2005

Thorsten Schneider

Overview

- Introduction
- Petri Nets – Overview
- Petri Nets – Example
- Protection by Obscurity and Obfuscation
- Example: Protection with Petri Nets
- Discussion
- Results

Known Problems



One Method of Resolution

Decrease Code Understanding

- Hardening through Obscurity
- Hardening through Obfuscation
- Hardening through Complexity
- Manipulation of Code Flow Graphs
- Manipulation of Information Flow Graphs
- ...
- Petri Nets!

```
> From: "Jim Coplien" <cope@research.bell-labs.com>
> Date: Tue, 22 Dec 1998 13:03:56 -0600
> To: Lalita Jagadeesan <lalita@research.bell-labs.com>, god, tball
> Subject: a program for your flow and testing tools
>
> /*
> * seriously -- run it :- )
> */
> #include <stdio.h>
> main(t,_,a)
> char *a;
> {
> return!0<t?t<3?main(-79,-13,a+main(-87,1-_,main(-86,0,a+1)+a):
> 1,t<_?main(t+1,_,a):3,main(-94,-27+t,a)&&t==2?_<13?
> main(2,_,+1,"%s %d %d\n"):9:16:t<0?t<-72?main(,t,
> "@n'+,#/*}w+/w#cdnr/+,}{r/*de}+,*{*,/w{%,/w#q#n+,#{l+,/n{n+,/+n+,/#\
> ,#q#n+/,+k#,*+,/r :d*3,}{w+K w'K:'+}e#';dq#'\ \
> q#'+d'K#!/+k#;q#r}eKK#}w'r}eKK{n}|/#;#q#n')}{#}w')}{n}|/+n';d}rw' i;#\
> }n}|/n{n#'; r{#w'r nc{n}|/#{l,+K {rw' ik{;{n}|/w#q#n'wk nw' \
> iwk{KK{n}|/w{%'###w#'# i; :{n}|/*{q#l'd;r}{nlwb!/*de}'c \
> ;;{nl'~}{rw}'/+,}##*}#nc,'#nw}'/+kd'+e}+;#rdq#w! nr'/' )}+}{rl#'{n'')# \
> }'+}##(!!/"
> :t<-50?_==*a?putchar(31[a]):main(-65,_,a+1):main((*a=='')+t,_,a+1)
> :0<t?main(2,2,"%s"):a=='/'||main(0,main(-61,*a,
> "!ek;dc i@bK'(q)-[w]*%n+r3#l,}{:\nuwloca-O;m .vpbks,fxntdCeghiry"),a+1);
> }
```

Petri Nets: Overview

- A *Petri Net* is a method, to represent *processes* in an abstract way
- Uninteresting Processes (for us):
 - Factory work flows
 - Business flows
 - Communication flows (protocols)
 - Device controls
 - Handicraft manuals
 - Biological Pathways (Bioinformatics)
- Interesting Processes:
 - Software Development Processes
 - Software Algorithms
 - Registration Number Schemes
- Petri Nets are graphs
- Advantage: Multithreaded processing!

Petri Net Types

- Discrete Petri Net
- Autonomous Petri Net
- Non-Autonomous T-Timed and P-Timed Petri Net
- Stochastic Petri Net
- Continuous Petri Net
- CSPN (Constant Speed Petri Net)
- VSPN (Variable Speed Petri Net)
- Hybrid Petri-Net

Petri Nets: Formal Definition

A Petri Net is a 6-Tupel (S, T, F, K, W, M_0) with:

- S: non-empty set of locations (Places)
 - T: non-empty set of Transitions
 - F: non-empty set of edges (Arcs)
-
- K: Capacity of Places for Tokens
 - W: Weight of Edges
 - M_0 : Startup Marking

$S = \{s_1, s_2, \dots, s_{|S|}\}$

$T = \{t_1, t_2, \dots, t_{|T|}\}$

$F \subseteq (S \times T \cup T \times S)$

$K: S \rightarrow \mathbb{N} \setminus \{0\}$

$W: F \rightarrow \mathbb{N} \setminus \{0\}$

$M_0: S \rightarrow \mathbb{N}$

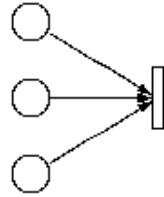
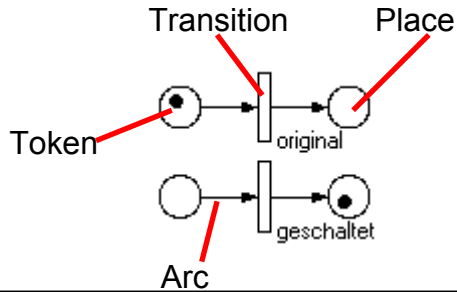
$0 \leq M(s) \leq K(s)$

$S \cap T = \emptyset$

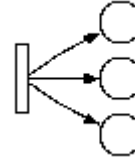
Very Simplified:

- 3 different objects: Places, Transitions and Tokens
 - No object (Places, Transitions) can belong to both sets
 - Between Places and Arcs there *might* be a relation (F)
 - Can simulate „something“
-
- So a Petri Net is a kind of a runnable process graph

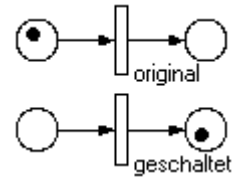
Petri Nets: Basics



Before Transition



After Transition



Firing Process

- Firing of *Transitions* changes network *Tokens* (located at the *Places*)
- Only one *Transition* can fire
- If so, a *Transition* removes as many *Tokens*, as the *Weight* of the *Arcs* defines
- The *Places* after the *Transition* receive the *Tokens*

- The *Places* *before* the *Transition* need to have enough *Tokens*
- The *Places* *after* the *Transition* need to have enough empty space for new *Tokens*.

- A *Transition* which is *able* to fire is called *activated*
 - But: *activated* does *not* mean that it is really fired!

- A *Petri Net* containing no *activated Transitions* is a *dead Petri Net*

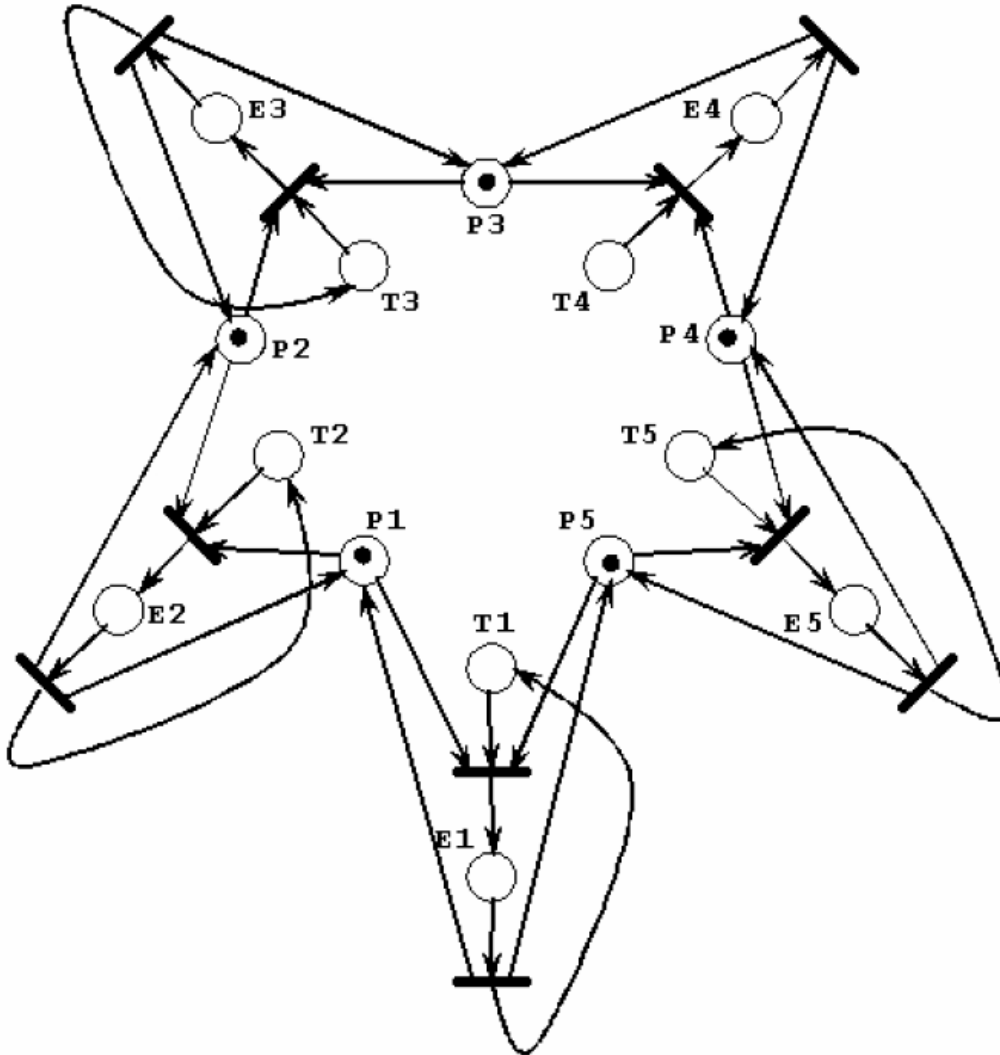
Petri Nets might save lives!

The Five Chinese Sages Problem [Dijkstra]:

Five Chinese sages are sitting at the circle table and have a dinner. Between of each two sages is only one stick. But for eating each of them needs two sticks in a moment. Obviously, if all sages takes sticks from the left side and waiting sticks from right side they all will die through starvation (dead loop).

DIJKSTRA, E.W.: *Co-operating sequential processes*. In *Programming Languages*, F. Genuys, Ed., 43–112, 1968.

Petri Nets might save lifes!



- Places $P_1 \dots P_5$ introduce sticks
- All sticks are on the table at the first moment
→ each place has a marker inside
- Transitions T_i and E_i introduce sages states:
 - $T_i \rightarrow$ sage_{*i*} thinks
 - $E_i \rightarrow$ sage_{*i*} eats.
- To pass from M_i state (obviously, no one can satisfy his hunger through his thoughts) to E_i state, both sticks (on left and right sides) must be on the table at one moment.

Petri Nets: Conflicts

- **Pre-Conflict:**

- 2 Transitions need the *same* Token to Fire
- Both Transitions are *activated*, but *only 1* can fire
- This is no erroneous Petri Net, but models the decision between 2 alternatives

- **Post-Conflict:**

- Similar to Pre-Conflict
- 2 Transitions produce Tokens, but the capacity of the Places is too low for all Tokens
- Solution is dependant on conflict strategy

- **Confusion:**

- Is a doubled conflict
- One Transition conflicts with two different Transitions

Protection by Obscurity and Obfuscation

P = Program, T = Transformation, S = Source Code

Given a program P and an obfuscated program P':

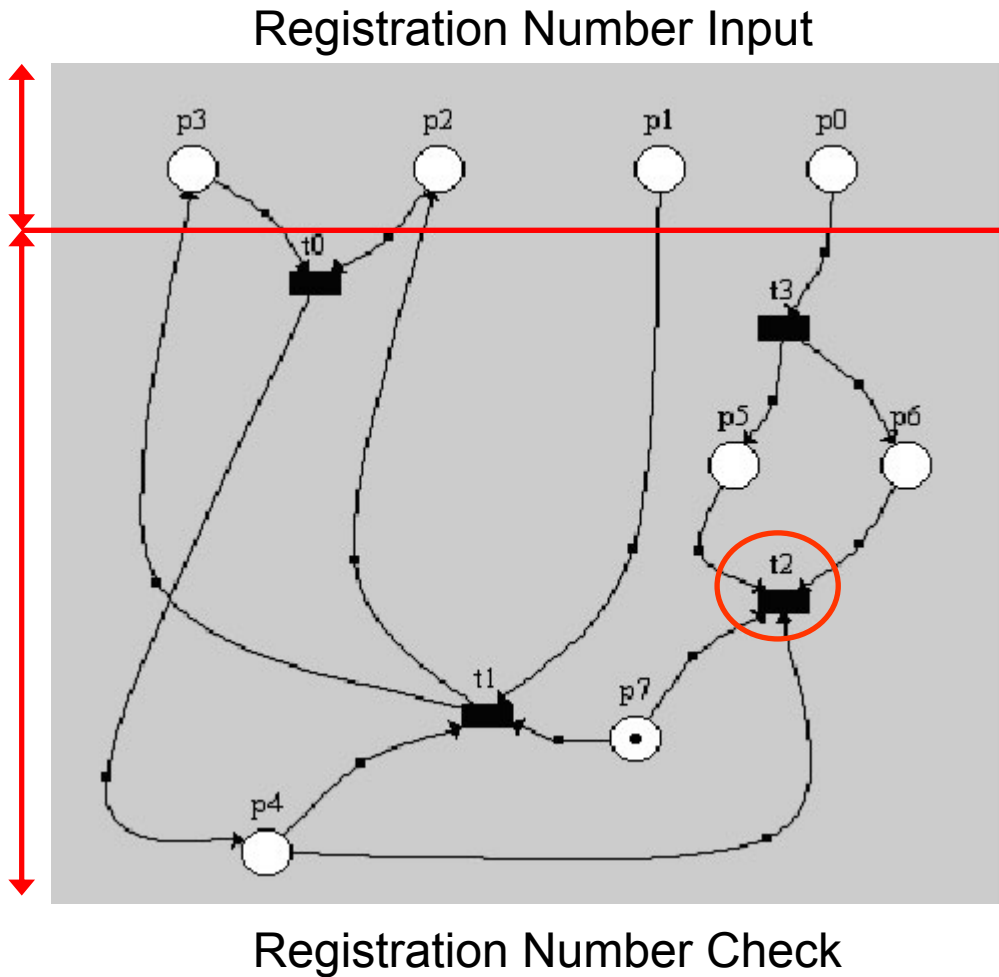
- P' has the same observable behavior as P, i.e., the transformations are semantics-preserving.
- The obscurity of P' maximized, i.e., understanding and reverse engineering P' will be strictly more time-consuming than understanding and reverse engineering P.
- The resilience of each transformation $T_i(S_j)$ is maximized, i.e., it will be difficult to construct an automatic tool to undo the transformations
- The stealth of each transformation $T_i(S_j)$ is maximized, i.e., the statistical properties of S'_j are similar to those of S_j .
- The cost (the execution time/space penalty incurred by the transformations) of P' is minimized.

Protection by Obscurity and Obfuscation

Code obfuscation is very similar to code optimization, except:

- with obfuscation, we are maximizing obscurity while minimizing execution time
- with optimization, we are just minimizing execution time.

Example: Protection with Petri Nets



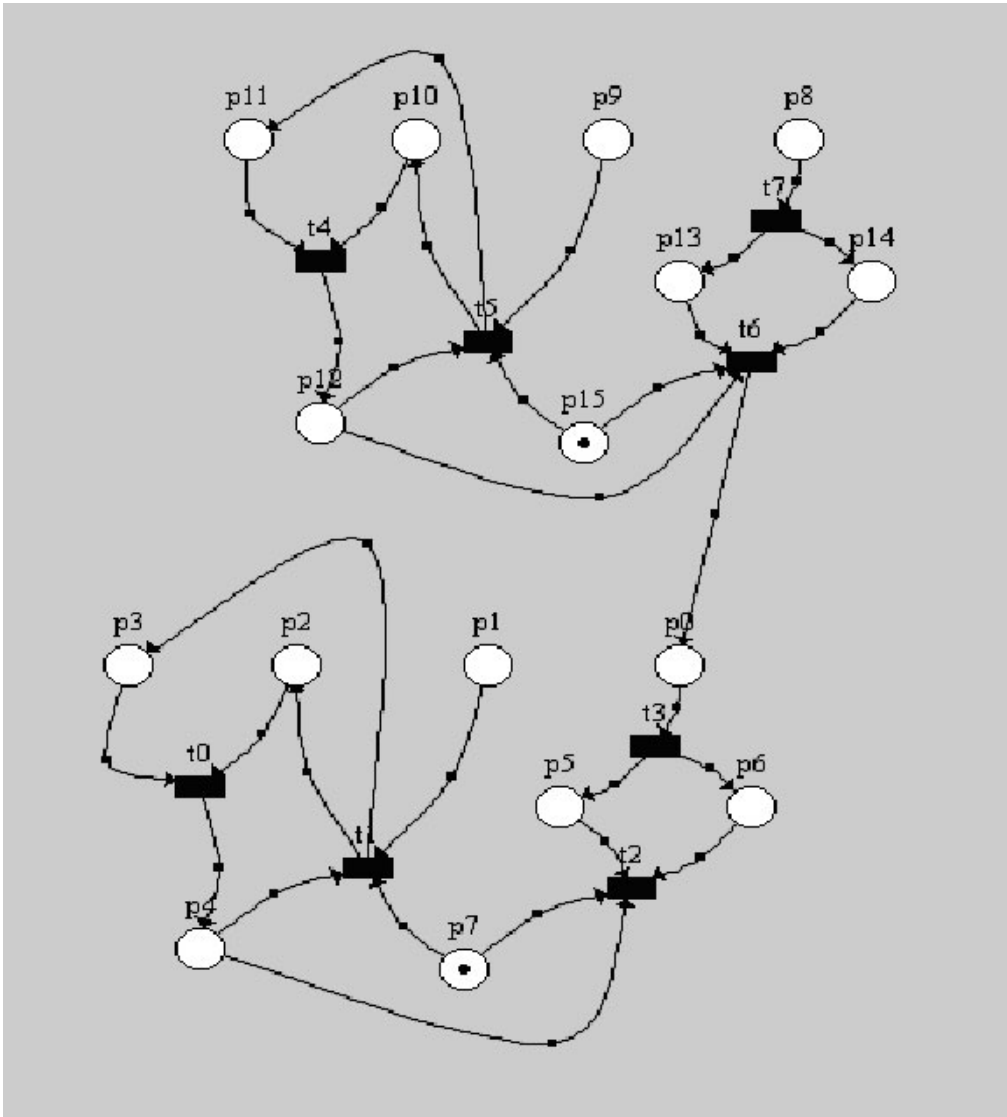
- Markable: P_0 to P_3
- P_7 is pre-marked !
- Lowest Priority to T_2
- Application gets registered when T_2 is fired

Solution:

$$\begin{aligned} P_0 &= 1 \\ P_1 &= 0 \\ P_2 &= 1 \\ P_3 &= 1 \end{aligned}$$

Solution Key Space: $2^4 = 16$ tries
→ Simple for Bruteforce
→ Simple for Brain

Example: Protection with Petri Nets



- Increasing Complexity
 - Decreasing Understanding
- Simple Copy & Paste possible
 - P_1 not possible
 - Reachability Problem of T_2
- Introduction of new Places and Transitions might be necessary

Solution:

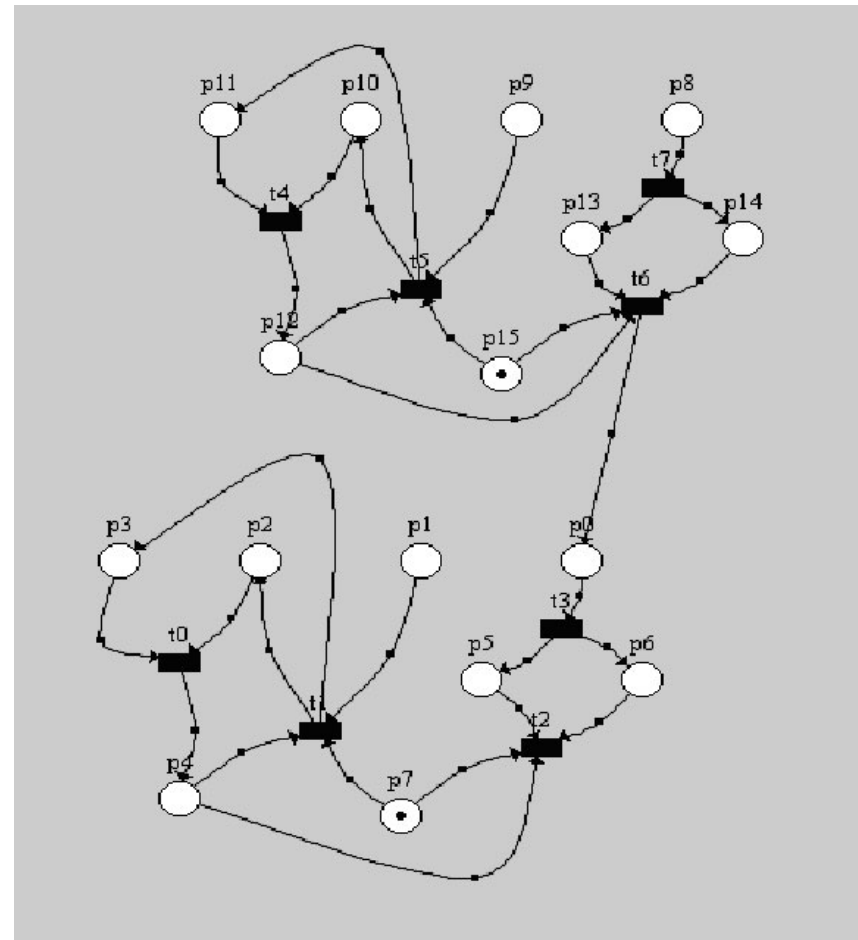
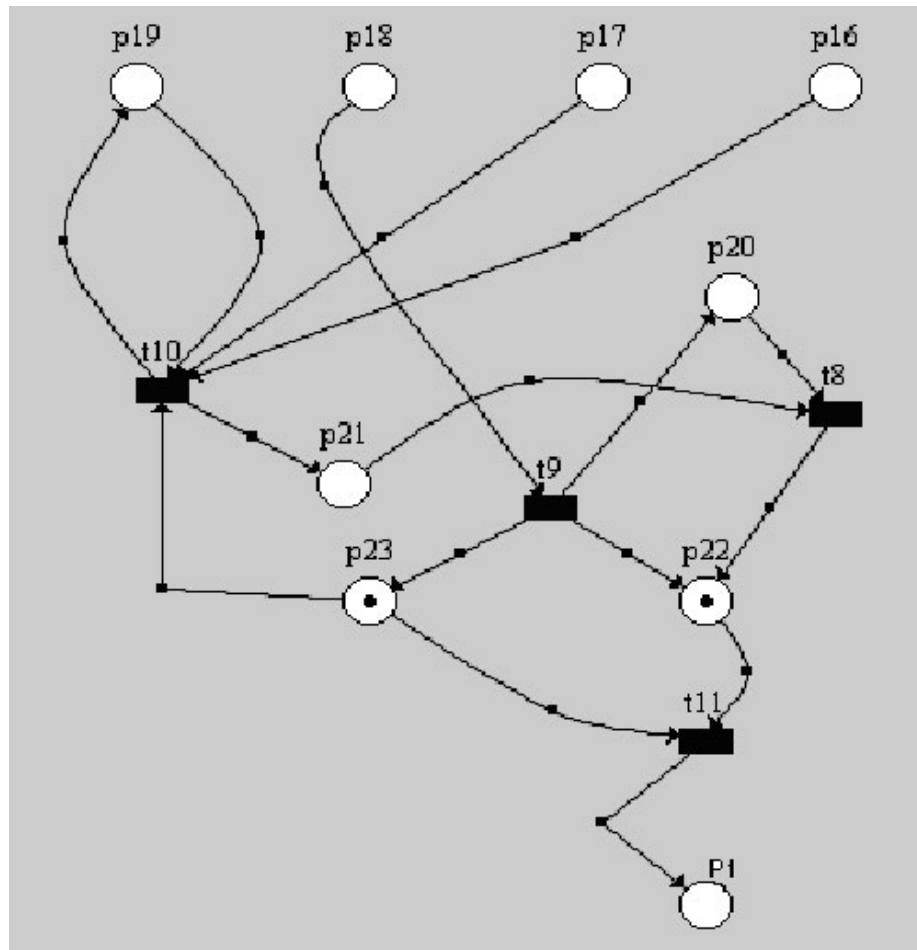
$$P_1 = 0, P_2 = 1, P_3 = 1, P_8 = 1, \\ P_9 = 0, P_{10} = 1, P_{11} = 1$$

Solution Key Space:

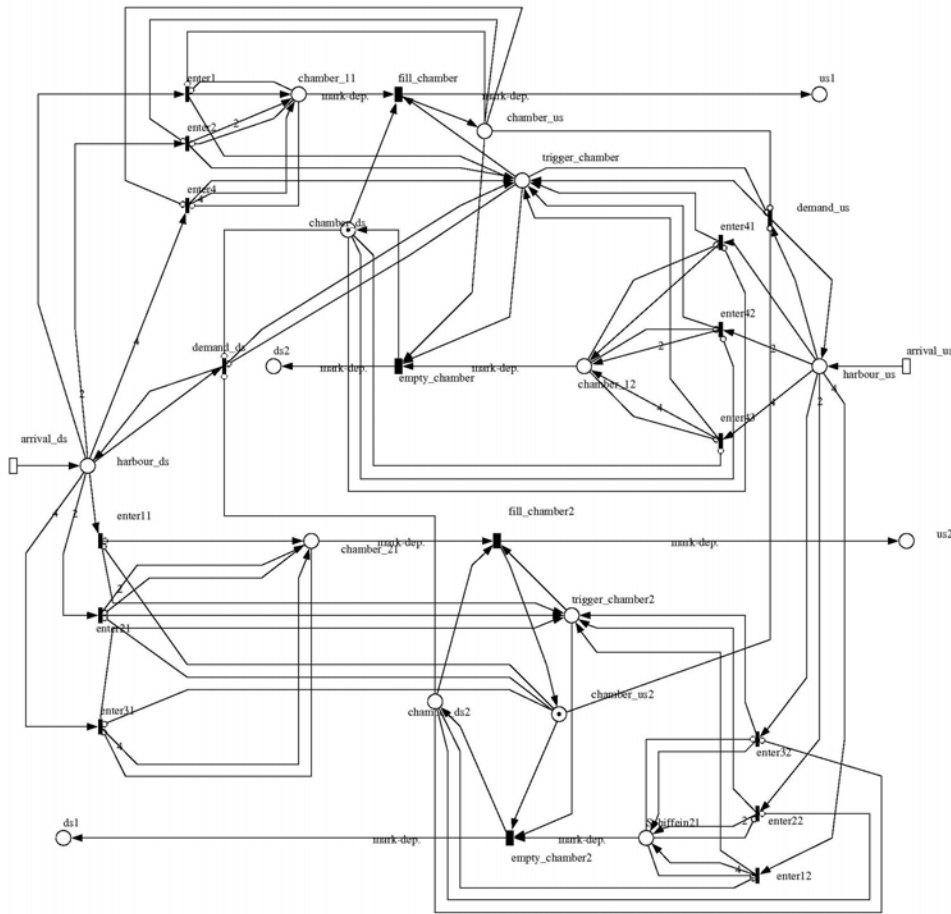
- $2^{16} = 65536$ tries
- Harder for Bruteforce
 - Harder for Brain

Example: Protection with Petri Nets

Changed Sub Petri Net to attach at Place P_1



Example: Protection with Petri Nets

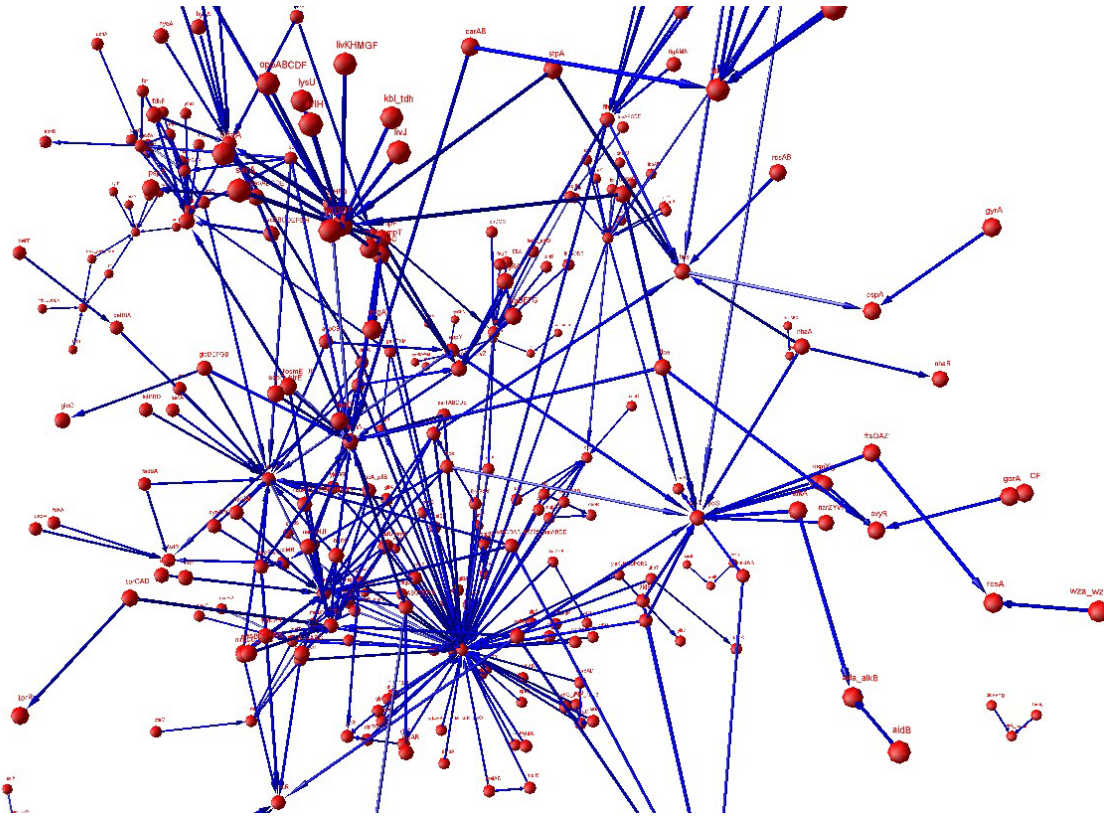


Now imagine:

- Each Transition is a Thread
- Each Place is a Thread
- Each Thread is protected different:
 - Anti-Debugging
 - Anti-Disassembly
 - Self-Encryption
 - Watchdogs
 - ... and much more ...
- Each Arc is an encrypted communication protocol
- Each Token is an encrypted Object

How would you analyse this ?

Example: Protection with Petri Nets



Now imagine:

1. You have a disassembly
2. You try to reconstruct the Petri Net from disassembly or debugging
3. You need to trace or debug *parallel* processes to understand the parallel processes
4. And all these ugly protection tricks within each Thread!

How would you analyse this ?

Problems of Complexity

- Research only focus on *decreasing* complexity
 - Many research groups
 - Much research
- We want to *increase* complexity
 - No research groups
 - No research yet
 - But (!): decreasing complexity can be inverted!
 - But: No algorithms yet

Discussion: Pro

- High complexity
- High obscurity
- Reconstruction of Petri Net from binary code is hard up to impossible
- Protection is hard up to impossible to understand

Discussion: Contra

- Once reconstructed, it is possible to simplify the Petri Net
- Once simplified, it is possible to run reduced Bruteforce Attacks
- Once bruteforced, it is possible to get a valid key or Keygen
- Protection still breakable (e.g. Patching) at the Input Layer of the protection
- Development of complex Petri Nets is very time consuming, no automatism yet
- Implementation very time-consuming, no automatism yet

Results

- Petri Nets are an efficient way to obscure and to complex processes
- Resistant against Bruteforcing
- But: Once analysed, they can be simplified
- Example source and binary available
- Fact: all software protection schemes have been cracked
- Fact: If a code is runnable, you can crack it!
- Further research necessary!

Example Code with Online Disassembly

(<http://pvdasm.reverse-engineering.net/PVPHP.php>)

Acknowledgments

- Robert Airapetyan (Polytechnical University of Odessa)
- RECON 2005 Team
- The anonymous reviewers
- The audience

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<http://knoppix-re.reverse-engineering.net>

Questions ?

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